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From subfactors to categories and topology I: Frobenius algebras in and Morita equivalence of tensor categories

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Abstract

We consider certain categorical structures that are implicit in subfactor theory. Making the connection between subfactor theory (at finite index) and category theory explicit sheds light on both subjects. Furthermore, it allows various generalizations of these structures, e.g. to arbitrary ground fields, and the proof of new results about topological invariants in three dimensions.

The central notion is that of a Frobenius algebra in a tensor category \mathcal{A} , which reduces to the classical notion if $\mathcal{A} = \mathbb{F}\text{-Vect}$, where \mathbb{F} is a field. An object $X \in \mathcal{A}$ with two-sided dual \bar{X} gives rise to a Frobenius algebra in \mathcal{A} , and under weak additional conditions we prove a converse: There exists a bicategory \mathcal{E} with $\text{Obj } \mathcal{E} = \{\mathcal{U}, \mathcal{B}\}$ such that $\text{End}_{\mathcal{E}}(\mathcal{U}) \cong^{\otimes} \mathcal{A}$ and such that there are $J, \bar{J}: \mathcal{B} \rightleftharpoons \mathcal{U}$ producing the given Frobenius algebra. Many properties (additivity, sphericity, semisimplicity,...) of \mathcal{A} carry over to the bicategory \mathcal{E} .

We define weak monoidal Morita equivalence of tensor categories, denoted $\mathcal{A} \approx \mathcal{B}$, and establish a correspondence between Frobenius algebras in \mathcal{A} and tensor categories $\mathcal{B} \approx \mathcal{A}$. While considerably weaker than equivalence of tensor categories, weak monoidal Morita equivalence $\mathcal{A} \approx \mathcal{B}$ has remarkable consequences: \mathcal{A} and \mathcal{B} have equivalent (as braided tensor categories) quantum doubles ('centers') and (if \mathcal{A}, \mathcal{B} are semisimple spherical or $*$ -categories) have equal dimensions and give rise the same state sum invariant of closed oriented 3-manifolds as recently defined by Barrett and Westbury. An instructive example is provided by finite-dimensional semisimple and cosemisimple Hopf algebras, for which we prove $H - \text{mod} \approx \hat{H} - \text{mod}$.

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The present formalism permits a fairly complete analysis of the center of a semisimple spherical category, which is the subject of the companion paper (J. Pure Appl. Algebra 180 (2003) 159–219).

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1. Introduction

Since tensor categories (or monoidal categories), in particular symmetric ones, have traditionally been part and parcel of the representation theory of groups it is hardly surprising that they continue to keep this central position in the representation theory of quantum groups, loop groups and of conformal field theories. See, e.g. [10,30]. The main new ingredient in these applications is the replacement of the symmetry by a braiding [28] which suggests connections with topology. Braided tensor categories have in fact served as an input in new constructions of invariants of links and 3-manifolds and of topological quantum field theories [68,31]. (Recently it turned out [4,20] that a braiding is not needed for the construction of the triangulation or ‘state sum’ invariant of 3-manifolds.)

A particular rôle in this context has been played by subfactor theory, see e.g. [24,52,26,16], which has led to the discovery of Jones’ polynomial invariant for knots [25]. Since the Jones polynomial was quickly reformulated in more elementary terms, and due to the technical difficulty of subfactor theory, the latter seems to have lost some of the attention of the wider public. This is deplorable, since operator algebraists continue to generate ideas whose pertinence extends beyond subfactor theory, e.g. in [53,15,73,23]. The series of papers of which this is the first aims at extracting the remarkable categorical structure which is inherent in subfactor theory, generalizing it and putting it to use for the proof of new results in categorical algebra and low-dimensional topology. As will be evident to experts, the series owes much to the important contributions of A. Ocneanu who, however, never advocated a categorical point of view. We emphasize that our works will not assume any familiarity with subfactor theory—they are in fact also meant to convey the author’s understanding of what (the more algebraic side of) the theory of finite index subfactors is about.

The present paper is devoted to the proof of several results relating two-sided adjoint 1-morphisms in 2-categories and Frobenius algebras in tensor categories. Here we are in particular inspired by A. Ocneanu’s notion of ‘paragroups’ and R. Longo’s description of type III subfactors in terms of ‘Q-systems’. In order to make the series accessible to readers with different backgrounds we motivate the constructions of the present paper by considerations departing from classical Frobenius theory, from category theory and from subfactor theory. While we will ultimately be interested in semisimple spherical categories, a sizable part of our considerations holds in considerably greater generality.

1.1. Classical Frobenius algebras

One of the many equivalent criteria for a finite-dimensional algebra A over a field \mathbb{F} to be a Frobenius algebra is the existence of a linear form $\phi: A \rightarrow \mathbb{F}$ for which

the bilinear form $b(a, b) = \phi(ab)$ is non-degenerate. (For a nice exposition of the present state of Frobenius theory we refer to [29].) Recent results of Quinn and Abrams [58, 1, 2] provide the following alternative characterization: A Frobenius algebra is a quintuple $(A, m, \eta, \Delta, \varepsilon)$, where (A, m, η) and (A, Δ, ε) are a finite-dimensional algebra and coalgebra, respectively, over \mathbb{F} , subject to the condition

$$m \otimes id_A \circ id_A \otimes \Delta = \Delta \circ m = id_A \otimes m \circ \Delta \otimes id_A.$$

In Section 6 we will say a bit more about the relation between these two definitions. Since an \mathbb{F} -algebra (coalgebra) is just a monoid (comonoid) in the tensor category $\mathbb{F}\text{-Vect}$, it is clear that the second definition of a Frobenius algebra makes sense in any tensor category. A natural problem therefore is to obtain examples of Frobenius algebras in categories other than $\mathbb{F}\text{-Vect}$ and to understand their significance. The aim of the next two subsections will be to show how Frobenius algebras arise in category theory and subfactor theory.

1.2. Adjoint functors and adjoint morphisms

We assume the reader to be conversant with the basic definitions of categories, functors and natural transformations, [44] being our standard reference. (In the next section we will recall some of the relevant definitions.) As is well known, the concept of adjoint functors is one of the most important ones not only in category theory itself but also in its applications to homological algebra and algebraic geometry. Before we turn to the generalizations which we will need to consider we recall the definition. Given categories \mathcal{C}, \mathcal{D} and functors $F: \mathcal{C} \rightarrow \mathcal{D}$, $G: \mathcal{D} \rightarrow \mathcal{C}$, F is a left adjoint of G , equivalently, G is a right adjoint of F , iff there are bijections

$$\phi_{X,Y}: \text{Hom}_{\mathcal{D}}(FX, Y) \cong \text{Hom}_{\mathcal{C}}(X, GY), \quad X \in \text{Obj } \mathcal{C}, Y \in \text{Obj } \mathcal{D},$$

that are natural w.r.t. X and Y . This is denoted $F \dashv G$. More convenient for the purpose of generalization is the equivalent characterization, according to which F is a left adjoint of G iff there are natural transformations $r: id_{\mathcal{C}} \rightarrow GF$ and $s: FG \rightarrow id_{\mathcal{D}}$ satisfying

$$id_G \otimes s \circ r \otimes id_G = id_G, \quad s \otimes id_F \circ id_F \otimes r = id_F.$$

Given an adjunction $F \dashv G$, the composite functor $T = GF$ is an object in the (strict tensor) category $\mathcal{X} = \text{Fun}(\mathcal{C})$ of endofunctors of \mathcal{C} . (When seen as an object of \mathcal{X} , we denote the identity functor of \mathcal{C} by $\mathbf{1}$.) By the above alternative characterization of adjoint pairs there are $\eta \equiv r \in \text{Hom}_{\mathcal{X}}(\mathbf{1}, T)$ and $m \in \text{Hom}_{\mathcal{X}}(T^2, T)$ defined by $m \equiv id_G \otimes s \otimes id_F$. These morphisms satisfy

$$m \circ m \otimes id_T = m \otimes id_T \otimes m,$$

$$m \circ \eta \otimes id_T = m \otimes id_T \otimes \eta = id_T;$$

thus (T, m, η) is a monoid in $\mathcal{X} = \text{Fun}(\mathcal{C})$, equivalently, a monad in \mathcal{C} . Similarly, one finds that $(U, \Delta, \varepsilon) \equiv (FG, id_F \otimes r \otimes id_G, s)$ is a comonoid in $\text{Fun}(\mathcal{D})$ or a comonad in \mathcal{D} .

Now, given a monad in a category \mathcal{C} one may ask whether it arises from an adjunction as above. This is always the case, there being two canonical solutions given by Kleisli and Eilenberg/Moore, respectively. See [44, Chapter VI] for the definitions and proofs. In fact, considering an appropriate category of all adjunctions yielding the given monad, the above particular solutions are initial and final objects, respectively.

So far, we have been considering the particular 2-category \mathcal{CAT} of small categories. Thanks to the second definition of adjoints most of the above considerations generalize to an arbitrary 2-category \mathcal{F} (or even a bicategory). See [44, Chapter 12] and [31,22] for introductions to 2- and bi-categories. Given objects (0-cells) \mathcal{C}, \mathcal{D} and 1-morphisms (1-cells) $F: \mathcal{C} \rightarrow \mathcal{D}, G: \mathcal{D} \rightarrow \mathcal{C}$ we say G is a right adjoint of F iff there are 2-morphisms (2-cells) r, s with the above properties. An important special case pertains if the 2-category \mathcal{F} has only one object, say \mathcal{C} . By the usual ‘dimension shift’ argument it then is the same as a (strict) tensor category, and the adjoint 1-morphisms become dual objects in the usual sense.

Now, a monad (comonad) in a general 2-category \mathcal{F} is most naturally defined [61] as an object \mathcal{C} in \mathcal{F} (the basis) together with a monoid (resp. comonoid) in the tensor category $END_{\mathcal{F}}(\mathcal{C})$. It is clear that an adjoint pair $F: \mathcal{C} \rightarrow \mathcal{D}, G: \mathcal{D} \rightarrow \mathcal{C}$ again gives rise to a monad (comonad) in \mathcal{F} with basis objects \mathcal{C} (\mathcal{D}). Again, the natural question arises whether every monad is produced by an adjoint pair of 1-morphisms. Without restrictive assumptions on \mathcal{F} this will in general not be true. (See [61] for a property of a 2-category which guarantees that every monad arises from an adjoint pair of 1-morphisms.) The aim of the present work is to explore a different aspect of this problem for which we need another preparatory discussion.

Assuming a 1-morphism in a 2-category \mathcal{F} (or an object in a tensor category) has both left and right adjoints, there is in general no reason why they should be isomorphic (i.e. related by an invertible 2-morphism). Yet there are tensor categories where every object has a two-sided dual, in particular rigid symmetric categories (=closed categories [44]), rigid braided ribbon categories (=tortile categories [27]), $*$ -categories [42], and most generally, pivotal categories. (Functors, i.e. 1-morphisms in \mathcal{CAT} , which have a two-sided adjoint are occasionally called ‘Frobenius functors’.) If a 1-morphism F happens to have a simultaneous left and right dual G , then not only GF gives rise to a monad and FG to a comonad, but also vice versa. But in fact, there is more structure. Let $p: id_{\mathcal{D}} \rightarrow FG$ and $q: GF \rightarrow id_{\mathcal{C}}$ be the 2-morphisms associated with the adjunction $G \dashv F$ and denote $T = GF$, $\varepsilon = q$, $\Delta = id_G \otimes p \otimes id_F$. Then $(T, m, \eta, \Delta, \varepsilon)$ is a monoid and a comonoid in $End(\mathcal{C})$, but in addition we have

$$id_T \otimes m \circ \Delta \otimes id_T = \Delta \circ m = m \otimes id_T \circ id_T \otimes \Delta.$$

The first half of this equation is proved diagrammatically by

and the other half similarly. Thus, a 1-morphism $F: \mathfrak{U} \rightarrow \mathfrak{B}$ with a two-sided dual gives rise to Frobenius algebras in the tensor categories $\text{End}(\mathfrak{U})$ and $\text{End}(\mathfrak{B})$. Again, one may ask whether there is a converse, i.e. if every Frobenius algebra in a tensor category arises as above. Already in the category $\mathcal{C} = \mathbb{F}\text{-Vect}$ one finds Frobenius algebras $Q = (Q, v, v', w, w')$ that are not of the form $Q = X\bar{X}$ for $X \in \mathcal{C}$. As one of our main results (Theorem 3.11), we will see that under certain conditions on a Frobenius algebra Q in a tensor category \mathcal{A} there is a solution if one embeds \mathcal{A} as a corner into a bicategory \mathcal{E} . I.e., there are a bicategory \mathcal{E} , objects $\mathfrak{U}, \mathfrak{B} \in \text{Obj } \mathcal{E}$ and a 1-morphism $J: \mathfrak{B} \rightarrow \mathfrak{U}$ with two-sided dual \bar{J} such that $\mathcal{A} \simeq \text{END}_{\mathcal{E}}(\mathfrak{U})$ and Q arises via $Q = J\bar{J}$, etc. If \mathcal{A} is (pre)additive, abelian, \mathbb{F} -linear, (semisimple) spherical or a $*$ -category then \mathcal{E} will have the same properties. Under certain conditions, the bicategory \mathcal{E} is equivalent to the bicategory \mathcal{F} containing the 1-morphism $J: \mathfrak{B} \rightarrow \mathfrak{U}$ which gave rise to the Frobenius algebra.

1.3. Subfactors

For some basic definitions concerning subfactors we refer to Section 6.4. For the purposes of this introduction it is sufficient to know that a factor is a complex unital $*$ -algebra with center $\mathbb{C}1$, usually infinite dimensional. (Every finite-dimensional factor is isomorphic to a matrix algebra $M_n(\mathbb{C})$ for some $n \in \mathbb{N}$.) A factor ‘has separable predual’ if it admits a faithful continuous representation on a separable Hilbert space. We restrict ourselves to such factors, which we (abusively) call separable. In his seminal work [24], Vaughan Jones introduced a notion of index $[M: N] \in [1, \infty]$, defined for every inclusion $N \subset M$ of type II_1 factors. The index was soon generalized to factors of arbitrary type by Kosaki [35]. The index shares some basic properties with the one for groups: $[M: N] = 1$ iff $M = N$, $[P: M] \cdot [M: N] = [P: N]$ whenever $N \subset M \subset P$, etc. Yet the index is not necessarily an integer, in fact every $\lambda \in [4, \infty]$ can occur, whereas in the interval $[1, 4]$ only the countable set $\{4 \cos^2(\pi/n), n=3, 4, \dots\}$ is realized. Given factors P, Q (in fact, for general von Neumann algebras), there is a notion of $P - Q$ bimodule, and for bimodules ${}_P\mathcal{H}_Q, {}_Q\mathcal{K}_R$ there is a relative tensor product ${}_P\mathcal{H}_Q \otimes {}_Q\mathcal{K}_R$. This gives rise to a bicategory \mathcal{BIM} whose objects are factors, the 1-morphisms being bimodules and the 2-morphisms being the bimodule homomorphisms. Another bicategory \mathcal{MOR} arising from factors has as objects all factors, as 1-morphisms the continuous unital $*$ -algebra homomorphisms $\rho: P \rightarrow Q$ and as 2-morphisms the intertwiners. Thus if $\rho, \sigma: P \rightarrow Q$ then

$$\text{Hom}_{\mathcal{MOR}}(\rho, \sigma) = \{x \in Q \mid x\rho(y) = \sigma(y)x \quad \forall y \in P\}.$$

Whereas the definition of the tensor product of bimodules is technically involved (in particular in the non-type II_1 case, cf. [11, Appendix V.B]), the composition of 1-morphisms $\rho: P \rightarrow Q, \sigma: Q \rightarrow R$ in \mathcal{MOR} is just their composition $\sigma \circ \rho$ as maps and the unit 1-morphisms are the identity maps. Note that the composition of 1-morphisms in \mathcal{MOR} is strictly associative, thus \mathcal{MOR} is a 2-category. Every subfactor $N \subset M$ gives rise to a distinguished 1-morphism $\iota_{N,M}: N \rightarrow M$, the embedding map $N \hookrightarrow M$.

