

Coreflective semilocalizing subcategories in triangulated categories with models

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Reflection and coreflections

- A full subcategory \mathcal{L} of a category \mathcal{T} is **reflective** if the inclusion $\mathcal{L} \hookrightarrow \mathcal{T}$ has a left adjoint $\mathcal{T} \rightarrow \mathcal{L}$.
- The composite $L: \mathcal{T} \rightarrow \mathcal{T}$ is called a **reflection** or a **localization** onto \mathcal{L} .
- There is a natural transformation $l: \text{Id} \rightarrow L$ such that $Ll: L \rightarrow LL$ is an isomorphism, lL is equal to Ll , and, for each X , the morphism $l_X: X \rightarrow LX$ is initial in \mathcal{T} among morphisms from X to objects in \mathcal{L} .
- Similarly, a full subcategory \mathcal{C} of \mathcal{T} is **coreflective** if the inclusion $\mathcal{C} \hookrightarrow \mathcal{T}$ has a right adjoint.
- The composite $C: \mathcal{T} \rightarrow \mathcal{T}$ is called a **coreflection** or a **colocalization** onto \mathcal{C} .
- There is a natural transformation $c: C \rightarrow \text{Id}$ such that Cc is an isomorphism, cC is equal to Cc , and for each X , the morphism $c_X: CX \rightarrow X$ is terminal among morphisms from objects in \mathcal{C} .

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Weak reflections

- A full subcategory \mathcal{L} of a category \mathcal{T} is called **weakly reflective** if for every object X of \mathcal{T} there is a morphism $l_X: X \rightarrow X^*$ with X^* in \mathcal{L} and such that the function

$$\mathcal{T}(l_X, Y): \mathcal{T}(X^*, Y) \longrightarrow \mathcal{T}(X, Y)$$

is surjective for all objects Y of \mathcal{L} .

- Every morphism from X to an object of \mathcal{L} factors through l_X , not necessarily in a unique way.
- If such a factorization is unique for all objects X , then the morphisms $l_X: X \rightarrow X^*$ for all X define together a reflection, so \mathcal{L} is then reflective.

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Closure properties for triangulated categories

Let \mathcal{T} be a triangulated category with products and coproducts. Let Σ be the shift or suspension operator. Triangles in \mathcal{T} will be denoted by

$$X \longrightarrow Y \longrightarrow Z \longrightarrow \Sigma X,$$

Let \mathcal{S} be a full subcategory of \mathcal{T} . We define the following closure properties in \mathcal{S} :

- (i) **Fibres**: If Y and Z are in \mathcal{S} then X is in \mathcal{S} .
- (ii) **Cofibres**: If X and Y are in \mathcal{S} then Z is in \mathcal{S} .
- (iii) **Extensions**: If X and Z are in \mathcal{S} then Y is in \mathcal{S} .

Definition

A full subcategory of \mathcal{T} is called **semilocalizing** if it is closed under coproducts, (ii) and (iii). It is called **localizing** if it is closed under coproducts, (i), (ii) and (iii). Dually, a full subcategory is called **semicolocalizing** if it is closed under products, (i) and (iii). It is called **colocalizing** if it is closed under products, (i), (ii) and (iii).

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Closure properties for triangulated categories

Let \mathcal{T} be a *tensor triangulated category*, i.e., \mathcal{T} has a closed symmetric monoidal structure with a unit object \mathcal{S} , tensor product \wedge and internal hom $F(-, -)$, compatible with the triangulated structure and such that $\mathcal{T}(X, F(Y, Z)) \cong \mathcal{T}(X \wedge Y, Z)$ naturally in all variables. In this case we have in addition the following closure properties in \mathcal{S} :

- (iv) **Tensor**: If X is in \mathcal{S} , then so is $X \wedge W$ for all W in \mathcal{T} .
- (v) **Internal hom**: If X is in \mathcal{S} , then so is $F(W, X)$ for all W in \mathcal{T} .

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A full subcategory of \mathcal{T} is called a **localizing ideal** if it is a localizing subcategory and closed under (iv). It is called a **colocalizing coideal** if it is a colocalizing subcategory and it is closed under (v).

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Closure properties for triangulated categories

- A reflection L on \mathcal{T} will be called **semiexact** if the subcategory of L -local objects is semicolocalizing, and **exact** if it is colocalizing.
- Dually, a coreflection C will be called **semiexact** if the subcategory of C -colocal objects is semilocalizing and **exact** if it is localizing.
- A (co)reflection is exact if and only if it preserves all triangles (or, equivalently, commutes with the suspension operator)

Examples

- Bousfield–Farjoun **localizations** are reflections.
- **Cellular approximations** are coreflections.
- **Nullifications** are semiexact reflections.
- **Homological localizations** are exact reflections

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Orthogonality and semiorthogonality

For a class of objects \mathcal{D} in \mathcal{T} we introduce the following notation:

$${}^{\perp}\mathcal{D} = \{X \mid \mathcal{T}(X, \Sigma^k D) = 0 \text{ for all } D \in \mathcal{D} \text{ and } k \leq 0\}$$

$$\mathcal{D}^{\perp} = \{Y \mid \mathcal{T}(\Sigma^k D, Y) = 0 \text{ for all } D \in \mathcal{D} \text{ and } k \geq 0\}$$

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Orthogonality and semiorthogonality

Theorem

In every triangulated category \mathcal{T} there is a bijective correspondence between semiexact reflections and semiexact coreflections such that, if a reflection L is paired with a coreflection C then:

- (i) *For every X , the morphisms $l_X: X \rightarrow LX$ and $c_X: CX \rightarrow X$ fit into a triangle*

$$CX \longrightarrow X \longrightarrow LX \longrightarrow \Sigma CX.$$

- (ii) *The class \mathcal{L} of L -local objects coincides with the class of C -acyclics, and the class \mathcal{C} of C -colocal objects coincides with the class of L -acyclics.*
- (iii) *The class \mathcal{C} is equal to ${}^{\perp}\mathcal{L}$, and \mathcal{L} is equal to \mathcal{C}^{\perp} .*
- (iv) *L is exact if and only if C is exact. In this case, $\mathcal{C} = {}^{\perp}\mathcal{L}$ and $\mathcal{L} = \mathcal{C}^{\perp}$.*