Both \mathcal{BIM} and \mathcal{MOR} are $*$ -bicategories, i.e. come with antilinear involutions on the 2-morphisms which reverse the direction. (For $s \in \text{Hom}_{\mathcal{BIM}}({}_P\mathcal{H}_Q, {}_P\mathcal{H}_Q) \subset \mathcal{B}(\mathcal{H}, \mathcal{H})$ and $t \in \text{Hom}_{\mathcal{MOR}}(\rho, \sigma) \subset Q$, where $\rho, \sigma: P \rightarrow Q$, s^* is given by the adjoint in $\mathcal{B}(\mathcal{H}, \mathcal{H})$ and in Q , respectively.) If one restricts oneself to separable type III factors the corresponding full sub-bicategories are equivalent in the sense of [7]: $\mathcal{BIM}_{\text{III}}^{\text{sep}} \simeq \mathcal{MOR}_{\text{III}}^{\text{sep}}$, cf. [11,39]. In the following discussion we will focus on the 2-category $\mathcal{MOR}_{\text{III}}^{\text{sep}}$.

The relevance of the concept of adjoints in 2-categories to subfactor theory is evident in view of the following result which is due to Longo [39], though originally not formulated in this way.

Theorem 1.1. *Let $N \subset M$ be an inclusion of separable type III factors. The embedding morphism $\iota: N \hookrightarrow M$ has a two-sided adjoint $\bar{\iota}: M \hookrightarrow N$ iff $[M : N] < \infty$. The dimensions of $\iota, \bar{\iota}$ (in the sense of [42]) are related to the index by $d(\iota) = d(\bar{\iota}) = [M : N]^{1/2}$.*

Since separable type III factors are simple, every morphism $P \rightarrow Q$ is in fact an embedding. We thus have the following more symmetric formulation.

Corollary 1.2. *Let $\rho: P \rightarrow Q$ be a morphism of separable type III factors. Then ρ has a two-sided adjoint $\bar{\rho}: Q \rightarrow P$ iff $[Q : \rho(P)] < \infty$. Then $d(\rho) = d(\bar{\rho}) = [P : \bar{\rho}(Q)]^{1/2} = [Q : \rho(P)]^{1/2}$.*

If these equivalent conditions are satisfied, an important object of study is the full sub-bicategory of $\mathcal{MOR}_{\text{III}}^{\text{sep}}$ generated by ρ and $\bar{\rho}$, which consists of all morphisms between P and Q which are obtained by composition of $\rho, \bar{\rho}$, the retracts and finite direct sums of such. (The type II_1 analog is Ocneanu’s paragroup [52] associated with $N \subset M$.)

In this situation, $\gamma = \bar{\rho} \circ \rho: P \rightarrow P$ and $\sigma = \rho \circ \bar{\rho}: Q \rightarrow Q$ are endomorphisms of P and Q , respectively, the so-called canonical endomorphisms. By the considerations of the preceding section we know that there is a Frobenius algebra (γ, v, v', w, w') in $\text{End}(P) := \text{END}_{\mathcal{MOR}}(P)$ and similarly for σ . Since we are in a $*$ -categorical setting we have $v' = v^*, w' = w^*$. In [40] such triples (γ, v, w) were called Q-systems and it was shown that for every Q-system where $\gamma \in \text{End}(M)$ there exists a subfactor $N \subset M$ such that (γ, v, w) arises as above. These results were clearly motivated by the notion of conjugates (duals) in tensor categories, but the constructions of Kleisli and Eilenberg/Moore do not seem to have played a rôle.

Longo’s results can be rephrased by saying that given a separable type III factor M there is a bijection between subfactors $N \subset M$ with $[M : N] < \infty$ and Q-systems (\simeq Frobenius algebras) in the tensor category $\text{End}(M)$. Given such a Frobenius algebra Q , we can on the one hand construct a subfactor N , an adjoint pair $\rho: M \rightarrow N$, $\bar{\rho}: N \rightarrow M$ and the 2-category generated by them. On the other hand we can regard $\text{End}(M)$ as an abstract $*$ -category and apply to Q the construction announced in the preceding subsection. It will turn out that these two procedures give rise to equivalent 2-categories.

1.4. Organization of the paper

In the following section we will discuss some preliminaries on (tensor) categories, mostly concerning the notions of duality and dimension. The emphasis in this paper is on spherical tensor categories, but we also consider $*$ -categories. In Section 3, we establish the connection between two-sided duals in bicategories and Frobenius algebras in tensor categories. This is used in Section 4 to define the notion of weak monoidal Morita equivalence. In Section 5 we consider categories that are linear over some field, in particular spherical and $*$ -categories. We show that the bicategory \mathcal{E} constructed from a Frobenius algebra in a $*$ -category or a (semisimple) spherical category is a $*$ -bicategory or (semisimple) spherical bicategory, and we prove the equality of certain dimensions. Section 6 is devoted to several examples. We begin with classical Frobenius algebras, i.e. Frobenius algebras in the category $\mathbb{F}\text{-Vect}$, and specialize to finite-dimensional Hopf algebras. Our main result is the weak monoidal Morita equivalence $H - \text{mod} \approx \bar{H} - \text{mod}$ whenever H is semisimple and cosemisimple. We also discuss in more detail the examples provided by subfactor theory. An application to the subject of quantum invariants of 3-manifolds is given in Section 7. We outline a proof of the fact that spherical categories that are weakly monoidally Morita equivalent define the same state sum invariant of 3-manifolds. In the last section, we briefly relate our results to previous works and conclude with some announcements of further results and open problems.

2. Categorical preliminaries

2.1. Some basic notions and notations

We assume that the standard definitions of tensor categories and bicategories are known, cf. [44]. For definiteness all categories in this paper are supposed small. (In the early stages essential smallness would suffice, whereas later on we will even require the number of isoclasses to be finite.) We use ‘tensor category’ and ‘monoidal categories’ interchangeably. Tensor categories will usually be assumed strict. (A tensor category is strict if the tensor product satisfies associativity $X \otimes (Y \otimes Z) = (X \otimes Y) \otimes Z$ on the nose and the unit object $\mathbf{1}$ satisfies $X \otimes \mathbf{1} = \mathbf{1} \otimes X = X \ \forall X$. Similarly, in a strict bicategory (= 2-category) the composition of 1-morphisms is associative.) Since every tensor category is equivalent to a strict one [44,28] and every bicategory to a 2-category this does not restrict the generality of our results.

Throughout, (2-)categories will be denoted by calligraphic letters $\mathcal{A}, \mathcal{B}, \mathcal{C}, \dots$ and objects and morphisms in 1-categories by capital and lowercase Latin letters, respectively. In 2-categories objects, 1- and 2-morphisms are denoted by Gothic ($\mathfrak{U}, \mathfrak{B}, \dots$), capital and lowercase Latin letters, respectively. Unit objects and unit morphisms in tensor categories are denoted by $\mathbf{1}$ and $id_X \in \text{Hom}(X, X)$, respectively. Similarly, the unit 1- and 2-morphisms in 2-categories are $\mathbf{1}_{\mathfrak{U}} \in \text{Hom}(\mathfrak{U}, \mathfrak{U})$ and $id_X \in \text{Hom}(X, X)$. If $\mathfrak{U}, \mathfrak{B}$ are objects in a bicategory \mathcal{E} then $\text{Hom}_{\mathcal{E}}(\mathfrak{U}, \mathfrak{B})$ and $\text{HOM}_{\mathcal{E}}(\mathfrak{U}, \mathfrak{B})$ denote the corresponding 1-morphisms as a set and as a (1-)category (whose morphisms are the

2-morphisms in \mathcal{E}), respectively. As is well known, $END_{\mathcal{E}}(\mathcal{U}) \equiv HOM_{\mathcal{E}}(\mathcal{U}, \mathcal{U})$ is a tensor category. (The composition \circ of $\mathcal{U} - \mathcal{U}$ -morphisms in \mathcal{E} is the tensor product of objects in \mathcal{A} and the compositions \circ, \otimes of 2-morphisms in \mathcal{E} are the compositions \circ, \otimes of morphisms in \mathcal{A} .) Since the use of the composition symbol \circ for $\mathcal{U} - \mathcal{U}$ -morphisms in \mathcal{E} as opposed to the tensor product \otimes of objects in \mathcal{A} might lead to confusion (and there is only one composition for these items anyway) we mostly omit the composition symbols \otimes for objects in \mathcal{A} and \circ for 1-morphisms in \mathcal{E} altogether.

We write $Y \prec X$ if Y is a retract of X , i.e. there are morphisms $e: Y \rightarrow X, f: X \rightarrow Y$ such that $f \circ e = id_Y$. We also say Y is a subobject of X , slightly incompatibly with common usage. In this situation $p = e \circ f \in End(X)$ is idempotent: $p^2 = p$. Categories in which every idempotent arises in this way are called ‘pseudo-abelian’ (Karoubi), ‘Karoubian’ (SGA) or ‘Cauchy-complete’ [9]. Others say that ‘idempotents split in \mathcal{A} ’ or ‘ \mathcal{A} has subobjects’ [12]. We consider none of these expressions particularly satisfactory. We will stick to the last alternative since it goes best with ‘ \mathcal{A} has direct sums’. Every category \mathcal{A} can be embedded as a full subcategory into one which has subobjects. The latter can be defined as solution to a universal problem, cf. [9], which implies uniqueness up to equivalence. There is, however, a well-known canonical solution, which we call $\tilde{\mathcal{A}}^p$, cf. e.g. [19]. Its objects are pairs (X, p) where $X \in Obj \mathcal{A}$ and $p = p^2 \in End(X)$. The morphisms are given by

$$Hom_{\tilde{\mathcal{A}}^p}((X, p), (Y, q)) = \{s: X \rightarrow Y \mid s = q \circ s \circ p\}.$$

If \mathcal{A} is a tensor category then so is $\tilde{\mathcal{A}}^p$ with $(X, p) \otimes (Y, q) = (X \otimes Y, p \otimes q)$.

A preadditive category (or Ab-category) is a category where all hom-sets are abelian groups and the composition is additive w.r.t. both arguments. A preadditive category \mathcal{A} is said to have direct sums (or biproducts) if for every $X_1, X_2 \in \mathcal{A}$ there are Y and morphisms $v_i \in Hom(X_i, Y), v'_i \in Hom(Y, X_i)$ such that $v_1 \circ v'_1 + v_2 \circ v'_2 = id_Y$ and $v'_i \circ v_j = \delta_{i,j} id_{X_i}$. We then write $Y \cong X_1 \oplus X_2$. Note that every $Y' \cong Y$ is a direct sum of X_1, X_2 , too. A preadditive category can always be embedded as a full subcategory into one with (finite) direct sums. There is a canonical such category, cf. e.g. [19], which we call $\tilde{\mathcal{A}}^{\oplus}$. Again, this construction is compatible with a monoidal structure on \mathcal{A} . Constructions completely analogous to $\tilde{\mathcal{A}}^p, \tilde{\mathcal{A}}^{\oplus}$ exist for a bicategory \mathcal{E} , cf. [42, Appendix]. (Thus in $\tilde{\mathcal{E}}^p$ all idempotent 2-morphisms split and $\tilde{\mathcal{E}}^{\oplus}$ has direct sums for parallel 1-morphisms.) A pre-additive category is additive if it has direct sums and a zero object. If \mathcal{A} is preadditive then $\tilde{\mathcal{A}}^p$ is additive: any $(X, 0)$ is a zero object. Furthermore, $\tilde{\mathcal{A}}^{p \oplus} \simeq \tilde{\mathcal{A}}^{\oplus p}$.

Given a commutative ring k , a (monoidal) category is k -linear if all hom-sets are finitely generated k -modules and the composition \circ (and tensor product \otimes) of morphisms are bilinear. Mostly, k will be a field \mathbb{F} . An object in a k -linear category is called simple if $End X = k id_X$. (This property is often called absolute simplicity or irreducibility. We drop the attribute ‘absolute’, see the remarks below.) A k -linear category \mathcal{C} is called semisimple if it has direct sums and subobjects and there are simple objects X_i labeled by a set I which are mutually disjoint ($i \neq j \Rightarrow Hom(X_i, X_j) = \{0\}$) such that the obvious map

$$\bigoplus_{i \in I} Hom(Y, X_i) \otimes_{\mathbb{F}} Hom(X_i, Z) \rightarrow Hom(Y, Z)$$

is an isomorphism. Then every object X is a finite direct sum of objects in $\{X_i, i \in I\}$ and is determined up to isomorphism by the function

$$I \rightarrow \mathbb{Z}_+, \quad i \mapsto N_i^X = \dim \operatorname{Hom}(X, X_i) = \dim \operatorname{Hom}(X_i, X).$$

If \mathcal{C} is monoidal we also require **1** to be simple.

In this paper we will not use the language of abelian categories, since we will not need the notions of kernels and cokernels. (However, when applied to an abelian category the constructions given below give rise to abelian categories. See Remark 3.18.) We briefly relate our definition of semisimplicity to more conventional terminology. An abelian category is semisimple if it satisfies the following equivalent properties: (i) all short exact sequences split, (ii) every monic is a retraction, (iii) every epi is a section. An object X in an abelian category is indecomposable if it is not a direct sum of two non-zero objects, equivalently, if $\operatorname{End} X$ contains no idempotents besides 0 and id_X . If \mathcal{C} is semisimple and $X \in \mathcal{C}$ is indecomposable then $\operatorname{End} X$ is a skew field. Thus, if \mathbb{F} is algebraically closed and \mathcal{C} is \mathbb{F} -linear then $X \in \mathcal{C}$ is indecomposable iff it is (absolutely) simple. Since an abelian category has direct sums and subobjects, \mathcal{C} is semisimple in our sense. Conversely, if \mathcal{C} is semisimple in our sense and has a zero object then it is semisimple abelian, whether the field \mathbb{F} is algebraically closed or not.

A subcategory $\mathcal{S} \subset \mathcal{C}$ is full iff $\operatorname{Hom}_{\mathcal{S}}(X, Y) = \operatorname{Hom}_{\mathcal{C}}(X, Y) \quad \forall X, Y \in \mathcal{S}$, thus it is determined by $\operatorname{Obj} \mathcal{S}$. A subcategory is replete iff $X \in \operatorname{Obj} \mathcal{S}$ implies $Y \in \operatorname{Obj} \mathcal{S}$ for all $Y \cong X$. (In the literature replete full subcategories are also called strictly full.) Most subcategories we consider will be replete full. Isomorphism of categories is denoted by \cong and equivalence by \simeq . For preadditive (k -linear) categories the functors establishing the equivalence/isomorphism are required to be additive (k -linear) and for tensor categories they must be monoidal. In principle one should use qualified symbols like $\cong_+, \simeq_k, \overset{\otimes}{\cong}, \overset{\otimes}{\simeq}_k$, etc., where $+, k$ and \otimes stand for additivity, k -linearity and monoidality, respectively, of the equivalence/isomorphism. We will drop the $+, k$, hoping that they are obvious from the context, but will write the \otimes .

The following definition is somewhat less standard. (Recall that we require the categories to be small.)

Definition 2.1. Two categories \mathcal{A}, \mathcal{B} are Morita equivalent, denoted $\mathcal{A} \cong \mathcal{B}$, iff the categories $\operatorname{Fun}(\mathcal{A}^{\operatorname{op}}, \operatorname{Sets}), \operatorname{Fun}(\mathcal{B}^{\operatorname{op}}, \operatorname{Sets})$ are equivalent. Preadditive categories are Morita equivalent iff the categories $\operatorname{Fun}_+(\mathcal{A}^{\operatorname{op}}, \operatorname{Ab}), \operatorname{Fun}_+(\mathcal{B}^{\operatorname{op}}, \operatorname{Ab})$ of additive functors are equivalent as preadditive categories. k -linear categories are Morita equivalent iff the categories $\operatorname{Fun}_k(\mathcal{A}^{\operatorname{op}}, k\text{-mod}), \operatorname{Fun}_k(\mathcal{B}^{\operatorname{op}}, k\text{-mod})$ of k -linear functors are equivalent as k -linear categories.

Proposition 2.2. Let \mathcal{A}, \mathcal{B} be categories. Then $\mathcal{A} \cong \mathcal{B}$ iff $\tilde{\mathcal{A}}^P \simeq \tilde{\mathcal{B}}^P$. If \mathcal{A}, \mathcal{B} are preadditive (k -linear) categories then $\mathcal{A} \cong \mathcal{B}$ iff $\tilde{\mathcal{A}}^{P\oplus} \simeq \tilde{\mathcal{B}}^{P\oplus}$ as preadditive (k -linear) categories.

Proof. For the first claim see [9, Section 6.5], for the others [19, Chapter 2]. \square

This result motivates the following definition of (strong) Morita equivalence for monoidal categories.

Definition 2.3. Two tensor categories \mathcal{A}, \mathcal{B} are strongly monoidally Morita equivalent ($\mathcal{A} \stackrel{\otimes}{\cong} \mathcal{B}$) iff $\tilde{\mathcal{A}}^P \stackrel{\otimes}{\cong} \tilde{\mathcal{B}}^P$. If \mathcal{A}, \mathcal{B} are preadditive (k -linear) tensor categories then they are strongly Morita equivalent iff $\tilde{\mathcal{A}}^{P \oplus} \stackrel{\otimes}{\cong} \tilde{\mathcal{B}}^{P \oplus}$ as preadditive (k -linear) tensor categories.

While useful for certain purposes, this definition is unsatisfactory in that we cannot offer an equivalent definition in terms of module-categories. In Section 4, we will define the notion of weak Morita equivalence $\mathcal{A} \approx \mathcal{B}$ of (preadditive, k -linear) tensor categories. It is genuine to tensor categories and satisfies $\mathcal{A} \stackrel{\otimes}{\cong} \mathcal{B} \Rightarrow \mathcal{A} \approx \mathcal{B}$. We speculate that $\mathcal{A} \approx \mathcal{B}$ iff suitably defined ‘representation categories’ of \mathcal{A}, \mathcal{B} are equivalent, but we will leave this problem for future investigations.

2.2. Duality in tensor categories and 2-categories

As explained in the Introduction, the notion of adjoint functors generalizes from \mathcal{CAT} to an arbitrary 2-category \mathcal{C} . Specializing to tensor categories, i.e. one-object 2-categories we obtain the following well-known notions. We recall that a tensor category \mathcal{C} is said to have left (right) duals if for every $X \in \mathcal{C}$ there is a *X (X^*) together with morphisms $e_X : \mathbf{1} \rightarrow X \otimes {}^*X$, $d_X : {}^*X \otimes X \rightarrow \mathbf{1}$ ($\varepsilon_X : \mathbf{1} \rightarrow X^* \otimes X$, $\eta_X : X \otimes X^* \rightarrow \mathbf{1}$) satisfying the usual duality equations

$$id_X \otimes d_X \circ e_X \otimes id_X = id_X, \quad d_X \otimes id_{\bar{X}} \circ id_{\bar{X}} \otimes e_X = id_{\bar{X}},$$

etc. Categories having left and right duals for every object are called autonomous. Since duals, if they exist, are unique up to isomorphism the conditions ${}^*X \cong X^*$ and $X^{**} \cong X$, which are easily seen to be equivalent, do not involve any choices. If $X^* \cong {}^*X$ we speak of a two-sided dual of X . We will exclusively consider categories with two-sided duals, for which we use the symmetric notation \bar{X} . Assume now that \mathcal{C} is linear over the commutative ring $k \equiv \text{End}(\mathbf{1})$. Then we have $\eta_X \circ e_X, d_X \circ \varepsilon_X \in k$ and we would like to consider them as dimensions of X or \bar{X} . Yet, if λ is a unit in k then replacing e_X and d_X by $\lambda e_X, \lambda^{-1} d_X$, respectively, the triangular equations are still verified while $\eta_X \circ e_X$ and $d_X \circ \varepsilon_X$ change. (The product is invariant, though.) We thus need a way to eliminate this indeterminacy. There are three known solutions to this problem. If \mathcal{C} has a braiding $c_{-, -}$ and a twist $\theta(-)$ we can determine the right duality in terms of the left duality by $\varepsilon_X = (id_{\bar{X}} \otimes \theta_X) \circ c_{X, \bar{X}} \circ e_X$ and $\eta_X = d_X \circ c_{X, \bar{X}} \circ (\theta_X \otimes id_{\bar{X}})$, allowing to unambiguously define $d(X) = \eta_X \circ e_X$. Since in this paper, we do not require the existence of braidings this approach is of no use to us. A fairly satisfactory way to eliminate the indeterminacy exists if \mathcal{C} has a $*$ -operation, see Section 2.4. The third solution, cf. the next subsection, is provided by the notion of spherical categories which was introduced by Barrett and Westbury [5,4], elaborating on earlier work on pivotal or sovereign categories. (For categories with braiding the latter approach is related to the first one, in that there is a one-to-one correspondence between twists and spherical structures, cf. [77].) We believe that this is the most general setting within which

results like Proposition 5.17 and those of [47] obtain without a fundamental change of the methods. (*-categories can be turned into spherical categories, cf. [74], but this is not necessarily the most convenient thing to do.)

The following result shows that the *square* of the dimension of a simple object in a linear category is well defined even in the absence of further structure.

Proposition 2.4. *Let \mathcal{C} be an \mathbb{F} -linear tensor category with simple unit. If X is simple and has a two-sided dual then $d^2(X) = (\eta_X \circ e_X)(d_X \circ \varepsilon_X) \in \mathbb{F}$ is a well-defined quantity. If X, Y, XY are simple then $d^2(XY) = d^2(X)d^2(Y)$. Whenever \mathcal{C} has a spherical or *-structure $d^2(X)$ coincides with $d(X)^2$ as defined using the latter.*

Proof. Let X be simple, \bar{X} a two-sided dual and $e_X, d_X, \varepsilon_X, \eta_X$ the corresponding morphisms. By simplicity of X we have $\text{Hom}(\mathbf{1}, X\bar{X}) \cong \text{Hom}(X, X) \cong \mathbb{F}$, thus $\text{Hom}(\mathbf{1}, X\bar{X}) = \mathbb{F}e_X$, $\text{Hom}(\mathbf{1}, \bar{X}X) = \mathbb{F}\varepsilon_X$, etc. Therefore, any other solution of the triangular equations is given by

$$\tilde{e}_X = \alpha e_X, \quad \tilde{d}_X = \alpha^{-1} d_X, \quad \tilde{\varepsilon}_X = \beta \varepsilon_X, \quad \tilde{\eta}_X = \beta^{-1} \eta_X,$$

where $\alpha, \beta \in \mathbb{F}^*$. Thus $(\eta_X \circ e_X)(d_X \circ \varepsilon_X)$ does not depend on the choice of the morphisms. The independence $d^2(X)$ of the choice of \bar{X} and the multiplicativity of d^2 (only if XY is simple) are obvious. The final claim will follow from the fact that in spherical and *-categories $d(X)$ is defined as $\eta_X \circ d_X$ for a certain choice of d_X, η_X and the fact that $d(X) = d(\bar{X})$. \square

By this result the square of the dimension of a simple object is independent of the chosen spherical or *-structure and can in fact be defined without assuming the latter. Yet, consistently choosing signs of the dimensions and extending $d(X)$ to an additive and multiplicative function for all objects is a non-trivial cohomological problem to which there does not seem to be a simple solution. A *-structure, when available, provides the most natural way out, spherical structures being the second (but more general) choice.

The proposition implies that the following definition makes sense.

Definition 2.5. Let \mathcal{C} be a semisimple \mathbb{F} -linear tensor category with simple unit and two-sided duals. If \mathcal{C} has finitely many isomorphism classes of simple objects then we define

$$\dim \mathcal{C} = \sum_X d^2(X) \in \mathbb{F},$$

where the summation is over the isomorphism classes of simple objects and $d^2(X)$ is as in the proposition. If \mathcal{C} has infinitely many simple objects then we formally posit $\dim \mathcal{C} = \infty$.

We recommend the reader to skip the next two subsections until the structures introduced there will be needed in Section 5.

2.3. Spherical categories

In spherical categories the problem mentioned above is solved by picking a two-sided dual \bar{X} for every object X and by specifying morphisms $\mathbf{1} \rightarrow X \otimes \bar{X}$, $X \otimes \bar{X} \rightarrow \mathbf{1}$ as part of the given data. This may look unnatural, but in important cases such assignments are handy. (If, e.g. H is an involutive Hopf algebra then $\pi \mapsto \pi^t \circ S$ gives an involution on the representation category of H .) We recall that we consider only strict tensor categories, and by the coherence theorem of [4] we may assume also strict duality. Since later we will have occasion to construct spherical structures out of other data we give a redundancy-free definition of spherical categories, which then will be proven equivalent to that of [4,5].

Definition 2.6. A strict tensor category \mathcal{C} is a strict pivotal category if there is a map $Obj \mathcal{C} \rightarrow Obj \mathcal{C}, X \mapsto \bar{X}$ such that

$$\bar{\bar{X}} = X, \quad \overline{X \otimes Y} = \bar{Y} \otimes \bar{X}, \quad \bar{\mathbf{1}} = \mathbf{1} \quad (2.1)$$

and there are morphisms $\varepsilon(X): \mathbf{1} \rightarrow X \otimes \bar{X}$, $\bar{\varepsilon}(X): X \otimes \bar{X} \rightarrow \mathbf{1}, X \in \mathcal{C}$ satisfying the following conditions.

(1) The composites

$$X \equiv X \otimes \mathbf{1} \xrightarrow{id \otimes \varepsilon(\bar{X})} X \otimes \bar{X} \otimes X \xrightarrow{\bar{\varepsilon}(X) \otimes id} \mathbf{1} \otimes X \equiv X,$$

$$X \equiv \mathbf{1} \otimes X \xrightarrow{\varepsilon(X) \otimes id} X \otimes \bar{X} \otimes X \xrightarrow{id \otimes \bar{\varepsilon}(\bar{X})} X \otimes \mathbf{1} \equiv X$$

coincide with id_X .

(2) The diagrams

$$\begin{array}{ccc} \mathbf{1} & \xrightarrow{\varepsilon(X)} & X \otimes \bar{X} \\ \varepsilon(X \otimes Y) \downarrow & & \downarrow id \otimes \varepsilon(Y) \otimes id \\ X \otimes Y \otimes \overline{X \otimes Y} & \equiv \equiv \equiv & X \otimes Y \otimes \bar{Y} \otimes \bar{X} \end{array}$$

and

$$\begin{array}{ccc} \mathbf{1} & \xleftarrow{\bar{\varepsilon}(X)} & X \otimes \bar{X} \\ \bar{\varepsilon}(X \otimes Y) \uparrow & & \uparrow id \otimes \bar{\varepsilon}(Y) \otimes id \\ X \otimes Y \otimes \overline{X \otimes Y} & \equiv \equiv \equiv & X \otimes Y \otimes \bar{Y} \otimes \bar{X} \end{array}$$

commute for all X, Y .

(3) For every $s: X \rightarrow Y$ the composites

$$\bar{Y} \equiv \bar{Y} \otimes \mathbf{1} \xrightarrow{id \otimes \varepsilon(X)} \bar{Y} \otimes X \otimes \bar{X} \xrightarrow{id \otimes s \otimes id} \bar{Y} \otimes Y \otimes \bar{X} \xrightarrow{\bar{\varepsilon}(\bar{Y}) \otimes id} \mathbf{1} \otimes \bar{X} \equiv \bar{X},$$

$$\bar{Y} \equiv \mathbf{1} \otimes \bar{Y} \xrightarrow{\varepsilon(\bar{X}) \otimes id} \bar{X} \otimes X \otimes \bar{Y} \xrightarrow{id \otimes s \otimes id} \bar{X} \otimes Y \otimes \bar{Y} \xrightarrow{id \otimes \bar{\varepsilon}(Y)} \bar{X} \otimes \mathbf{1} \equiv \bar{X}$$

coincide in $Hom(\bar{Y}, \bar{X})$.

For every X and $s: X \rightarrow X$ we define morphisms in $End(\mathbf{1})$ by

$$tr_L(s) : \mathbf{1} \xrightarrow{\varepsilon(X)} X \otimes \bar{X} \xrightarrow{s \otimes id} X \otimes \bar{X} \xrightarrow{\bar{\varepsilon}(X)} \mathbf{1},$$

$$tr_R(s) : \mathbf{1} \xrightarrow{\varepsilon(\bar{X})} \bar{X} \otimes X \xrightarrow{id \otimes s} \bar{X} \otimes X \xrightarrow{\bar{\varepsilon}(\bar{X})} \mathbf{1}.$$

Now \mathcal{C} is called spherical iff $tr_L(s) = tr_R(s)$ for all s .

Remark 2.7. 1. If for $s \in Hom(X, Y)$ we now define $\bar{s} \in Hom(\bar{Y}, \bar{X})$ by the formulas in (3) we easily verify $\bar{\bar{s}} = s$ and $\overline{s \circ t} = \bar{t} \circ \bar{s}$. Thus the maps $X \mapsto \bar{X}, s \mapsto \bar{s}$ constitute an involutive contravariant endofunctor of \mathcal{C} . Condition (1) can now be expressed as $\overline{id_X} = id_{\bar{X}} \forall X$.

2. Using conditions (1) one verifies that $\bar{\varepsilon}(X) = \overline{\varepsilon(\bar{X})}$ and thus consistency of our notation.

3. The definition of the map $s \mapsto \bar{s}$ implies that the square

$$\begin{array}{ccc} \mathbf{1} & \xrightarrow{\varepsilon(X)} & X \otimes \bar{X} \\ \varepsilon(Y) \downarrow & & \downarrow s \otimes id \\ Y \otimes \bar{Y} & \xrightarrow{id \otimes \bar{s}} & Y \otimes \bar{X} \end{array} \quad (2.2)$$

commutes for every $s: X \rightarrow Y$. Thus, our axioms for strict pivotal categories imply those of [4,5], and the converse is obvious if we put $\bar{\varepsilon}(X) := \overline{\varepsilon(\bar{X})}$.

4. Even in those applications where the pivotal category \mathcal{C} under study is strict monoidal, like in type III subfactor theory, the duals are rarely strict. In general one has natural isomorphisms $\tau_X: X \rightarrow \bar{\bar{X}}, \gamma_{X,Y}: \bar{X} \otimes \bar{Y} \rightarrow \overline{Y \otimes X}$ and $v: \mathbf{1} \rightarrow \bar{\mathbf{1}}$ which must satisfy a number of compatibility conditions. See [5, Theorem 1.9], where it is proven that such a category is monoidally equivalent to a strict pivotal category as defined above. To be sure, in practise one does not really want to strictify the categories under consideration in order to work with them. Just as the well-known coherence results on (braided) tensor categories [44,28], [5, Theorem 1.9] can be rephrased as follows: All computations in a strict pivotal or spherical category remain valid in the non-strict case, provided one inserts the morphisms $\tau_X, \gamma_{X,Y}, v$ wherever needed. If this is possible in different ways no possible result depends on the choices one makes in this process.