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- (ii) *The class \mathcal{L} of L -local objects coincides with the class of C -acyclics, and the class \mathcal{C} of C -colocal objects coincides with the class of L -acyclics.*
- (iii) *The class \mathcal{C} is equal to ${}^{\perp}\mathcal{L}$, and \mathcal{L} is equal to \mathcal{C}^{\perp} .*
- (iv) *L is exact if and only if C is exact. In this case, $\mathcal{C} = {}^{\perp}\mathcal{L}$ and $\mathcal{L} = \mathcal{C}^{\perp}$.*

Orthogonality and semiorthogonality

Theorem

In every triangulated category \mathcal{T} there is a bijective correspondence between semiexact reflections and semiexact coreflections such that, if a reflection L is paired with a coreflection C then:

- (i) *For every X , the morphisms $l_X: X \rightarrow LX$ and $c_X: CX \rightarrow X$ fit into a triangle*

$$CX \longrightarrow X \longrightarrow LX \longrightarrow \Sigma CX.$$

- (ii) *The class \mathcal{L} of L -local objects coincides with the class of C -acyclics, and the class \mathcal{C} of C -colocal objects coincides with the class of L -acyclics.*
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Orthogonality and semiorthogonality

Proposition

Let \mathcal{D} be any class of objects in a triangulated category with products and coproducts.

- (i) If $\text{scoloc}(\mathcal{D})$ is reflective, then $\text{scoloc}(\mathcal{D}) = ({}^{\perp}\mathcal{D})^{\perp}$, and if $\text{sloc}(\mathcal{D})$ is coreflective, then $\text{sloc}(\mathcal{D}) = {}^{\perp}(\mathcal{D}^{\perp})$.
- (ii) If $\text{coloc}(\mathcal{D})$ is reflective, then $\text{coloc}(\mathcal{D}) = ({}^{\perp}\mathcal{D})^{\perp}$, and if $\text{loc}(\mathcal{D})$ is coreflective, then $\text{loc}(\mathcal{D}) = {}^{\perp}(\mathcal{D}^{\perp})$.

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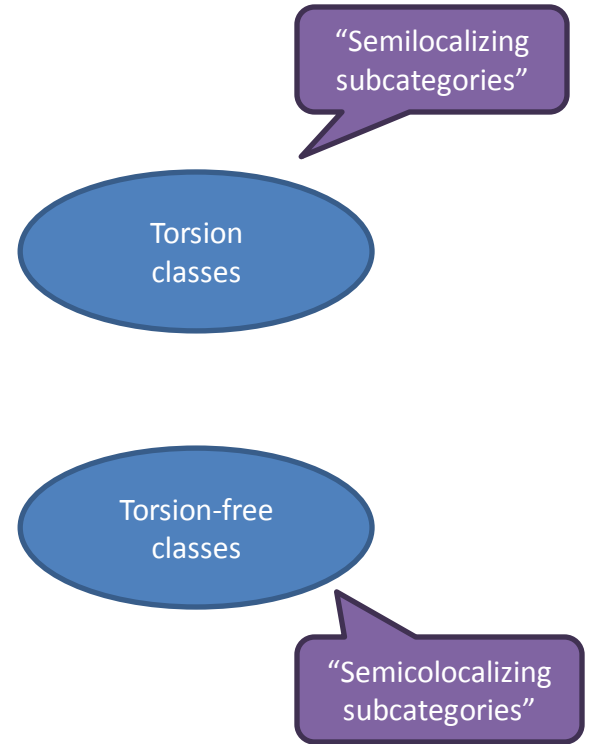
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Torsion theories
in abelian categories
(e.g., abelian groups)

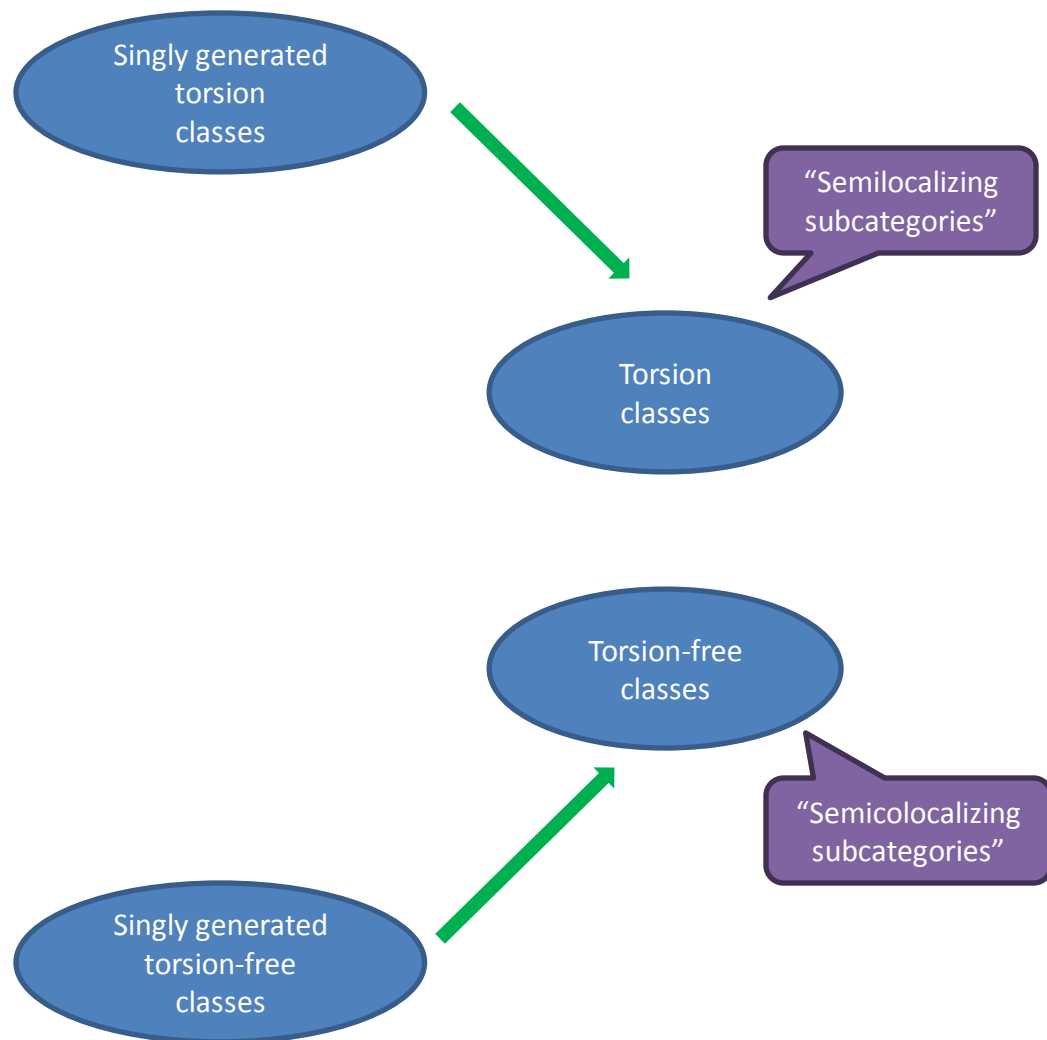
Torsion
classes

Torsion-free
classes

Torsion theories
in abelian categories
(e.g., abelian groups)