In order to travel with slightly lighter luggage we may therefore calmly stick to the strict case.

5. A pivotal category is called non-degenerate if for all X, Y the pairing

$$\text{Hom}(X, Y) \times \text{Hom}(Y, X) \rightarrow \mathbb{F}, \quad s \times t \mapsto \langle s, t \rangle = \text{tr}_X(s \circ t) = \text{tr}_Y(t \circ s)$$

between $\text{Hom}(X, Y)$ and $\text{Hom}(Y, X)$ is non-degenerate. In [20, Lemma 3.1] it is proved that semisimple pivotal categories are non-degenerate.

If \mathcal{C} is \mathbb{F} -linear pivotal with simple unit $\mathbf{1}$ then we define dimensions by

$$d(X) \text{id}_{\mathbf{1}} = \text{tr}_L(X) = \bar{\varepsilon}(X) \circ \varepsilon(X).$$

If \mathcal{C} is spherical then

$$d(X) = d(\bar{X}) \quad \forall X.$$

We will show that in the semisimple case this is equivalent to sphericity. Let $X \cong \bigoplus_{i \in J} N_i X_i$, thus there are $u_i^\alpha: X_i \rightarrow X, u_i': X \rightarrow X_i^\alpha$ such that

$$\sum_i \sum_{\alpha=1}^{N_i} u_i^\alpha \circ u_i'^\alpha = \text{id}_X, \quad u_i'^\alpha \circ u_j^\beta = \delta_{i,j} \delta_{\alpha,\beta} \text{id}_{X_i}. \quad (2.3)$$

Using the conjugation functor—we define $v_i^\alpha := \overline{u_i'^\alpha}: X \rightarrow X_i, v_i'^\alpha := \overline{u_i^\alpha}: X_i \rightarrow X$, and clearly

$$\sum_i \sum_{\alpha=1}^{N_i} v_i^\alpha \circ v_i'^\alpha = \text{id}_X, \quad v_i'^\alpha \circ v_j^\beta = \delta_{i,j} \delta_{\alpha,\beta} \text{id}_{X_i}.$$

Lemma 2.8. *Let \mathcal{C} be a semisimple pivotal tensor category with simple unit. Then \mathcal{C} is spherical iff $d(X) = d(\bar{X})$ for all X .*

Proof. We compute

$$\begin{aligned} \text{tr}_L(s) &= \bar{\varepsilon}(X) \circ s \otimes \text{id}_{\bar{X}} \circ \varepsilon(X) \\ &= \sum_{i,\alpha} \sum_{j,\beta} \bar{\varepsilon}(X) \circ (u_i^\alpha \circ u_i'^\alpha \circ s \circ u_j^\beta \circ u_j'^\beta) \otimes \text{id}_{\bar{X}} \circ \varepsilon(X) \\ &= \sum_{i,\alpha} \sum_{j,\beta} \bar{\varepsilon}(X) \circ (u_i^\alpha \circ u_i'^\alpha \circ s \circ u_j^\beta) \otimes v_j^\beta \circ \varepsilon(X_j) \\ &= \sum_{i,\alpha} \sum_{j,\beta} \bar{\varepsilon}(X_j) \circ (u_j'^\beta \circ u_i^\alpha \circ u_i'^\alpha \circ s \circ u_j^\beta) \otimes \text{id}_{\bar{X}_j} \circ \varepsilon(X_j) \\ &= \sum_{i,\alpha} \bar{\varepsilon}(X_i) \circ (u_i'^\alpha \circ s \circ u_i^\alpha) \otimes \text{id}_{\bar{X}_i} \circ \varepsilon(X_i) \\ &= \sum_{i,\alpha} s_i^\alpha \bar{\varepsilon}(X_i) \circ \varepsilon(X_i) = \sum_{i,\alpha} s_i^\alpha d(X_i), \end{aligned}$$

where $s_i^\alpha id_{X_i} = u_i'^\alpha \circ s \circ u_i^\alpha \in \text{Hom}(X_i, X_i)$. We have used (2.2) and the fact that the X_i are simple. In a similar way one computes

$$tr_R(s) = \sum_{i,\alpha} s_i^\alpha \tilde{\varepsilon}(\bar{X}_i) \circ \varepsilon(\bar{X}_i) = \sum_{i,\alpha} s_i^\alpha d(\bar{X}_i)$$

and the result follows from the assumption. \square

We will need the following facts concerning the behavior of sphericity under certain categorical constructions:

Lemma 2.9. *Let \mathcal{A}, \mathcal{B} be pivotal (spherical). Then $\mathcal{A}^{\text{op}}, \mathcal{A} \otimes_{\mathbb{F}} \mathcal{B}$ (the product in the sense of enriched category theory) and the completions $\tilde{\mathcal{A}}^\oplus, \tilde{\mathcal{A}}^p$ are pivotal (spherical).*

Proof. We only sketch the proof, restricting us to the strict pivotal/spherical case. For the opposite category the claimed facts are obvious, choose $\varepsilon(X^{\text{op}}) = \tilde{\varepsilon}(X)$, etc. As to $\tilde{\mathcal{A}}^p$, recall that its objects are given by pairs (X, p) , $X \in \text{Obj } \mathcal{A}$, $p = p^2 \in \text{End}(X)$ and those of $\tilde{\mathcal{A}}^\oplus$ by finite sequences (X_1, \dots, X_l) of objects in \mathcal{A} . We define the duality maps on $\mathcal{A}^{\text{op}}, \mathcal{A} \otimes_{\mathbb{F}} \mathcal{B}, \tilde{\mathcal{A}}^p, \tilde{\mathcal{A}}^\oplus$ by

$$\overline{X^{\text{op}}} = \tilde{X}^{\text{op}}, \quad \overline{X \boxtimes Y} = \tilde{X} \boxtimes \tilde{Y}, \quad \overline{(X, p)} = (\tilde{X}, \tilde{p}), \quad \overline{(X_1, \dots, X_l)} = (\tilde{X}_1, \dots, \tilde{X}_l),$$

respectively. Conditions (2.1) are clearly satisfied. We define further

$$\varepsilon(X^{\text{op}}) = \tilde{\varepsilon}(X)^{\text{op}}, \quad \varepsilon(X \boxtimes Y) = \varepsilon(X) \boxtimes \varepsilon(Y), \quad \varepsilon((X, p)) = p \otimes \tilde{p} \circ \varepsilon(X)$$

and

$$\varepsilon((X_1, \dots, X_l)) = \sum_{i=1}^l u_i \otimes \bar{u}_i \circ \varepsilon(X_i),$$

where the $u_i : X_i \rightarrow (X_1, \dots, X_l)$ are the injections which together with the $\{u'_i\}$ satisfy (2.3). The easy verification of the axioms is left to the reader. It is clear that the new categories are spherical if \mathcal{A}, \mathcal{B} are spherical. \square

The definition of strict pivotal tensor categories can be generalized to 2-categories \mathcal{E} in a straightforward manner. Every 1-morphism $X : \mathcal{U} \rightarrow \mathcal{V}$ has a two-sided dual $\bar{X} : \mathcal{V} \rightarrow \mathcal{U}$. This map has properties which are obvious generalizations of the monoidal situation. In particular, $\overline{1_{\mathfrak{X}}} = 1_{\bar{\mathfrak{X}}}$ for all objects \mathfrak{X} . For $X \in \text{Hom}(\mathcal{U}, \mathcal{V})$ there are $\varepsilon(X) : 1_{\mathcal{V}} \rightarrow X\bar{X}$ and $\tilde{\varepsilon}(X) : X\bar{X} \rightarrow 1_{\mathcal{U}}$ satisfying conditions which are analogous to those in pivotal categories. Again the conjugation can be extended to a 2-functor $\bar{} : \mathcal{E}^{\text{op}} \rightarrow \mathcal{E}$, where \mathcal{E}^{op} has the same objects as \mathcal{E} and 1- and 2-morphisms are reversed. (This functor acts trivially on the objects.) All this can be obtained from [43] by ignoring the monoidal structure on \mathcal{E} considered there. There is one difference, however, which requires attention. For $X : \mathcal{U} \rightarrow \mathcal{V}$ and $s \in \text{End}(X)$ the two traces

$$tr_L(s) = \tilde{\varepsilon}(X) \circ s \otimes id_{\bar{X}} \circ \varepsilon(X), \quad tr_R(s) = \tilde{\varepsilon}(\bar{X}) \circ id_X \otimes s \circ \varepsilon(\bar{X})$$

take values in different commutative monoids, viz. $\text{End}(\mathbf{1}_{\mathfrak{B}})$ and $\text{End}(\mathbf{1}_{\mathfrak{U}})$, respectively. Thus, defining spherical 2-categories as satisfying $\text{tr}_L(s) = \text{tr}_R(s)$ for all 2-morphisms $s \in \text{End}(X)$ —not only those where X has equal source and range—makes sense only if there is a way to identify the $\text{End}(\mathbf{1}_{\mathfrak{X}})$ for different \mathfrak{X} . We will restrict ourselves to 2-categories where $\mathbf{1}_{\mathfrak{X}}$ is absolutely simple for all objects \mathfrak{X} , thus all $\text{End}(\mathbf{1}_{\mathfrak{X}})$ are canonically isomorphic to \mathbb{F} via $c \mapsto c \text{ id}_{\mathbf{1}_{\mathfrak{X}}}$. (The bicategories \mathcal{BIM} and \mathcal{MOR} of bimodules and $*$ -homomorphisms of factors discussed in the introductions provide examples.) Under this condition every 1-morphism has an \mathbb{F} -valued dimension with the expected properties.

2.4. Duality in $*$ -categories

In this section, we posit $\mathbb{F} = \mathbb{C}$. Many complex linear categories have an additional piece of structure: a positive $*$ -operation. See [12,70] for two important classes of examples. A $*$ -operation on a \mathbb{C} -linear category is map which assigns to every morphism $s: X \rightarrow Y$ a morphism $s^*: Y \rightarrow X$. This map has to be antilinear, involutive ($(s^{**} = s)$), contravariant ($((s \circ t)^* = t^* \circ s^*)$) and monoidal ($((s \otimes t)^* = s^* \otimes t^*)$) if the category is monoidal. A $*$ -operation is called positive if $s^* \circ s = 0$ implies $s = 0$. A *tensor $*$ -category* is a \mathbb{C} -linear tensor category with a positive $*$ -operation. (For a braided tensor $*$ -category one often requires unitarity of the braiding, but there are examples where this is not satisfied by a naturally given braiding.) Some relevant references, where infinite-dimensional *Hom*-sets are admitted, are [21,12,42], the latter reference containing a very useful discussion of 2- $*$ -categories. In [46, Proposition 2.1] we showed that \mathbb{C} -linear categories with positive $*$ -operation and finite-dimensional *Hom*-sets are C^* - and W^* -categories in the sense of [21,12]. In W^* -categories one has a polar decomposition theorem for morphisms [21, Corollary 2.7], which implies, e.g. that if $\text{Hom}(Y, X)$ contains an invertible morphism then it contains a unitary morphism ($u \circ u^* = u^* \circ u = \text{id}$).

The notion of duality in $*$ -categories as considered in [12,42] has two peculiar features. First, it is automatically two-sided. Secondly, there is no compelling reason to fix a duality map $X \mapsto \bar{X}$ and to choose morphisms $\mathbf{1} \rightarrow X\bar{X}$, etc. Rather it is sufficient to assume that all objects *have* a conjugate. We stick to the term ‘conjugate’ from [12,42] in order to underline the conceptual difference.

An object \bar{X} is said to be a conjugate of X if there are $r_X \in \text{Hom}(\mathbf{1}, \bar{X}X)$, $\bar{r}_X \in \text{Hom}(\mathbf{1}, X\bar{X})$ satisfying the *conjugate equations*:

$$\bar{r}_X^* \otimes \text{id}_X \circ \text{id}_X \otimes r_X = \text{id}_X, \quad r_X^* \otimes \text{id}_{\bar{X}} \circ \text{id}_{\bar{X}} \otimes \bar{r}_X = \text{id}_{\bar{X}}. \quad (2.4)$$

A category \mathcal{C} has conjugates if every object $X \in \mathcal{C}$ has a conjugate $\bar{X} \in \mathcal{C}$. The triple $(\bar{X}, r_X, \bar{r}_X)$ is called a solution of the conjugate equations. It is called normalized if $r_X^* \circ r_X = \bar{r}_X^* \circ \bar{r}_X \in \text{Hom}(\mathbf{1}, \mathbf{1})$. (Since \mathbb{C} is algebraically closed every solution can be normalized.) A solution of the conjugate equations is called standard if $r_X = \sum_i \bar{w}_i \otimes w_i \circ r_i$ where w_i, \bar{w}_i are isometries effecting decompositions $X = \oplus_i X_i$, $\bar{X} = \oplus_i \bar{X}_i$ into simple objects, and $(\bar{X}_i, r_i, \bar{r}_i)$ are normalized solutions of (2.4) for X_i . For any object X we define a dimension by $d(X) = r_X^* \circ r_X$, where $(\bar{X}, r_X, \bar{r}_X)$ is a normalized standard solution. Then $d(X)$ is well defined and satisfies $d(X) = d(\bar{X})$, $d(X \oplus Y) = d(X) + d(Y)$

and $d(X \otimes Y) = d(X)d(Y)$ for all $X, Y \in \mathcal{C}$. The dimension takes values in the set $\{2 \cos \pi/n, n=3, 4, \dots\} \cup [2, \infty)$. If (\bar{X}, r, \bar{r}) is any normalized solution of the conjugate equations for X then $d(X) \leq r^* \circ r$ and equality holds iff (\bar{X}, r, \bar{r}) is standard. For the proofs we refer to [42].

Furthermore, a braided $*$ -category with conjugates automatically [46] has a canonical twist [28]. Together with the fact [77] that in braided tensor categories there is a one-to-one correspondence between sovereign structures and twists this implies that every braided $*$ -category has a canonical sovereign structure. It is not unreasonable to believe that this is true even in the absence of a braiding. In fact, in [72] Yamagami considered spherical structures (‘ ε -structures’) compatible with a given $*$ -structure, and in [74] he shows that every $*$ -category can be equipped with an essentially unique spherical structure. See also [17] for similar considerations. In any case, in the $*$ -case one does not need a spherical structure since every scalar quantity (i.e. morphism $\mathbf{1} \rightarrow \mathbf{1}$) is unambiguously defined if one sticks to the above normalization rules for the duality morphisms r_X, \bar{r}_X , cf. [42].

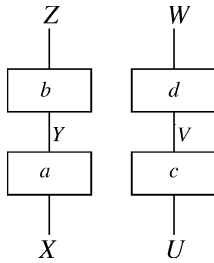
2.5. Graphical notation

In this paper, we will often represent computations with morphisms in a tensor category in terms of tangle diagrams rather than by formulas or commutative diagrams. Since this notation is well known, cf. e.g. [30], we just explain our conventions.