Torsion theories
in abelian categories
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Torsion theories
in abelian categories
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Hereditary
torsion
theories

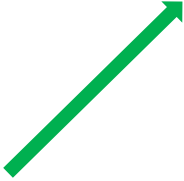
Singly generated
torsion
classes



Torsion
classes

“Semilocalizing
subcategories”

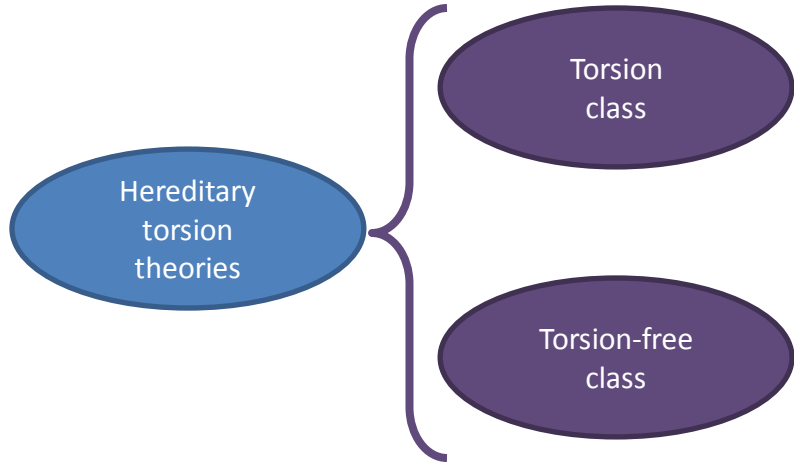
Torsion-free
classes



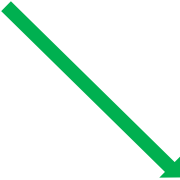
Singly generated
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Torsion theories
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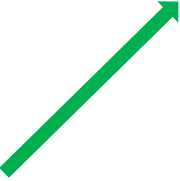
Singly generated torsion classes



Torsion classes

“Semilocalizing subcategories”