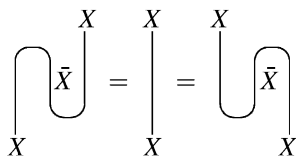
Our diagrams are to be read upwards, thus with $a: X \rightarrow Y, b: Y \rightarrow Z, c: U \rightarrow V, d: V \rightarrow W$ we represent

$$b \otimes d \circ a \otimes c = (b \circ a) \otimes (d \circ c) \in \text{Hom}(X \otimes U, Z \otimes W)$$

by



Representing $\varepsilon(X)$ by $\begin{array}{c} X \\ \curvearrowright \\ \bar{X} \end{array}$ and $\bar{\varepsilon}(X)$ by $\begin{array}{c} \bar{X} \\ \curvearrowleft \\ X \end{array}$ we depict condition (1) of Definition 2.6 by



3.1. From two-sided duals to Frobenius algebras

Definition 3.1. Let \mathcal{A} be a (strict) tensor category. A Frobenius algebra in \mathcal{A} is a quintuple (Q, v, v', w, w') , where Q is an object in \mathcal{A} and $v: \mathbf{1} \rightarrow Q, v': Q \rightarrow \mathbf{1}, w: Q \rightarrow Q^2, w': Q^2 \rightarrow Q$ are morphisms satisfying the following conditions:

$$w' \otimes id_O \circ id_O \otimes w = w \circ w' = id_O \otimes w' \circ w \otimes id_O. \quad (3.5)$$

$$\begin{array}{c} \cup \\ | \end{array} = w, \quad \begin{array}{c} | \\ \cap \end{array} = w', \quad \begin{array}{c} | \\ \circ \end{array} = v, \quad \begin{array}{c} \circ \\ | \end{array} = v'.$$

For the tangle diagrams corresponding to the above conditions see Figs. 1 and 2.

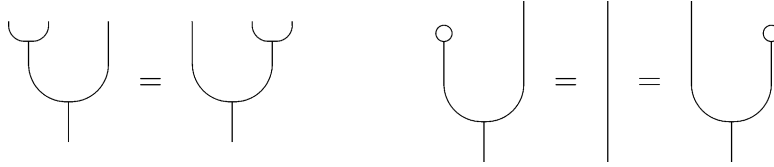


Fig. 1. Comonoids in a category.

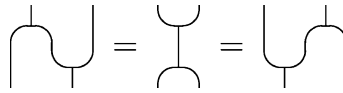


Fig. 2. The Frobenius condition.

2. Eqs. (3.1)–(3.4) amount to requiring that (Q, w', v) and (Q, w, v') are a monoid and a comonoid, respectively, in the category \mathcal{C} . The new ingredient is the Frobenius condition (3.5), cf. Fig. 2, which can be interpreted as expressing that $w: Q \rightarrow Q^2$ is a map of Q – Q bimodules. This must not be confused with the more familiar bialgebra condition. The latter makes sense only if \mathcal{C} comes with a braiding, which we do not assume. For this reason we avoid the usual symbols $m, \Delta, \varepsilon, \eta$.
3. To the best of our knowledge (3.5) makes its first appearance in [58, Appendix A.3], where it is part of an alternative characterization of symmetric algebras (in Vect). This is a special case of the alternative characterization of Frobenius algebras mentioned in the Introduction and discussed in more detail in Section 6.
4. If \mathcal{A} is a $*$ -category we will later on require $w' = w^*$, $v' = v^*$, which obviously renders (3.2), (3.4) redundant.

Definition 3.3. Frobenius algebras (Q, v, v', w, w') , $(\tilde{Q}, \tilde{v}, \tilde{v}', \tilde{w}, \tilde{w}')$ in the (strict) tensor category \mathcal{A} are isomorphic if there is an isomorphism $s': Q \rightarrow \tilde{Q}$ such that

$$s \circ v = \tilde{v}, \quad v' = \tilde{v}' \circ s, \quad s \otimes s \circ w = \tilde{w} \circ s, \quad s \circ w' = \tilde{w}' \circ s \otimes s.$$

The above definitions are vindicated by the following.

Lemma 3.4. Let $J: \mathfrak{B} \rightarrow \mathfrak{A}$ be a 1-morphism in a 2-category \mathcal{C} and let $\bar{J}: \mathfrak{A} \rightarrow \mathfrak{B}$ be a two-sided dual with duality 2-morphisms $d_J, e_J, \varepsilon_J, \eta_J$. Positing $Q = J\bar{J}: \mathcal{A} \rightarrow \mathcal{A}$ there are v, v', w, w' such that (Q, v, v', w, w') is a Frobenius algebra in the tensor category $\mathcal{A} = \text{HOM}_{\mathcal{C}}(\mathfrak{A}, \mathfrak{A})$.

Proof. Since J, \bar{J} are mutually two-sided duals there are

$$e_J: \mathbf{1}_{\mathfrak{A}} \rightarrow J\bar{J}, \quad \eta_J: J\bar{J} \rightarrow \mathbf{1}_{\mathfrak{A}}, \quad \varepsilon_J: \mathbf{1}_{\mathfrak{B}} \rightarrow \bar{J}J, \quad d_J: \bar{J}J \rightarrow \mathbf{1}_{\mathfrak{B}}$$

satisfying

$$\begin{aligned} id_J \otimes d_J \otimes e_J \otimes id_J &= id_J = \eta_J \otimes id_J \otimes id_J \otimes \eta_J, \\ id_{\bar{J}} \otimes \eta_J \otimes \varepsilon_J \otimes id_{\bar{J}} &= id_{\bar{J}} = d_J \otimes id_{\bar{J}} \otimes id_{\bar{J}} \otimes e_J. \end{aligned} \tag{3.6}$$

Defining

$$\begin{aligned}
 v &= e_J \in \text{Hom}_{\mathcal{E}}(\mathbf{1}_{\mathfrak{U}}, J\bar{J}), \\
 v' &= \eta_J \in \text{Hom}_{\mathcal{E}}(J\bar{J}, \mathbf{1}_{\mathfrak{U}}), \\
 w &= id_J \otimes \varepsilon_J \otimes id_{\bar{J}} \in \text{Hom}_{\mathcal{E}}(J\bar{J}, J\bar{J}J\bar{J}), \\
 w' &= id_J \otimes d_J \otimes id_{\bar{J}} \in \text{Hom}_{\mathcal{E}}(J\bar{J}J\bar{J}, J\bar{J}),
 \end{aligned} \tag{3.7}$$

we have $v \in \text{Hom}_{\mathcal{A}}(\mathbf{1}, Q)$, $w \in \text{Hom}_{\mathcal{A}}(Q, Q^2)$, $v' \in \text{Hom}_{\mathcal{A}}(Q, \mathbf{1})$ and $w' \in \text{Hom}_{\mathcal{A}}(Q^2, Q)$. Now (3.1) follows simply from the interchange law

and (3.3) from the duality equations (3.6):

Conditions (3.2), (3.4) are analogous, and the Frobenius algebra condition (3.5) follows from

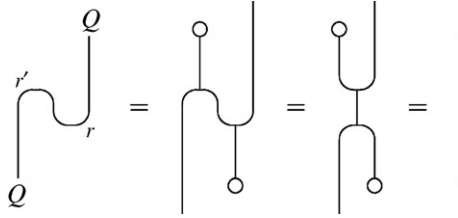
Remark 3.5. 1. By duality we obviously also obtain a Frobenius algebra $\hat{Q} = (\hat{Q} = \bar{J}J, \dots)$ in $\mathcal{B} = \text{HOM}_{\mathcal{E}}(\mathfrak{B}, \mathfrak{B})$.

2. Considering a tensor category \mathcal{C} as a 2-category with a single object \mathfrak{U} we obtain the special case of an object X in \mathcal{C} with a two-sided dual \bar{X} .

3. Since duals are unique up to isomorphism a different choice of \bar{J} changes Q only within its isomorphism class. Yet it is in general not true that the Frobenius algebra (Q, v, v', w, w') is well defined up to isomorphism. \square

Lemma 3.6. *The object Q of a Frobenius algebra is two-sided self-dual, i.e. there are $r \in \text{Hom}(\mathbf{1}, Q^2)$, $r' \in \text{Hom}(Q^2, \mathbf{1})$ such that the duality equations are satisfied with $e_Q = \varepsilon_Q = r$, $d_Q = \eta_Q = r'$.*

Proof. Set $r = w \circ v$, $r' = v' \circ w'$. Then the duality equations follow by



Here the first equality holds by definition, the second follows from (3.5) and the last from (3.3). \square

3.2. A universal construction

Now we prove a converse of Lemma 3.4 by embedding a given tensor category \mathcal{A} with a Frobenius algebra (Q, v, v', w, w') into a suitably constructed 2-category such that $Q = J\bar{J}$.

Definition 3.7. An almost-2-category is defined as a 2-category [31] except that we do not require the existence of a unit 1-morphism $\mathbf{1}_{\mathfrak{X}}$ for every object \mathfrak{X} .

Proposition 3.8. *Let \mathcal{A} be a strict tensor category and $Q = (Q, v, v', w, w')$ a Frobenius algebra in \mathcal{A} . Then there is an almost-2-category \mathcal{E}_0 satisfying:*

1. $\text{Obj } \mathcal{E}_0 = \{\mathfrak{U}, \mathfrak{B}\}$.
2. There is an isomorphism $I : \mathcal{A} \rightarrow \text{Hom}_{\mathcal{E}_0}(\mathfrak{U}, \mathfrak{U})$ of tensor categories.
3. There are 1-morphisms $J : \mathfrak{B} \rightarrow \mathfrak{U}$ and $\bar{J} : \mathfrak{U} \rightarrow \mathfrak{B}$ such that $J\bar{J} = I(Q)$.

If \mathcal{A} is \mathbb{F} -linear then so is \mathcal{E}_0 . Isomorphic Frobenius algebras give rise to isomorphic almost 2-categories.

Proof. The proof is constructive, the definition of the objects obviously being forced upon us: $\text{Obj } \mathcal{E}_0 = \{\mathfrak{U}, \mathfrak{B}\}$.

1-morphisms: We define formally

$$\begin{aligned}
 \text{Hom}_{\mathcal{E}_0}(\mathfrak{U}, \mathfrak{U}) &= \text{Obj } \mathcal{A}, \\
 \text{Hom}_{\mathcal{E}_0}(\mathfrak{B}, \mathfrak{U}) &= \{“XJ”, X \in \text{Obj } \mathcal{A}\}, \\
 \text{Hom}_{\mathcal{E}_0}(\mathfrak{U}, \mathfrak{B}) &= \{“\bar{J}X”, X \in \text{Obj } \mathcal{A}\}, \\
 \text{Hom}_{\mathcal{E}_0}(\mathfrak{B}, \mathfrak{B}) &= \{“\bar{J}XJ”, X \in \text{Obj } \mathcal{A}\}.
 \end{aligned} \tag{3.8}$$

For the moment this simply means that $\text{Hom}_{\mathcal{E}_0}(\mathcal{U}, \mathfrak{B})$, $\text{Hom}_{\mathcal{E}_0}(\mathfrak{B}, \mathcal{U})$ and $\text{Hom}_{\mathcal{E}_0}(\mathfrak{B}, \mathfrak{B})$ are isomorphic to $\text{Hom}_{\mathcal{A}}(\mathcal{U}, \mathcal{U})$ as sets. In particular, with $X = \mathbf{1} \in \text{Obj } \mathcal{A}$ we obtain the following distinguished 1-morphisms:

$$J = \text{“}\mathbf{1}J\text{”} \in \text{Hom}_{\mathcal{E}_0}(\mathfrak{B}, \mathcal{U}),$$

$$\bar{J} = \text{“}\bar{J}\mathbf{1}\text{”} \in \text{Hom}_{\mathcal{E}_0}(\mathcal{U}, \mathfrak{B}).$$

Composition of 1-morphisms: Wherever legal, it is defined by juxtaposition, followed by replacing a possibly occurring composite $J\bar{J}$ by the underlying object Q of the Frobenius algebra. (The latter is the case whenever one considers $X \circ Y$ where $Y \in \text{Hom}_{\mathcal{E}_0}(\mathfrak{B}, \mathfrak{B})$, $X \in \text{Hom}_{\mathcal{E}_0}(\mathfrak{B}, \mathfrak{B})$.) We refrain from tabulating all the possibilities and give instead a few examples:

$$\text{“}\bar{J}X\text{”} \circ \text{“}YJ\text{”} := \text{“}\bar{J}(XY)J\text{”} \in \text{Hom}_{\mathcal{E}_0}(\mathfrak{B}, \mathfrak{B}),$$

$$\text{“}XJ\text{”} \circ \text{“}\bar{J}YJ\text{”} := \text{“}(XQY)J\text{”} \in \text{Hom}_{\mathcal{E}_0}(\mathfrak{B}, \mathcal{U}).$$

\mathcal{A} being strict by assumption, the composition of 1-morphisms is obviously strictly associative. With this definition the set of 1-morphisms is the free semigroupoid (= small category) generated by $\text{Obj } \mathcal{A} \cup \{J, \bar{J}\}$ modulo $J\bar{J} = Q$ and the relations in \mathcal{A} . If we had to consider only 1-morphisms we could drop the quotes in (3.8) since now J, \bar{J} are legal 1-morphisms and, e.g. “ XJ ” is just the composition $X \circ J$ of X and J . But in order to define the 2-morphisms and their compositions and to verify that we obtain a 2-category we must continue to distinguish between “ X ” “ J ” and “ XJ ” for a while.

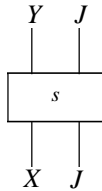
2-morphisms as sets: Since we want $\text{HOM}(\mathcal{U}, \mathcal{U}) \cong \mathcal{A}$, we clearly have to set

$$\text{Hom}_{\mathcal{E}_0}(X, Y) = \text{Hom}_{\mathcal{A}}(X, Y), \quad X, Y \in \text{Hom}_{\mathcal{E}_0}(\mathcal{U}, \mathcal{U}) = \text{Obj } \mathcal{A}.$$

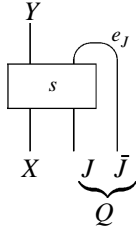
In order to identify the remaining 2-morphisms we appeal to duality which in the end should hold in \mathcal{E} . Applied to $\text{Hom}_{\mathcal{E}_0}(\mathfrak{B}, \mathcal{U})$ this means

$$\begin{aligned} \text{Hom}_{\mathcal{E}_0}(\text{“}X \circ J\text{”}, \text{“}Y \circ J\text{”}) &\cong \text{Hom}_{\mathcal{E}_0}(\text{“}X \circ J \circ \bar{J}\text{”}, Y) = \text{Hom}_{\mathcal{E}_0}(XQ, Y) \\ &= \text{Hom}_{\mathcal{A}}(XQ, Y). \end{aligned}$$

This means that the elements of $\text{Hom}_{\mathcal{E}_0}(\text{“}XJ\text{”}, \text{“}YJ\text{”})$



where $X, Y \in \text{Obj } \mathcal{A}$, are represented by those of $\text{Hom}_{\mathcal{A}}(XQ, Y)$, etc:



Thus we define

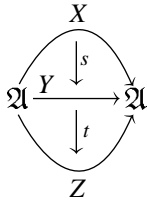
$$\text{Hom}_{\mathcal{E}_0}("XJ", "YJ") = \text{Hom}_{\mathcal{A}}(XQ, Y),$$

$$\text{Hom}_{\mathcal{E}_0}("JX", "JY") = \text{Hom}_{\mathcal{A}}(X, QY),$$

$$\text{Hom}_{\mathcal{E}_0}("JXJ", "JYJ") = \text{Hom}_{\mathcal{A}}(XQ, QY),$$

as sets. Now we must define the vertical and horizontal compositions of the 2-morphisms in \mathcal{E}_0 , which we denote \bullet and \times , respectively, in order to avoid confusion with the compositions in \mathcal{A} .