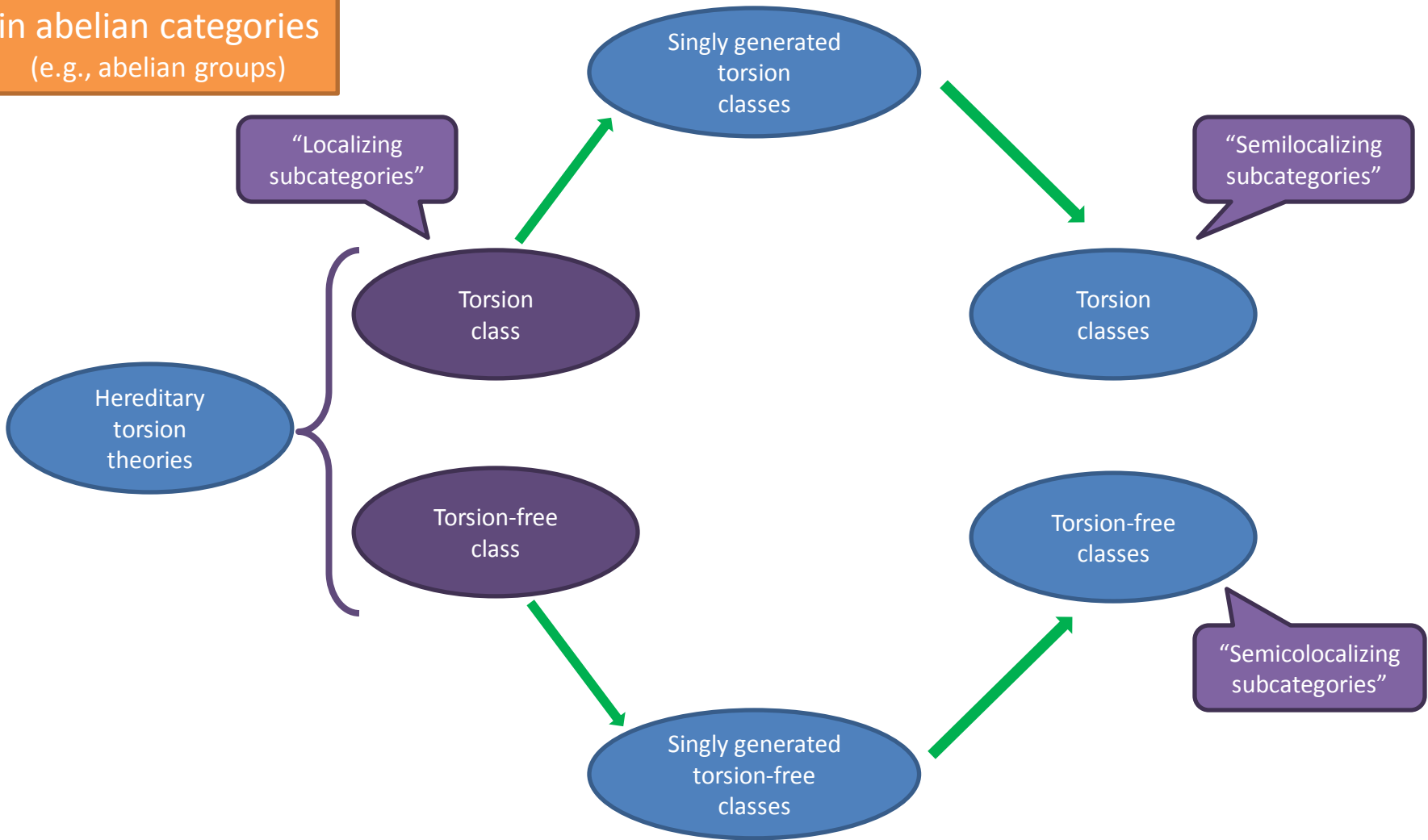
Singly generated torsion-free classes



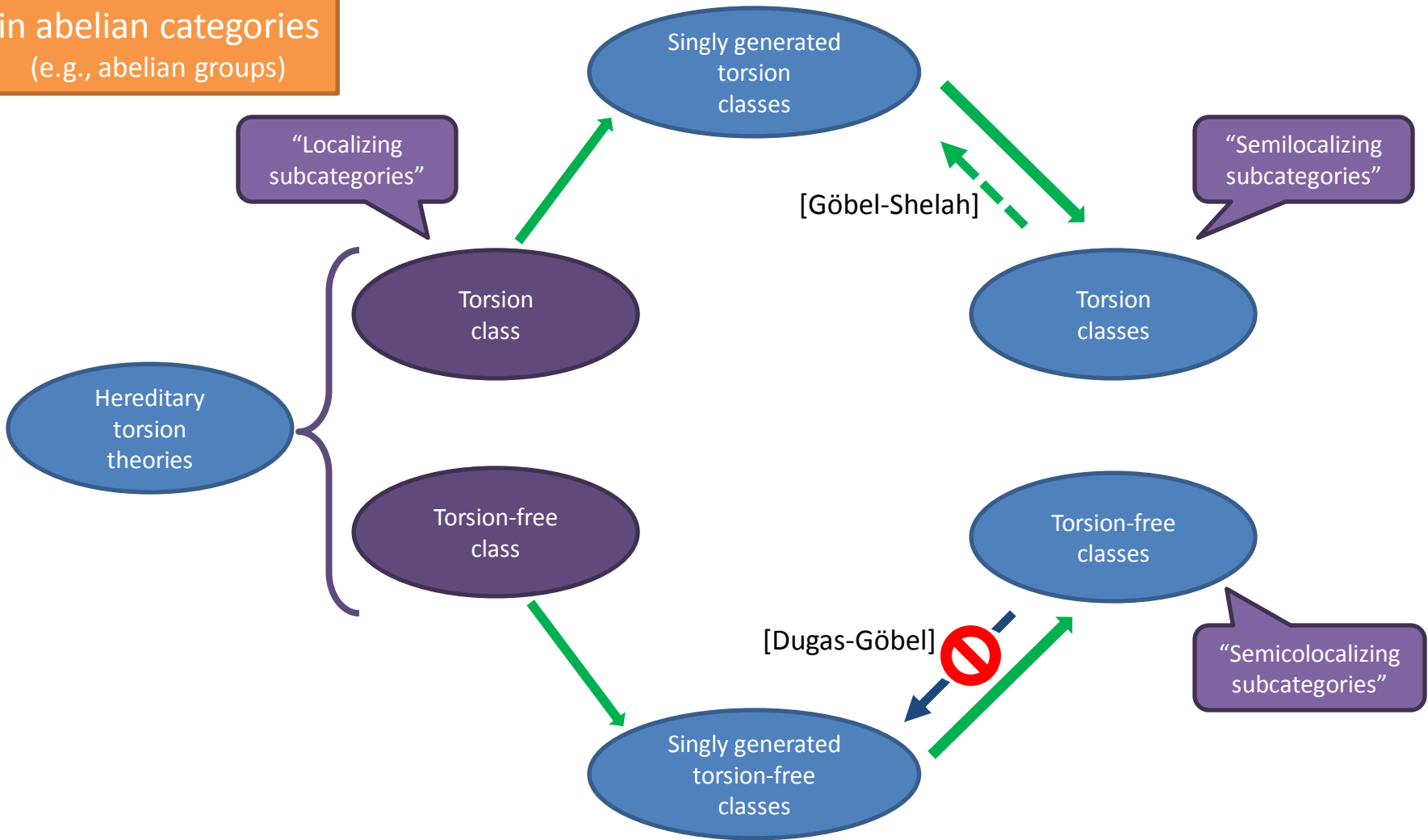
Torsion-free classes

“Semicolocalizing subcategories”

Torsion theories
in abelian categories
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----- Vopěnka's principle

Triangulated categories (with combinatorial models)

Semilocalizing
subcategories

Semicolocalizing
subcategories

Triangulated categories (with combinatorial models)

Closed
semilocalizing
subcategories

Semilocalizing
subcategories

Closed
semicolocalizing
subcategories

Semicolocalizing
subcategories

Triangulated categories (with combinatorial models)

Coreflective
semilocalizing
subcategories

Closed
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subcategories

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subcategories

Reflective
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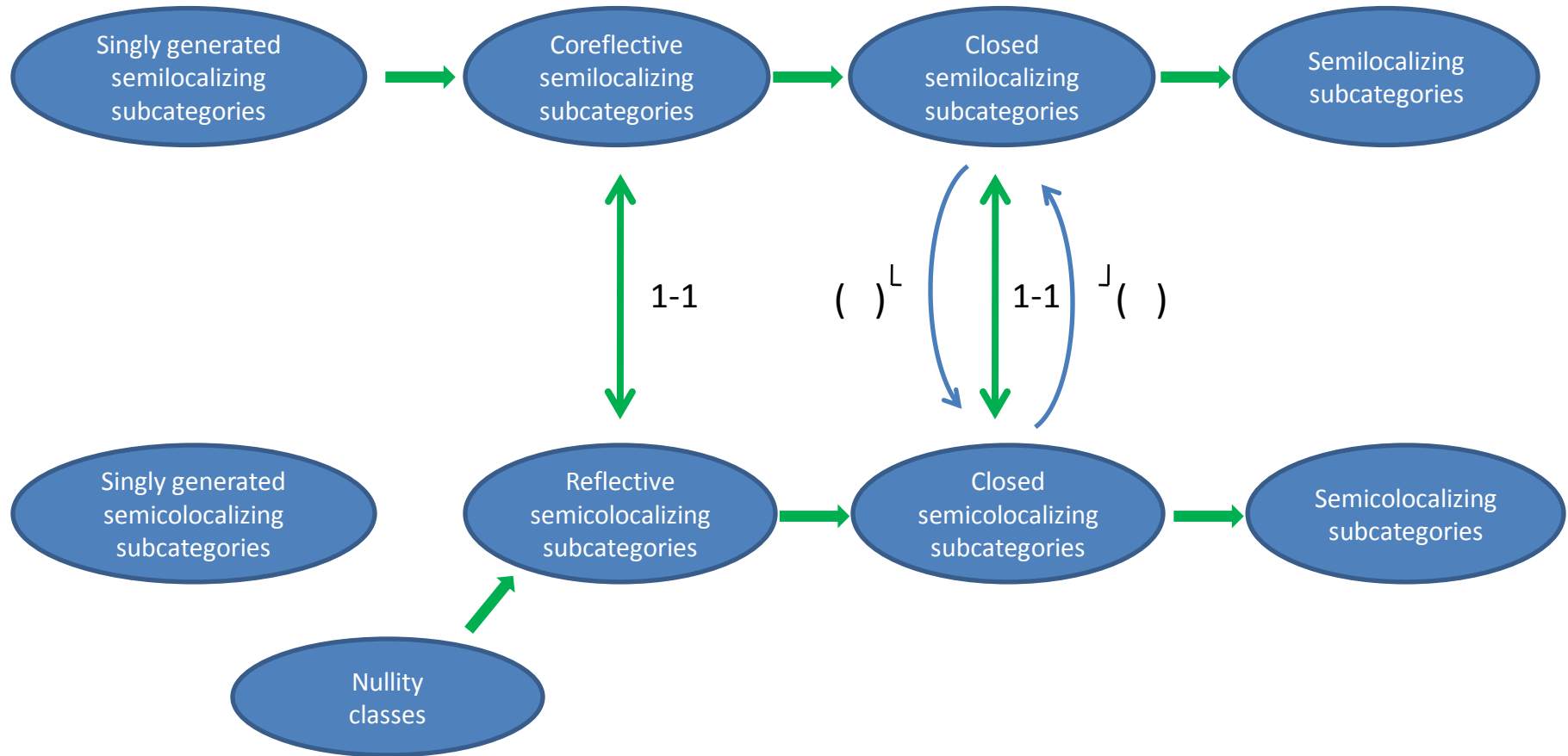
Reflective
semicolocalizing
subcategories