Vertical (\bullet -) Composition of 2-morphisms: Let $X, Y, Z \in \text{Obj } \mathcal{A}$. For $s \in \text{Hom}_{\mathcal{E}_0}(X, Y) \equiv \text{Hom}_{\mathcal{A}}(X, Y)$, $t \in \text{Hom}_{\mathcal{E}_0}(Y, Z) \equiv \text{Hom}_{\mathcal{A}}(Y, Z)$ it is clear that $t \bullet s = t \circ s$.



Let now $s \in \text{Hom}_{\mathcal{E}_0}("X \circ J", "Y \circ J") \equiv \text{Hom}_{\mathcal{A}}(XQ, Y)$, $t \in \text{Hom}_{\mathcal{E}_0}("Y \circ J", "Z \circ J") \equiv \text{Hom}_{\mathcal{A}}(YQ, Z)$. Then we define $t \bullet s \in \text{Hom}_{\mathcal{E}_0}("X \circ J", "Z \circ J") \equiv \text{Hom}_{\mathcal{A}}(XQ, Z)$ by

$$t \bullet s = t \circ s \otimes \text{id}_Q \circ \text{id}_X \otimes w.$$

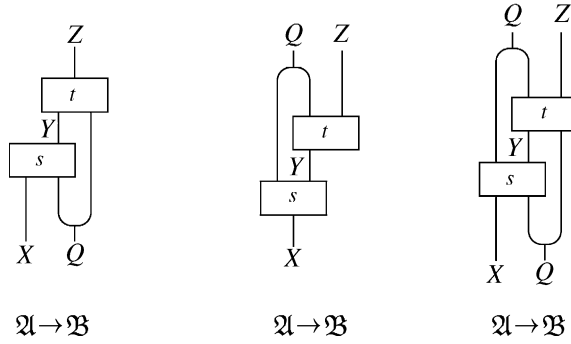
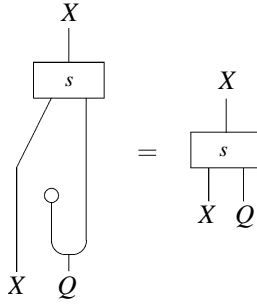
Similarly, for $s \in \text{Hom}_{\mathcal{E}_0}("J \circ X", "J \circ Y") \equiv \text{Hom}_{\mathcal{A}}(X, QY)$, $t \in \text{Hom}_{\mathcal{E}_0}("J \circ Y", "J \circ Z") \equiv \text{Hom}_{\mathcal{A}}(Y, QZ)$ we define $t \bullet s \in \text{Hom}_{\mathcal{E}_0}("J \circ X", "J \circ Z") \equiv \text{Hom}_{\mathcal{A}}(X, QZ)$ by

$$t \bullet s = w' \otimes \text{id}_Z \circ \text{id}_Q \otimes t \circ s.$$

Finally, for $s \in \text{Hom}_{\mathcal{E}_0}("J \circ X \circ J", "J \circ Y \circ J") \equiv \text{Hom}_{\mathcal{A}}(XQ, QY)$, $t \in \text{Hom}_{\mathcal{E}_0}("J \circ Y \circ J", "J \circ Z \circ J") \equiv \text{Hom}_{\mathcal{A}}(YQ, QZ)$ we define $t \bullet s \in \text{Hom}_{\mathcal{E}_0}("J \circ X \circ J", "J \circ Z \circ J") \equiv \text{Hom}_{\mathcal{A}}(XQ, QZ)$ by

$$t \bullet s = w' \otimes \text{id}_Z \circ \text{id}_Q \otimes t \circ s \otimes \text{id}_Q \circ \text{id}_X \otimes w. \quad (3.9)$$

See Fig. 3 for diagrams corresponding to these definitions. Associativity of the \bullet -compositions is easily verified using coassociativity (3.1) and associativity (3.2) of the Frobenius algebra (Q, v, v', w, w') .

Fig. 3. Vertical (\bullet -) composition of 2-morphisms in \mathcal{E}_0 .Fig. 4. Unit 2-morphism id_{XJ} .

Unit 2-arrows: It is clear $id_X \in Hom_{\mathcal{A}}(X, X)$ for $X \in Obj \mathcal{A}$ is the unit 2-arrow for the 1-morphism $X : \mathcal{U} \rightarrow \mathcal{U}$. Furthermore, using the above rules for the \bullet -composition of 2-morphisms and Eqs. (3.3), (3.4), it is easily verified that

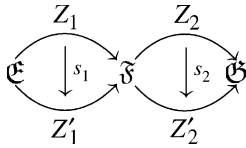
$$id_{XJ} = id_X \otimes v' \in Hom_{\mathcal{A}}(XQ, X) \equiv Hom_{\mathcal{E}_0}("XJ", "XJ")$$

is in fact the unit 2-arrow id_{XJ} . Diagrammatically, the equation $s \bullet id_{XJ} = s$ with $s \in Hom_{\mathcal{A}}(XQ, X) \equiv Hom_{\mathcal{E}_0}("XJ", "XJ")$ looks as in Fig. 4. Similarly, we have

$$id_{JX} = v \otimes id_X \in Hom_{\mathcal{A}}(X, QX) \equiv Hom_{\mathcal{E}_0}("JX", "JX"),$$

$$id_{JXJ} = v \otimes id_X \otimes v' \in Hom_{\mathcal{A}}(XQ, QX) \equiv Hom_{\mathcal{E}_0}("JXJ", "JXJ").$$

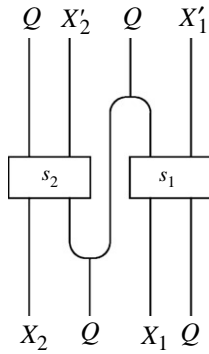
Horizontal (\times -) Composition of 2-morphisms. Let $\mathfrak{E}, \mathfrak{F}, \mathfrak{G} \in \{\mathcal{U}, \mathcal{B}\}$, $Z_1, Z'_1 \in Hom_{\mathcal{E}_0}(\mathfrak{E}, \mathfrak{F})$, $Z_2, Z'_2 \in Hom_{\mathcal{E}_0}(\mathfrak{F}, \mathfrak{G})$ and $s_i \in Hom_{\mathcal{E}_0}(Z_i, Z'_i)$, $i = 1, 2$.



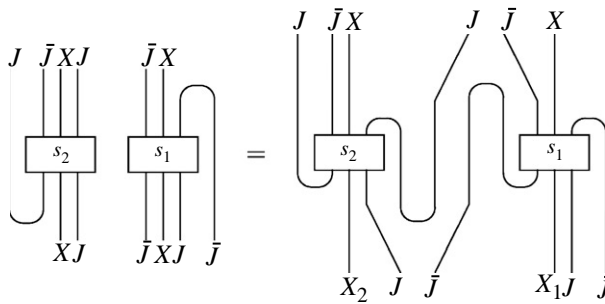
Then we must define $s_2 \times s_1 \in \text{Hom}_{\mathcal{E}_0}(Z_2 \circ Z_1, Z'_2 \circ Z'_1)$. To this purpose we consider, as before, s_i as morphisms in \mathcal{A} . Thus, for $\mathfrak{E} = \mathfrak{F} = \mathfrak{U}$ we have $Z_1, Z'_1 \in \text{Obj } \mathcal{A}$ and $s_1 \in \text{Hom}_{\mathcal{A}}(Z_1, Z'_1)$, and for $\mathfrak{E} = \mathfrak{B}, \mathfrak{F} = \mathcal{A}$ we have $Z_1 = X_1 J, Z'_1 = X'_1 J$ with $X_1, X'_1 \in \text{Obj } \mathcal{A}$ and $s_1 \in \text{Hom}_{\mathcal{E}}(X_1 J, X'_1 J) \equiv \text{Hom}_{\mathcal{A}}(X_1 Q, X'_1)$, etc. We define

$$\begin{aligned} s_2 \times s_1 &= s_2 \otimes s_1 & \text{if } \mathfrak{F} = \mathfrak{U}, \\ s_2 \times s_1 &= id_{\gamma} \otimes w' \otimes id_{\gamma} \circ s_2 \otimes id_Q \otimes s_1 \circ id_{\gamma} \otimes w \otimes id_{\gamma} & \text{if } \mathfrak{F} = \mathfrak{B}. \end{aligned} \quad (3.10)$$

To illustrate the second equation, consider the case $\mathfrak{E} = \mathfrak{F} = \mathfrak{G} = \mathfrak{B}$, thus $Z_i^{(\prime)} = \bar{J} \circ X_i^{(\prime)} \circ J, i = 1, 2$, with $X_i^{(\prime)} \in \mathcal{A}$ and $s_i \in \text{Hom}_{\mathcal{A}}(X_i Q, Q X_i^{(\prime)})$. Then (3.10) looks like



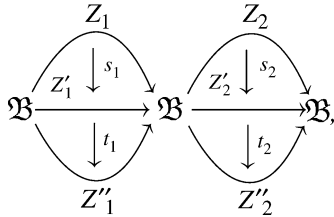
Our task is now twofold. On the one hand, we must convince the reader that this is the ‘right’ definition. We remark that definition (3.10) is motivated by the interpretation of the 2-morphisms of \mathcal{E}_0 in terms of morphisms in \mathcal{A} . For the horizontal composition of 2-morphisms where the intermediate object \mathfrak{F} is \mathfrak{B} this looks as follows:



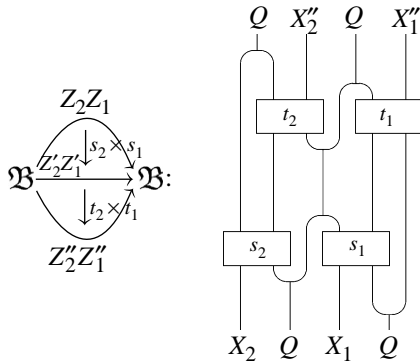
It should be clear that for $\mathfrak{F} = \mathfrak{U}$ the first formula in (3.10) is the correct definition.

The second task is, of course, to prove that our definition renders \mathcal{E}_0 an almost-2-category. This means the horizontal composition of three 2-morphisms must be associative for all legal compositions. Luckily, this is quite obvious from associativity of the \otimes -composition in the tensor category \mathcal{A} and we refrain from formalizing this. It remains to show the interchange law, which again we do only for $\mathfrak{E} = \mathfrak{F} = \mathfrak{G} = \mathfrak{B}$, all

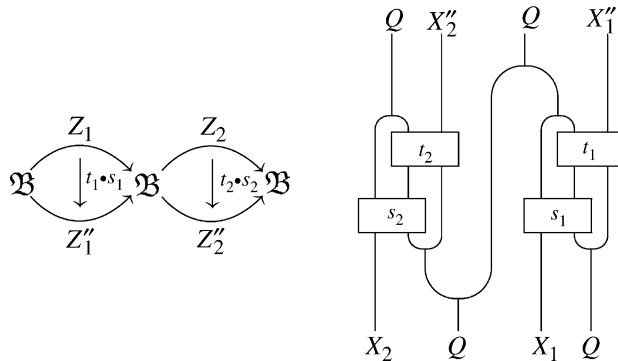
other cases being easier. Let thus $Z_i = JX_i\bar{J}$, $s_1 \in \text{Hom}_{\mathcal{E}}(Z_1, Z'_1) \equiv \text{Hom}_{\mathcal{A}}(X_1Q, QX'_1)$, etc. We compute the composition



which is in $\text{Hom}_{\mathcal{E}}(Z_2Z_1, Z_2Z'_1) \equiv \text{Hom}_{\mathcal{A}}(X_2QX_1Q, QX_2QX'_1)$, in two different ways. We can first do the horizontal compositions and obtain



Beginning with the vertical compositions we arrive at



That the two expressions coincide is again verified easily using (3.1)–(3.5).

Assume now that \mathcal{A} is \mathbb{F} -linear. Clearly, the spaces of 2-morphisms in \mathcal{E}_0 are \mathbb{F} -vector spaces and the compositions \bullet, \times are bilinear. Thus \mathcal{E}_0 is \mathbb{F} -linear. Finally, let $s: Q \rightarrow \tilde{Q}$ be an isomorphism between the Frobenius algebras Q, \tilde{Q} and consider

the almost 2-categories $\mathcal{E}_0, \tilde{\mathcal{E}}_0$ constructed from Q, \tilde{Q} . The objects and 1-morphisms (as well as the composition of the latter) of $\mathcal{E}_0, \tilde{\mathcal{E}}_0$ not depending on the Frobenius algebra, there are obvious bijections. Furthermore, there is a bijection between, e.g. $Hom_{\mathcal{E}_0}("XJ", "YJ") \equiv Hom_{\mathcal{A}}(XQ, Y)$ and $Hom_{\tilde{\mathcal{E}}_0}("XJ", "YJ") \equiv Hom_{\mathcal{A}}(X\tilde{Q}, Y)$ given by

$$Hom_{\mathcal{A}}(XQ, Y) \ni t \mapsto t \circ s^{-1} \otimes id_Y \in Hom_{\mathcal{A}}(X\tilde{Q}, Y).$$

Since these isomorphisms commute with the compositions \bullet, \times of 2-morphisms in $\mathcal{E}_0, \tilde{\mathcal{E}}_0$ we have an isomorphism of almost 2-categories. \square

Remark 3.9. 1. It is obvious how to modify the proposition if \mathcal{A} is non-strict: The definition of the objects and 1-morphisms and the composition of the latter are unchanged. Since we may still require $J\tilde{J} = Q$, the associativity constraint of \mathcal{A} gives rise to that of \mathcal{E}_0 . As to the 2-morphisms, the only change are appropriate insertions of associativity morphisms in the definitions of \bullet, \times .

2. If $Q = (Q, \dots)$ is a Frobenius algebra in a tensor category \mathcal{A} then the functor $F = Q \otimes -$ is part of a Frobenius algebra in $End \mathcal{A}$, thus in particular of a monoid in $End \mathcal{A}$, equivalently a monad $(Q \otimes -, \{w' \otimes id_X\}, \{v \otimes id_X\})$ in \mathcal{A} . It is easy to verify that our construction of the category $HOM_{\mathcal{E}_0}(\mathfrak{U}, \mathfrak{B})$ is precisely the Kleisli construction [44, Section VI.5] starting from \mathcal{A} and the monad (F, \dots) . (Alternatively one may invoke the Kleisli type construction for monoids in 2-categories, cf. e.g. [22].) Similar statements hold for the categories $HOM_{\mathcal{E}}(\mathfrak{B}, \mathfrak{U})$ and $HOM_{\mathcal{E}}(\mathfrak{B}, \mathfrak{B})$. But our way of pasting everything together in order to obtain a bicategory seems to be new.

3. Just as the Kleisli category is the smallest solution to a certain problem (i.e. an initial object in the category of all adjunctions producing the given monad), it is intuitively clear that \mathcal{E}_0 is an initial object in the category of all solutions of our problem. Consider a 2-category \mathcal{E}' with $\{\mathfrak{U}', \mathfrak{B}'\} \subset Obj \mathcal{E}'$ such that $HOM_{\mathcal{E}'}(\mathfrak{U}', \mathfrak{U}') \cong \mathcal{A}$ (with given isomorphism) and with mutually dual $J' \in Hom_{\mathcal{E}'}(\mathfrak{B}', \mathfrak{U}')$, $\tilde{J}' \in Hom_{\mathcal{E}'}(\mathfrak{U}', \mathfrak{B}')$, such that $Q = J' \tilde{J}'$. Then there is a unique functor K of almost-2-categories $K : \mathcal{E}_0 \rightarrow \mathcal{E}'$ such that $F(\mathfrak{U}) = \mathfrak{U}'$, $F(\mathfrak{B}) = \mathfrak{B}'$, $F(J) = J'$, $F(\tilde{J}) = \tilde{J}'$. We omit the details, since in Theorem 3.17 we will prove a more useful uniqueness result.