Closed
semicolocalizing
subcategories

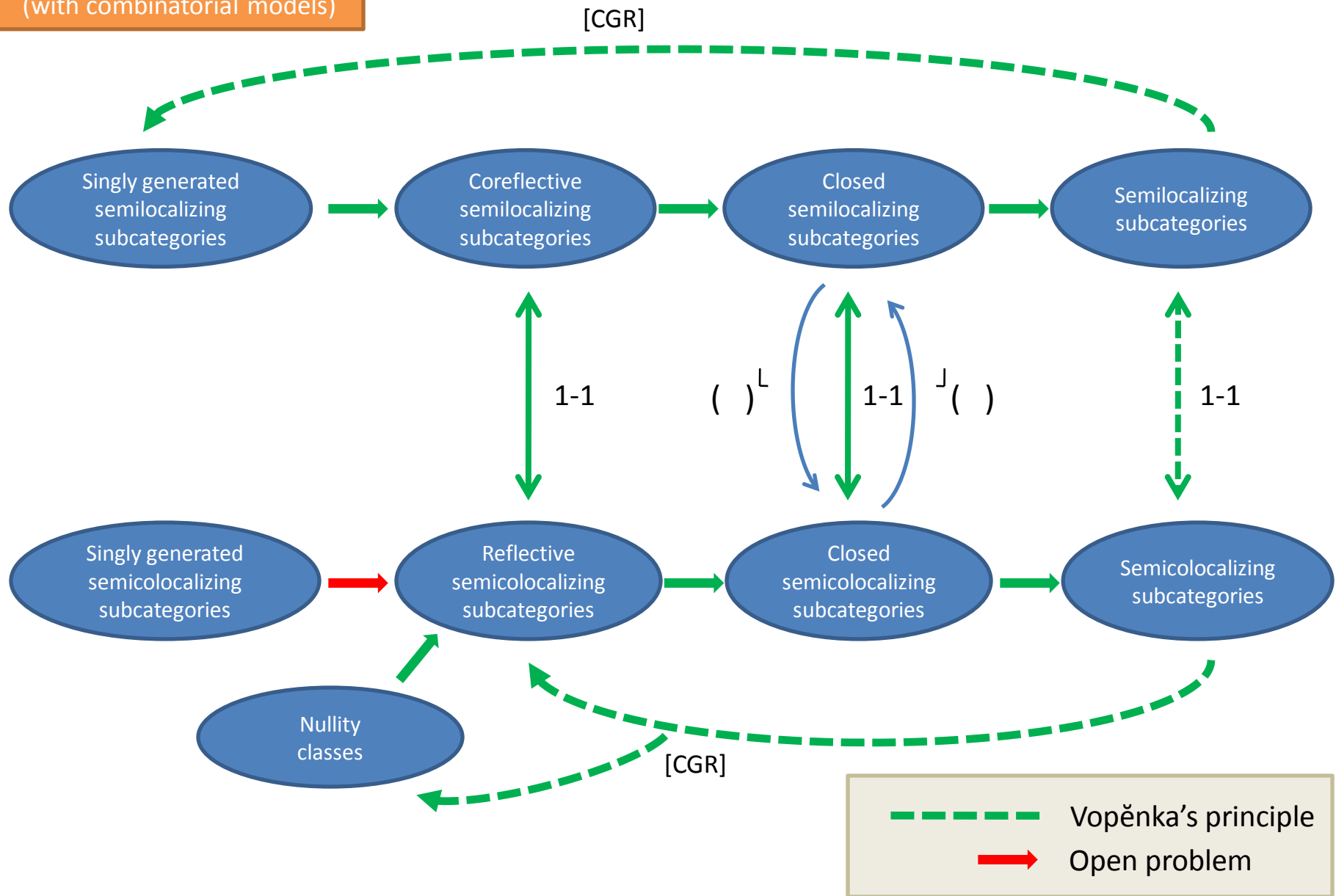
Semicolocalizing
subcategories

Nullity
classes

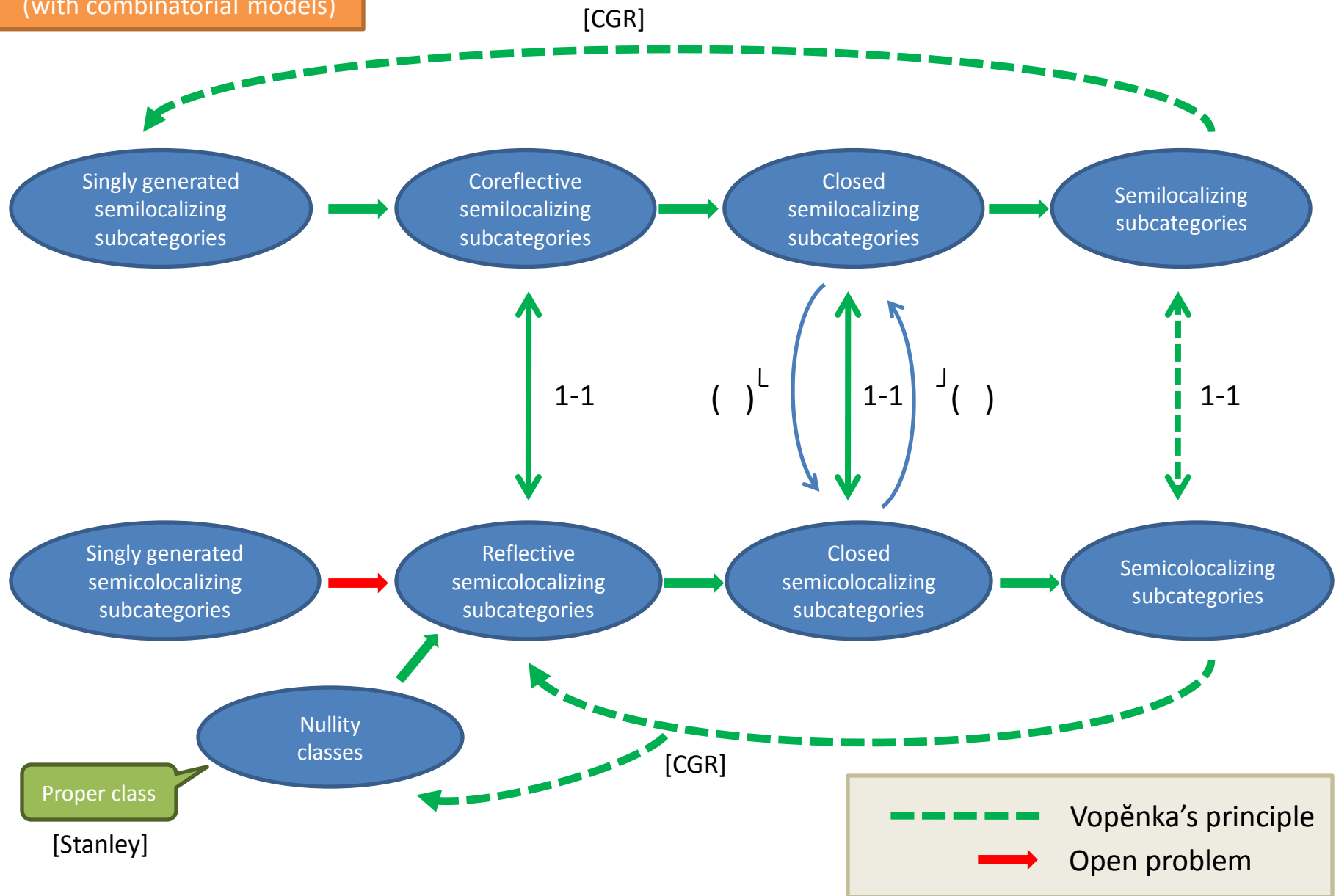
Triangulated categories (with combinatorial models)



Triangulated categories
(with combinatorial models)



Triangulated categories
(with combinatorial models)



Tensor triangulated
categories
(with combinatorial models)

Localizing
ideals

Colocalizing
coideals

Tensor triangulated
categories
(with combinatorial models)

Closed
localizing
ideals

Localizing
ideals

Closed
colocalizing
coideals

Colocalizing
coideals

Tensor triangulated
categories
(with combinatorial models)

Coreflective
localizing
ideals

Closed
localizing
ideals

Localizing
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Reflective
colocalizing
coideals

Closed
colocalizing
coideals

Colocalizing
coideals

Tensor triangulated
categories
(with combinatorial models)

Singly generated
localizing
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Coreflective
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Closed
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Localizing
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Singly generated
colocalizing
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coideals

Closed
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coideals

Colocalizing
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Homological
Bousfield classes

Cohomological
Bousfield classes

Tensor triangulated
categories
(with combinatorial models)

Singly generated
localizing
ideals

Coreflective
localizing
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Closed
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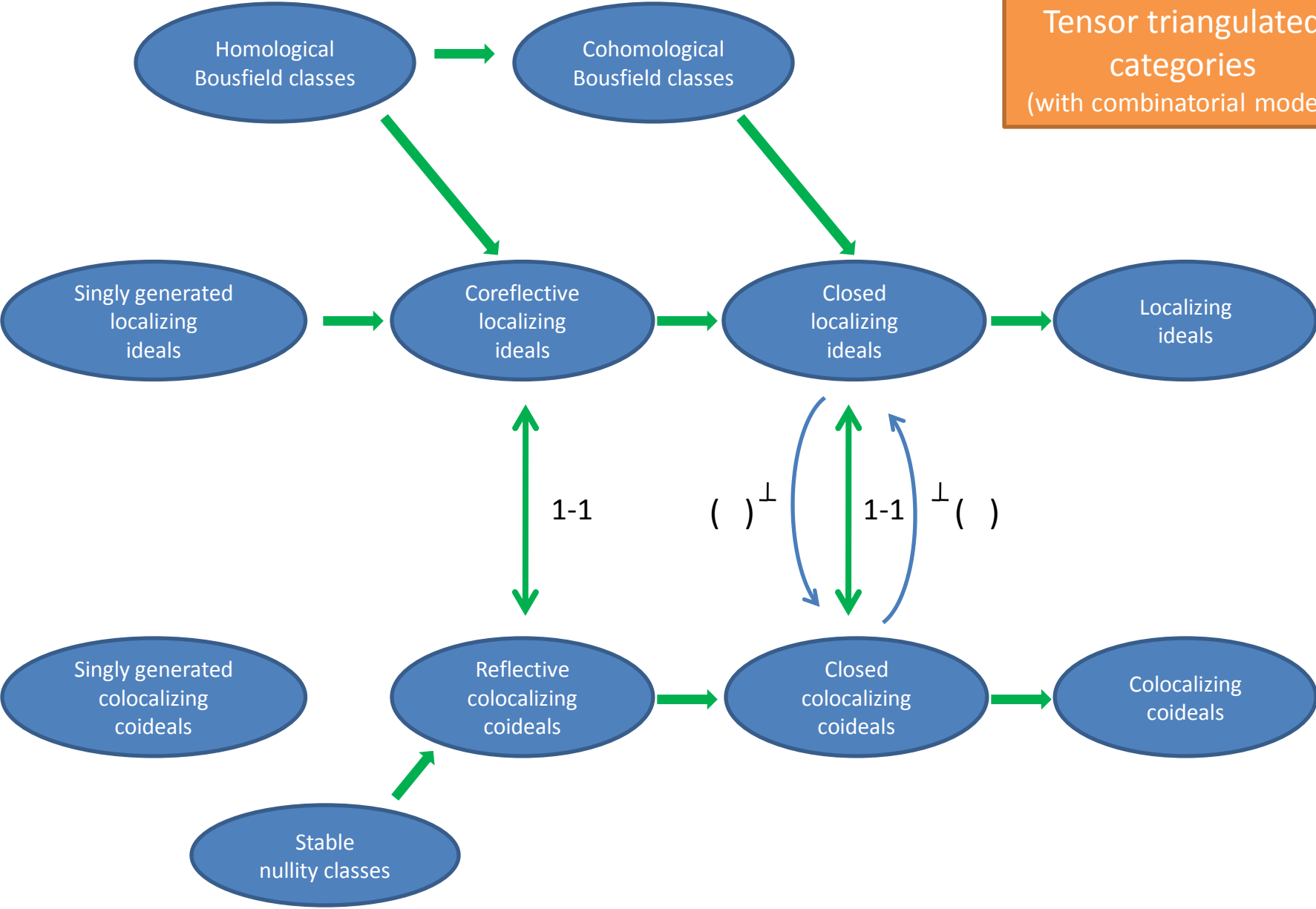
Reflective
colocalizing
coideals

Closed
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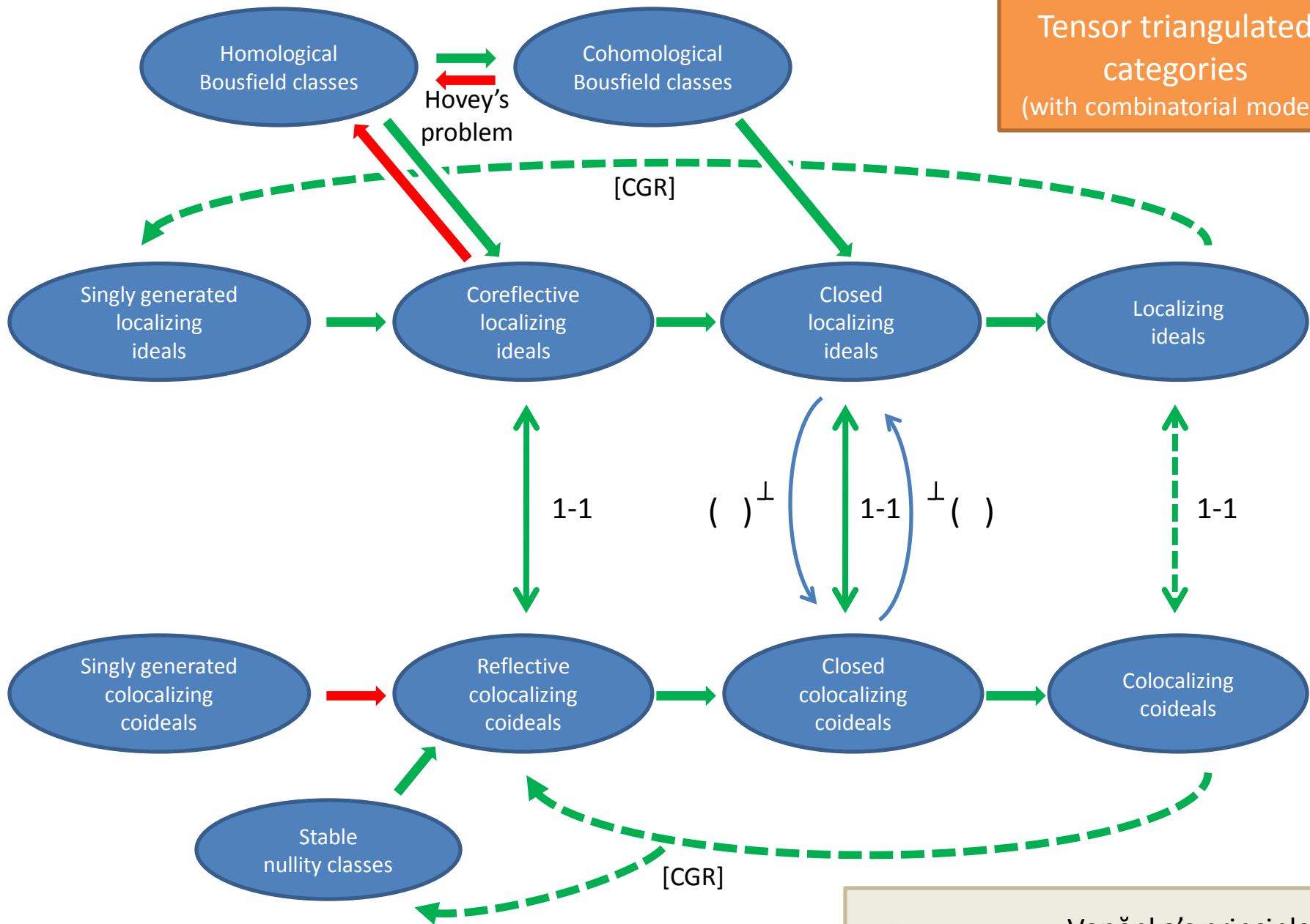
Colocalizing
coideals