With the preceding constructions \mathcal{E}_0 is a 2-category up to one defect: there is no unit- $\mathfrak{B} - \mathfrak{B}$ -morphism. We could, of course, try to add one by hand but that would be difficult to do in a consistent manner. Fortunately, it turns out that taking the closure of \mathcal{E}_0 in which idempotent 2-morphisms split automatically provides us with a (non-strict) $\mathfrak{B} - \mathfrak{B}$ -unit $1_{\mathfrak{B}}$ provided the Frobenius algebra (Q, v, v', w, w') satisfies an additional condition.

In all applications we are going to discuss \mathcal{A} is linear over a field \mathbb{F} and $End(\mathbf{1}) \cong \mathbb{F}$. Yet we wish to emphasize the generality of our basic construction. This is why we give the following definition, motivated by considerations in [68, pp. 72–73].

Definition-Proposition 3.10. *A strict (not necessarily preadditive) tensor category \mathcal{A} is $End(\mathbf{1})$ -linear if $\lambda \otimes s = s \otimes \lambda =: \lambda s$ for all $\lambda \in End(\mathbf{1})$ and $s : X \rightarrow Y$. It then follows*

that $\lambda(s \circ t) = (\lambda s) \circ t = s \circ (\lambda t)$ and $\lambda(s \otimes t) = (\lambda s) \otimes t = s \otimes (\lambda t)$. (All this generalizes to non-strict categories.)

Theorem 3.11. *Let \mathcal{A} be a strict tensor category and $Q = (Q, v, v', w, w')$ a Frobenius algebra in \mathcal{A} . Assume that one of the following conditions is satisfied:*

- (a) $w' \circ w = id_Q$.
- (b) \mathcal{A} is $End(\mathbf{1})$ -linear and

$$w' \circ w = \lambda_1 id_Q,$$

where λ_1 is an invertible element of the commutative monoid $End(\mathbf{1})$.

Then the completion $\mathcal{E} = \bar{\mathcal{E}}^p$ of the \mathcal{E}_0 defined in Proposition 3.8 is a bicategory such that

1. $Obj \mathcal{E} = \{\mathfrak{U}, \mathfrak{B}\}$.
2. There is a fully faithful tensor functor $I: \mathcal{A} \rightarrow HOM_{\mathcal{E}}(\mathfrak{U}, \mathfrak{U})$ such that for every $Y \in HOM_{\mathcal{E}}(\mathfrak{U}, \mathfrak{U})$ there is $X \in \mathcal{A}$ such that Y is a retract of $I(X)$. (Thus I is an equivalence if \mathcal{A} has subobjects.)
3. There are 1-morphisms $J: \mathfrak{B} \rightarrow \mathfrak{U}$ and $\bar{J}: \mathfrak{U} \rightarrow \mathfrak{B}$ such that $Q = J\bar{J}$.
4. J and \bar{J} are mutual two-sided duals, i.e. there are 2-morphisms

$$e_J: \mathbf{1}_{\mathfrak{U}} \rightarrow J\bar{J}, \quad \varepsilon_J: \mathbf{1}_{\mathfrak{B}} \rightarrow \bar{J}J, \quad d_J: \bar{J}J \rightarrow \mathbf{1}_{\mathfrak{B}}, \quad \eta_J: J\bar{J} \rightarrow \mathbf{1}_{\mathfrak{U}}$$

satisfying the usual relations.

5. We have the identity

$$I(Q, v, v', w, w') = (J\bar{J}, e_J, \eta_J, id_J \otimes \varepsilon_J \otimes id_{\bar{J}}, id_J \otimes d_J \otimes id_{\bar{J}})$$

of Frobenius algebras in $END_{\mathcal{E}}(\mathfrak{U})$. (In particular, $d_J \circ \varepsilon_J = \lambda_1 id_{\mathbf{1}_{\mathfrak{B}}}$.)

6. If \mathcal{A} is a preadditive (\mathbb{F} -linear) category then \mathcal{E} is a preadditive (\mathbb{F} -linear) 2-category.
7. If \mathcal{A} has direct sums then \mathcal{E} has direct sums of 1-morphisms.

Isomorphic Frobenius algebras Q, \tilde{Q} give rise to isomorphic bicategories $\mathcal{E}, \tilde{\mathcal{E}}$.

Proof. If we are in case (a) put $\lambda_1 = 1$ in the sequel. Then $End(\mathbf{1})$ -linearity will not be needed. We define the bicategory \mathcal{E} as the completion $\bar{\mathcal{E}}^p$. Thus, $Obj \mathcal{E} = \{\mathfrak{U}, \mathfrak{B}\}$ and for $\mathfrak{X}, \mathfrak{Y} \in \{\mathfrak{U}, \mathfrak{B}\}$ the 1-morphisms are

$$Hom_{\mathcal{E}}(\mathfrak{X}, \mathfrak{Y}) = \{(X, p) | X \in Hom_{\mathcal{E}_0}(\mathfrak{X}, \mathfrak{Y}), p = p \bullet p \in Hom_{\mathcal{E}_0}(X, X)\}.$$

Furthermore,

$$\begin{aligned} Hom_{\mathcal{E}}((X, p), (Y, q)) &= \{s \in Hom_{\mathcal{E}_0}(X, Y) | s \bullet p = q \bullet s = s\} \\ &= q \bullet Hom_{\mathcal{E}_0}(X, Y) \bullet p. \end{aligned}$$

In order to alleviate the notation we allow X to denote also (X, id_X) . With this definition it is clear that $(X, p) \prec X \equiv (X, id_X)$. ($p \in Hom_{\mathcal{E}}(X, X)$ is an invertible morphism from (X, p) to X , since also $p \in Hom_{\mathcal{E}}(X, (X, p))$ and $p \bullet p = p = id_{(X, p)}$.)

Exhibiting a unit $\mathfrak{B} - \mathfrak{B}$ -morphism: Recall that $\text{End}_{\mathcal{E}}(\bar{J}J) \equiv \text{End}_{\mathcal{A}}(Q)$ as a vector space and consider the morphism $p_1 \in \text{End}_{\mathcal{E}}(\bar{J}J)$ represented by $\lambda_1^{-1} id_Q \in \text{End}_{\mathcal{A}}(Q)$. (Note that id_Q is the unit of the monoid $\text{End}_{\mathcal{A}}(Q)$, but not of $\text{End}_{\mathcal{E}}(\bar{J}J)$, whose unit is $v \circ v'!$!) As a consequence of condition (b) we see that $p_1 \bullet p_1 = \lambda_1^{-2} w' \circ w = \lambda_1^{-1} id_Q = p_1$. Thus, p_1 is idempotent and $(\bar{J}J, p_1)$ is a $\mathfrak{B} - \mathfrak{B}$ -morphism in \mathcal{E} . We claim that $(\bar{J}J, p_1)$ is a (non-strict) unit $\mathfrak{B} - \mathfrak{B}$ -morphism. In order to see that $\mathbf{1}_{\mathfrak{B}}$ is a right unit we have to show that there are isomorphisms

$$r((XJ, p)) : (XJ, p)\mathbf{1}_{\mathfrak{B}} = (XQJ, p \times p_1) \rightarrow (XJ, p),$$

$$r((\bar{J}XJ, p)) : (\bar{J}XJ, p)\mathbf{1}_{\mathfrak{B}} = (\bar{J}XQJ, p \times p_1) \rightarrow (\bar{J}XJ, p)$$

for all $(XJ, p) : \mathfrak{B} \rightarrow \mathfrak{U}$ and $(\bar{J}XJ, p) : \mathfrak{B} \rightarrow \mathfrak{B}$, respectively. We consider only $r((\bar{J}XJ, p))$ and leave the other case to the reader. (For $r((XJ, p))$ the only change is that the upper left Q -leg of p disappears.)

By definition of the horizontal composition of 2-morphisms in \mathcal{E} we have

$$p \times p_1 = \lambda^{-1} \begin{array}{c} Q \quad X \quad Q \\ \boxed{p} \\ X \quad Q \quad Q \end{array} = \lambda^{-1} \begin{array}{c} Q \quad X \quad Q \\ \boxed{p} \\ X \quad Q \quad Q \end{array}$$

Consider

$$\begin{aligned} r'((\bar{J}XJ, p)) &= \begin{array}{c} Q \quad X \\ \boxed{p} \\ X \quad Q \quad Q \end{array} = \begin{array}{c} Q \quad X \\ \boxed{p} \\ XQ \quad Q \end{array} \in \text{Hom}_{\mathcal{A}}(XQ, QX) \equiv \text{Hom}_{\mathcal{E}}(\bar{J}XQJ, \bar{J}XJ) \\ r'((\bar{J}XJ, p)) &= \begin{array}{c} Q \quad X \quad Q \\ \boxed{p} \\ X \quad Q \end{array} = \begin{array}{c} Q \quad X \quad Q \\ \boxed{p} \\ X \quad Q \end{array} \in \text{Hom}_{\mathcal{A}}(XQ, QXQ) \equiv \text{Hom}_{\mathcal{E}}(\bar{J}XJ, \bar{J}XQJ) \end{aligned}$$

Using $p \bullet p = p$ and the rules of computation in \mathcal{E} it is easy to verify that $r \bullet r' = \lambda_1 p = \lambda_1 id_{(\bar{J}XJ, p)}$ and $r' \bullet r = \lambda_1 p \times p_1 = \lambda_1 id_{(\bar{J}XQJ, p \times p_1)}$, such that $r((\bar{J}XJ, p))$ is an isomorphism from $(\bar{J}XQJ, p \times p_1)$ to $(\bar{J}XJ, p)$. We leave this computation to the reader as an exercise. That $r((\bar{J}XJ, p))$ is natural w.r.t. $(\bar{J}XJ, p)$ is easy, and the coherence law connecting the unit constraint with the tensor product becomes almost obvious since \mathcal{E} is strict except for the unit morphism under study. That $\mathbf{1}_{\mathfrak{B}}$ is a left unit is shown by a similar argument defining $l((\bar{J}XJ, p)) \in Hom_{\mathcal{E}}(\bar{J}QXJ, \bar{J}XJ) \equiv Hom_{\mathcal{A}}(QXQ, QX)$ analogously. Finally, one sees that for $(\bar{J}XJ, p) = \mathbf{1}_{\mathfrak{B}}$ the left and right unit constraints coincide: $l(\mathbf{1}_{\mathfrak{B}}) = r(\mathbf{1}_{\mathfrak{B}})$.

Duality of J and \bar{J} : We refer to [22, Section I.6] for a discussion of adjoint 1-morphisms in (non-strict) bicategories. In order to show that J, \bar{J} are two-sided dual 1-morphisms we must exhibit morphisms

$$e_J : \mathbf{1}_{\mathfrak{U}} \rightarrow J\bar{J}, \quad d_J : \bar{J}J \rightarrow \mathbf{1}_{\mathfrak{B}}, \quad \varepsilon_J : \mathbf{1}_{\mathfrak{B}} \rightarrow \bar{J}J, \quad \eta_J : J\bar{J} \rightarrow \mathbf{1}_{\mathfrak{U}}$$

satisfying the usual triangular equations. Motivated by $Q = J\bar{J}$ and by Lemma 3.4 we set

$$\begin{aligned} e_J = v &\in Hom_{\mathcal{A}}(\mathbf{1}, Q) \equiv Hom_{\mathcal{E}}(\mathbf{1}_{\mathfrak{U}}, J\bar{J}), \\ \eta_J = v' &\in Hom_{\mathcal{A}}(Q, \mathbf{1}) \equiv Hom_{\mathcal{E}}(J\bar{J}, \mathbf{1}_{\mathfrak{U}}). \end{aligned} \quad (3.11)$$

Now we observe that

$$\begin{aligned} Hom_{\mathcal{E}}(\mathbf{1}_{\mathfrak{B}}, \bar{J}J) &\subset Hom_{\mathcal{E}_0}(\bar{J}J, \bar{J}J) \equiv Hom_{\mathcal{A}}(Q, Q), \\ Hom_{\mathcal{E}}(\bar{J}J, \mathbf{1}_{\mathfrak{B}}) &\subset Hom_{\mathcal{E}_0}(\bar{J}J, \bar{J}J) \equiv Hom_{\mathcal{A}}(Q, Q). \end{aligned}$$

(By construction of \mathcal{E} , $\mathbf{1}_{\mathfrak{B}}$ is the retract of $\bar{J}J : \mathfrak{B} \rightarrow \mathfrak{B}$ corresponding to the idempotent $\lambda_1^{-1} id_Q \in End_{\mathcal{A}}(Q) \equiv End_{\mathcal{E}_0}(\bar{J}J)$.) Thus, it is reasonable to consider the following candidates for d_J and ε_J (both of which live in $Hom_{\mathcal{A}}(Q, Q)$):

$$d_J = id_Q, \quad \lambda_1^{-1} \varepsilon_J = id_Q. \quad (3.12)$$

(Whether d_J or ε_J contains the factor λ_1^{-1} is immaterial, but the normalization of the left and right unit constraints l and r depends on this choice.) With this definition we have

$$d_J \bullet \varepsilon_J = \lambda_1^{-1} w' \circ w = id_Q = \lambda_1 p_1 = \lambda_1 id_{\mathbf{1}_{\mathfrak{B}}}$$

as desired. In the verification of the triangular equations we must be aware that \mathcal{E} is only a bicategory since there are non-trivial unit constraints for $\mathbf{1}_{\mathfrak{B}}$. The computations tend to be somewhat confusing. We prove only one of the four equations, namely that

$$J \equiv \mathbf{1}_{\mathfrak{U}} J \xrightarrow{e_J \otimes id_J} J\bar{J}J \xrightarrow{id_J \otimes d_J} J\mathbf{1}_{\mathfrak{B}} \xrightarrow{r(J)} J \quad (3.13)$$

is the identity 2-morphism id_J , the computation being completely analogous in the other cases. With $id_J = v' \in Hom_{\mathcal{A}}(Q, \mathbf{1}) \equiv Hom_{\mathcal{E}}(J, J)$, (3.11), and (3.12) we

compute

$$\begin{aligned}
 e_J \times \text{id}_J &= \text{diagram} \in \text{Hom}_{\mathcal{A}}(Q, Q) \equiv \text{Hom}_{\mathcal{E}}(J, J\bar{J}J) \\
 \text{id}_J \times d_J &= \text{diagram} = \text{diagram} \in \text{Hom}_{\mathcal{A}}(Q^2, Q) \equiv \text{Hom}_{\mathcal{E}}(J\bar{J}J, J\bar{J}J)
 \end{aligned}$$

•-Composing these 2-morphisms between $\mathfrak{B} \rightarrow \mathfrak{U}$ -morphisms according to the rules of Proposition 3.8 we obtain

$$\text{diagram} = \text{diagram}$$

This is precisely the isomorphism $r'(J): J\mathbf{1}_{\mathfrak{B}} \rightarrow J$ given by $r'(J) = \text{id}_Q \in \text{Hom}_{\mathcal{A}}(Q, Q) \equiv \text{Hom}_{\mathcal{E}}(J\bar{J}J, J)$ provided by Theorem 3.11. Now $r(J), r'(J)$ are mutually inverse, which proves that (3.13) gives the unit morphism id_J . The last statement is obvious since isomorphic (almost) bicategories $\mathcal{E}_0, \tilde{\mathcal{E}}_0$ have isomorphic completions $\overline{\mathcal{E}_0^p}, \tilde{\mathcal{E}_0^p}$. \square

Remark 3.12. 1. The bicategory \mathcal{E} fails to be strict (thus a 2-category) only due to the presence of non-trivial unit constraints for $\mathbf{1}_{\mathfrak{B}}$. This defect could be repaired by adding a strict unit 1-morphism for \mathfrak{B} which is isomorphic to $(\bar{J}J, p_1)$. There will, however, be no compelling reason to do so.