Stable
nullity classes

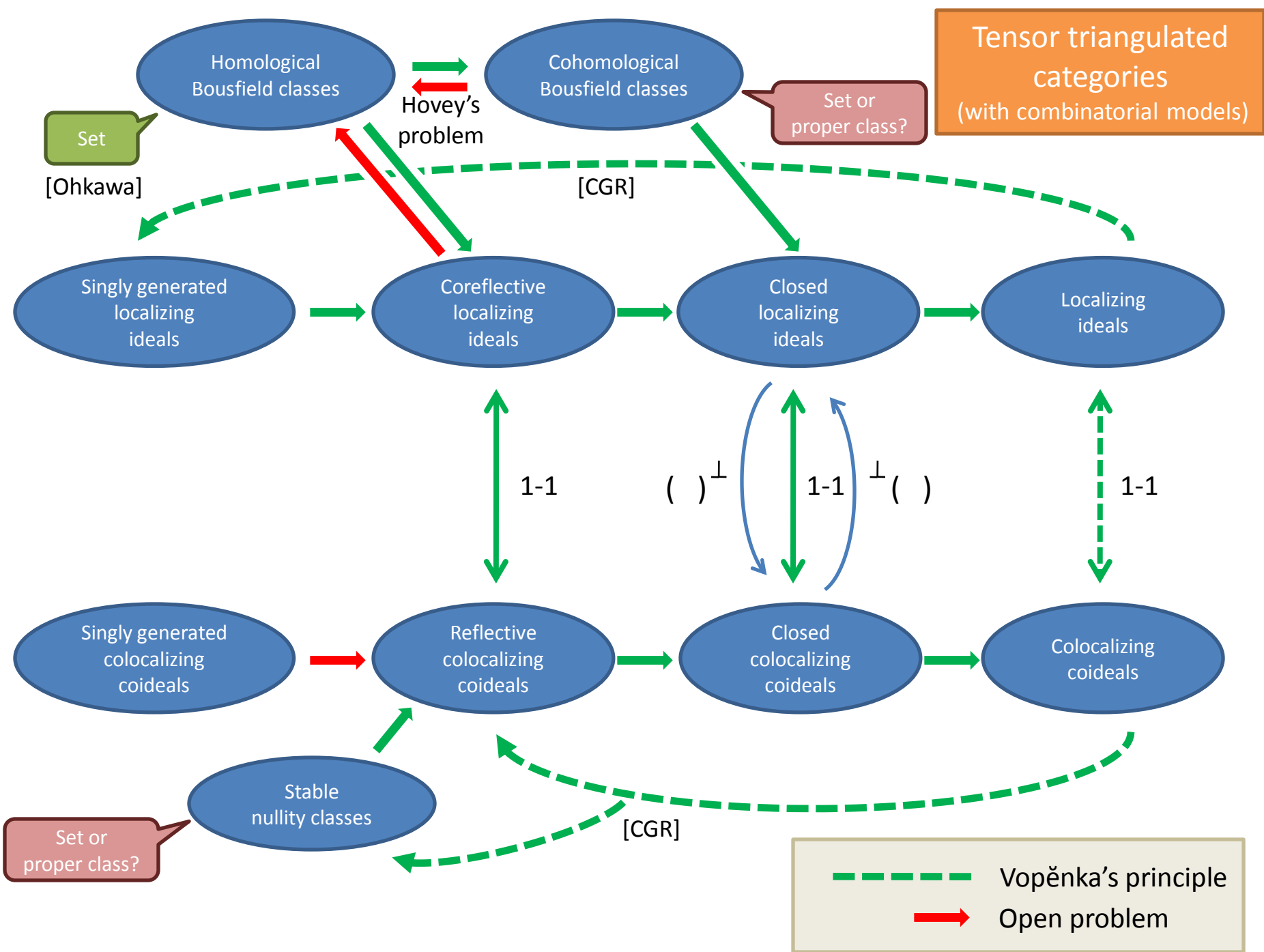
Tensor triangulated categories
(with combinatorial models)



Tensor triangulated categories
(with combinatorial models)



--- Vopěnka's principle
 → Open problem



Triangulated categories with combinatorial models

A model category \mathcal{K} is called

- (i) *Combinatorial* if it is locally presentable and cofibrantly generated.
- (ii) *Stable* if it is pointed and the suspension and loop operator are inverse equivalences on the homotopy category $\mathrm{Ho}(\mathcal{K})$. In this case $\mathrm{Ho}(\mathcal{K})$ is triangulated.

We are interested in triangulated categories that appear as homotopy categories of combinatorial stable (monoidal) model categories. Such triangulated categories are well-generated [Rosicky].

Examples

- The homotopy category of spectra
- The derived category of a commutative ring

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Vopěnka's principle

Vopěnka's principle (Categorical formulation)

Given any family of objects X_s of an accessible category indexed by the class of all ordinals, there is a morphism $X_s \rightarrow X_t$ for some ordinal $s < t$.

If Vopěnka's principle holds, then every full subcategory of a locally presentable category closed under λ -filtered colimits for some regular cardinal λ is accessible.

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Main results

Theorem 1 (CGR)

Let \mathcal{K} be a locally presentable category with a stable model category structure. If Vopěnka's principle holds, then every full subcategory \mathcal{L} of $\mathrm{Ho}(\mathcal{K})$ closed under fibres and products is reflective. If \mathcal{L} is semicolocalizing, then the reflection is semiexact. If \mathcal{L} is colocalizing, then the reflection is exact.

Corollary

Let \mathcal{K} be a locally presentable stable model category. If Vopěnka's principle holds, then every closed semilocalizing subcategory of $\mathrm{Ho}(\mathcal{K})$ is coreflective.

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Main results

Corollary

Let \mathcal{T} be a triangulated category with combinatorial models. Assuming Vopěnka's principle, every semicolocalizing subcategory of \mathcal{T} is equal to E^\perp for some object E (i.e., a nullity class) and every colocalizing subcategory is equal to E^\perp for some E (i.e., a stable nullity class).

Summary

Assuming **Vopěnka's principle**:

- Every semilocalizing subcategory of a triangulated category with combinatorial models is part of a t -structure and the same happens with every semicolocalizing subcategory.
- In every triangulated category with combinatorial models there is a bijective correspondence between the class of (semi)localizing subcategories and (semi)colocalizing subcategories.
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