2. The condition (a/b) in Theorem 3.11 was crucial for identifying the unit 1-morphism $\mathbf{1}_{\mathfrak{B}}$ as a retract of $\bar{J}J$. Furthermore, we obtained a distinguished retraction/section $\mathbf{1}_{\mathfrak{B}} \leftrightarrow \bar{J}J$. So far, our assumptions are not symmetric in that they do not imply $\mathbf{1}_{\mathfrak{U}} \prec J\bar{J}$, let alone provide a canonical retraction and section. This is achieved by the following definition.

Definition 3.13. Let \mathcal{A} be an $\text{End}(\mathbf{1})$ -linear (but not necessarily a preadditive) category. A Frobenius algebra $Q = (Q, v, v', w, w')$ in \mathcal{A} is ‘strongly separable’ iff

$$w' \circ w = \lambda_1 \text{id}_Q, \quad (3.14)$$

$$v' \circ v = \lambda_2, \quad (3.15)$$

where $\lambda_1, \lambda_2 \in \text{End}(\mathbf{1})$ are invertible. If $\lambda_1 = \lambda_2$ then Q is called normalized.

Remark 3.14. 1. The term ‘strongly separable’ will find its justification by the classical case in Proposition 6.5.

2. If $\alpha, \beta \in \text{End}(\mathbf{1})^*$ and $Q = (Q, v, v', w, w')$ is a strongly separable Frobenius algebra then clearly also $\tilde{Q} = (Q, \alpha v, \beta^{-1} v', \beta w, \alpha^{-1} w')$ is one. \tilde{Q} is isomorphic to Q (in the sense of Definition 3.3) iff $\alpha = \beta$, in which case an isomorphism is given by $s = \alpha \text{id}_Q : Q \rightarrow \tilde{Q} = Q$. Yet we consider Frobenius algebras related by this renormalization as equivalent. Note that $\tilde{\lambda}_1 \tilde{\lambda}_2 = \lambda_1 \lambda_2$, thus $v' \circ w' \circ w \circ v \in \text{End}(\mathbf{1})$ is invariant under isomorphism and renormalization.

From now on all tensor categories are assumed to be $\text{End}(\mathbf{1})$ -linear. Thus, if we state this explicitly it is only for emphasis.

Proposition 3.15. *Let \mathcal{E} be a bicategory and $J : \mathfrak{B} \rightarrow \mathfrak{A}$ a 1-morphism with two-sided dual $\bar{J} : \mathfrak{A} \rightarrow \mathfrak{B}$. Assume that the corresponding 2-morphisms $d_J, e_J, \varepsilon_J, \eta_J$ can be chosen such that $\eta_J \circ e_J$ and $d_J \circ \varepsilon_J$ are invertible in the monoids $\text{End}(\mathbf{1}_{\mathfrak{A}}), \text{End}(\mathbf{1}_{\mathfrak{B}})$, respectively. Then the functor $F = - \otimes J : \text{Hom}_{\mathcal{E}}(\mathfrak{A}, \mathfrak{A}) \rightarrow \text{Hom}_{\mathcal{E}}(\mathfrak{B}, \mathfrak{A})$ is faithful and dominant in the sense that every $X : \mathfrak{B} \rightarrow \mathfrak{A}$ is a retract of $Y \circ J$ for some $Y : \mathfrak{A} \rightarrow \mathfrak{A}$. The same holds for the other seven functors given by composition with J or \bar{J} from the left or right.*

Proof. Our conditions obviously imply that $e_X : \mathbf{1}_{\mathfrak{A}} \rightarrow J\bar{J}$ and $\varepsilon_X : \mathbf{1}_{\mathfrak{A}} \rightarrow \bar{J}J$ are retractions, viz. have left inverses. Thus, $\mathbf{1}_{\mathfrak{A}} \prec J\bar{J}$ and therefore $X \prec X \circ (\bar{J} \circ J) \cong (X \circ \bar{J}) \circ J$ for any $X : \mathfrak{B} \rightarrow \mathfrak{A}$. Since $X \circ \bar{J}$ is a $\mathfrak{A} - \mathfrak{A}$ morphism this implies the dominance of F . Faithfulness can be proved using [44, Theorem IV.3.1], but we prefer to give a direct argument. Let $X, Y : \mathfrak{A} \rightarrow \mathfrak{A}$ and $s \in \text{Hom}_{\mathcal{E}}(X, Y)$. If $s \otimes \text{id}_J = 0$ then also $s \otimes \text{id}_J \otimes \text{id}_{\bar{J}} = s \otimes \text{id}_Q = 0$. Sandwiching between $\text{id}_Y \otimes v'$ and $\text{id}_X \otimes v$ gives $s \otimes (v' \circ v) = \lambda_2 s = 0$ and thus $s = 0$ by invertibility of λ_2 . \square

Corollary 3.16. *Let \mathcal{A} be $\text{End}(\mathbf{1})$ -linear and (Q, v, v', w, w') a strongly separable Frobenius algebra in \mathcal{A} . Then the bicategory \mathcal{E} defined above is such that $\eta_J \circ e_J$ and $d_J \circ \varepsilon_J$ are invertible in the monoids $\text{End}(\mathbf{1}_{\mathfrak{A}}), \text{End}(\mathbf{1}_{\mathfrak{B}})$, respectively. Conversely, if $J : \mathfrak{B} \rightarrow \mathfrak{A}$ has a two-sided dual \bar{J} such that $e_J, d_J, \varepsilon_J, \eta_J$ satisfy these conditions then the Frobenius algebras $(J\bar{J}, \dots)$ and $(\bar{J}J, \dots)$ in $\text{END}_{\mathcal{E}}(\mathfrak{A}), \text{END}_{\mathcal{E}}(\mathfrak{B})$, respectively, are strongly separable.*

Proof. By Theorem 3.11, $d_J \circ \varepsilon_J = \lambda_1 \text{id}_Q$ with $\lambda_1 \in \text{End}(\mathbf{1})^*$. On the other hand, $\eta_J \circ e_J = v' \circ v$ which is invertible in $\text{End}(\mathbf{1})$ since Q is strongly separable. The converse is obvious in view of Lemma 3.4. \square

Now we are in a position to consider the uniqueness of our bicategory \mathcal{E} .

Theorem 3.17. *Let \mathcal{A} be $\text{End}(\mathbf{1})$ -linear and (Q, v, v', w, w') a strongly separable Frobenius algebra in \mathcal{A} . Let \mathcal{E} be as constructed in Theorem 3.11 and let $\tilde{\mathcal{E}}$ be any bicategory such that:*

1. $\text{Obj } \tilde{\mathcal{E}} = \{\mathfrak{A}, \mathfrak{B}\}$.
2. Idempotent 2-morphisms in $\tilde{\mathcal{E}}$ split.

3. There is a fully faithful tensor functor $\tilde{I}: \mathcal{A} \rightarrow \text{END}_{\tilde{\mathcal{E}}}(\mathfrak{U})$ such that every object of $\text{END}_{\tilde{\mathcal{E}}}(\mathfrak{U})$ is a retract of $\tilde{I}(X)$ for some $X \in \mathcal{A}$.
4. There are mutually two-sided dual 1-morphisms $\tilde{J}: \mathfrak{B} \rightarrow \mathfrak{U}$, $\tilde{\tilde{J}}: \mathfrak{U} \rightarrow \mathfrak{B}$ and an isomorphism $\tilde{s}: I(Q) \rightarrow \tilde{J}\tilde{\tilde{J}}$ between the Frobenius algebras $I(Q, v, v', w, w')$ and $(\tilde{J}\tilde{\tilde{J}}, \tilde{e}_{\tilde{J}}, \dots)$ in $\text{End}_{\tilde{\mathcal{E}}}(\mathfrak{U})$.

Then there is an equivalence $E: \mathcal{E} \rightarrow \tilde{\mathcal{E}}$ of bicategories such that there is a monoidal natural isomorphism between the tensor functors \tilde{I} and $(E \upharpoonright \text{END}_{\mathcal{E}}(\mathfrak{U})) \circ I$.

Proof. Replacing $\mathcal{E}, \tilde{\mathcal{E}}$ by equivalent bicategories, we may assume that the functors $I: \mathcal{A} \rightarrow \text{END}_{\mathcal{E}}(\mathfrak{U})$, $\tilde{I}: \mathcal{A} \rightarrow \text{END}_{\tilde{\mathcal{E}}}(\mathfrak{U})$ are injective on the objects. Thus $\text{END}_{\mathcal{E}}(\mathfrak{U})$ and $\text{END}_{\tilde{\mathcal{E}}}(\mathfrak{U})$ contain \mathcal{A} as a full subcategory. In view of the coherence theorem for bicategories we may replace $\mathcal{E}, \tilde{\mathcal{E}}$ by equivalent strict bicategories or 2-categories and we suppress the symbols I, \tilde{I} . (As a consequence of these replacements, we will no more have the identity $I(Q) = J\tilde{J}$ but only an isomorphism $s: Q \equiv I(Q) \rightarrow J\tilde{J}$ compatible with the Frobenius algebra structures.) In view of Proposition 3.15, every $Y \in \text{Hom}_{\mathcal{E}}(\mathfrak{B}, \mathfrak{U})$ is a retract of $Y\tilde{J}J$ and therefore of XJ for $X = Y\tilde{J} \in \mathcal{A}$. Similarly, every $Z \in \text{Hom}_{\mathcal{E}}(\mathfrak{U}, \mathfrak{B})$ ($Z \in \text{Hom}_{\tilde{\mathcal{E}}}(\mathfrak{B}, \mathfrak{B})$) is a retract of $\tilde{J}X$ ($\tilde{J}XJ$) for some $X \in \mathcal{A}$, and similarly for $\tilde{\mathcal{E}}$. Let \mathcal{E}_0 be the full sub 2-category of \mathcal{E} with objects $\{\mathfrak{U}, \mathfrak{B}\}$ and 1-morphisms $X, XJ, \tilde{J}X, \tilde{J}XJ$ with $X \in \mathcal{A}$, and similarly for $\tilde{\mathcal{E}}$. Now we can define $E: \mathcal{E}_0 \rightarrow \tilde{\mathcal{E}}_0$ as the identity on objects and 1-morphisms. Composing the obvious isomorphisms $\text{Hom}_{\mathcal{E}_0}(XJ, YJ) \cong \text{Hom}_{\mathcal{A}}(XQ, Y) \cong \text{Hom}_{\tilde{\mathcal{E}}_0}(X\tilde{J}, Y\tilde{J})$, etc., provided by the duality of J, \tilde{J} and $\tilde{J}, \tilde{\tilde{J}}$ we can define the functor E_0 on the 2-morphisms. That E commutes with the horizontal and vertical compositions is obvious by the isomorphism $(J\tilde{J}, e_J, \dots) \cong (Q, v, \dots) \cong (\tilde{J}\tilde{\tilde{J}}, \tilde{e}_{\tilde{J}}, \dots)$ of Frobenius algebras. In order to obtain a (non-strict) isomorphism E of (strict) bicategories we need to define invertible 2-cells $\phi_{gf}: Eg \circ Ef \rightarrow E(g \circ f)$ satisfying the usual conditions [7]. When $\text{Ran } f = \text{Src } g = \mathfrak{U}$ we choose them to be identities and for $\text{Ran } f = \text{Src } g = \mathfrak{B}$ we use the isomorphism $\tilde{s} \circ s^{-1}: J\tilde{J} \rightarrow \tilde{J}\tilde{\tilde{J}}$. The verification of the coherence conditions is straightforward but very tedious to write down, and therefore omitted. In view of $\mathcal{E} \simeq \overline{\mathcal{E}_0}^P \cong \tilde{\mathcal{E}}0^P \simeq \tilde{\mathcal{E}}$ the isomorphism $E_0: \mathcal{E}_0 \rightarrow \tilde{\mathcal{E}}_0$ extends to an equivalence $E: \mathcal{E} \rightarrow \tilde{\mathcal{E}}$ which has all desired properties. \square

Remark 3.18. 1. The construction of the bicategory \mathcal{E} given in this section reflects the author's understanding as of 1999. More recently, it has become clear that there exists an alternative construction which can be stated quite succinctly. Namely, define a bicategory $\tilde{\mathcal{E}}$ with $\text{Obj } \tilde{\mathcal{E}} = \{\mathfrak{U}, \mathfrak{B}\}$ by positing

$$\text{HOM}_{\tilde{\mathcal{E}}}(\mathfrak{U}, \mathfrak{U}) = \mathcal{C},$$

$$\text{HOM}_{\tilde{\mathcal{E}}}(\mathfrak{U}, \mathfrak{B}) = Q - \text{mod},$$

$$\text{HOM}_{\tilde{\mathcal{E}}}(\mathfrak{B}, \mathfrak{U}) = \text{mod} - Q,$$

$$\text{HOM}_{\tilde{\mathcal{E}}}(\mathfrak{B}, \mathfrak{B}) = Q - \text{mod} - Q,$$

where $\mathcal{Q} - \text{mod}, \text{mod} - \mathcal{Q}, \mathcal{Q} - \text{mod} - \mathcal{Q}$ are the categories of left and right \mathcal{Q} -modules and of $\mathcal{Q} - \mathcal{Q}$ -bimodules in \mathcal{C} , respectively. It is very easy to prove that these categories are abelian if \mathcal{C} is abelian, in particular they ‘have subobjects’, i.e. idempotents split. The compositions of 1-morphisms are defined as suitable quotients like in ring theory. The two different constructions are related to each other as the constructions of Kleisli and Eilenberg/Moore. That they lead to equivalent bicategories is, at least morally, due to the completion w.r.t. subobjects which we apply in our construction. (Module categories automatically have subobjects.) To prove this rigorously it is, in view of Theorem 3.17, sufficient to prove that $\tilde{\mathcal{E}}$ satisfies the requirements of the latter. Inspired by [18] (which in turn was influenced by ideas of the author), the alternative definition of $\tilde{\mathcal{E}}$ was proposed also by S. Yamagami in [75], which he kindly sent me.

2. By the above remark, our category $\mathcal{B} = \text{END}_{\mathcal{E}}(\mathfrak{B})$ is equivalent to the bimodule category $\mathcal{Q} - \text{mod} - \mathcal{Q}$. Under certain technical assumptions, which are satisfied if \mathcal{A} is semisimple spherical, [59, Theorem 3.3] then implies the braided monoidal equivalence $\mathcal{Z}(\mathcal{A}) \stackrel{\otimes}{\simeq}_{br} \mathcal{Z}(\mathcal{B})$ claimed in the abstract. We hope to discuss these matters in more detail in [48].

4. Weak monoidal Morita equivalence ‘ \approx ’

Definition 4.1. A Morita context is a bicategory \mathcal{F} satisfying

1. $\text{Obj } \mathcal{F} = \{\mathfrak{U}, \mathfrak{B}\}$.
2. Idempotent 2-morphisms in \mathcal{F} split.
3. There are mutually two-sided dual 1-morphisms $J: \mathfrak{B} \rightarrow \mathfrak{U}$, $\bar{J}: \mathfrak{U} \rightarrow \mathfrak{B}$ such that the compositions $\eta_J \circ e_J \in id_{1_{\mathfrak{U}}}$ and $d_J \circ \varepsilon_J \in id_{1_{\mathfrak{B}}}$ are invertible.

Definition 4.2. Two (preadditive, k -linear) tensor categories \mathcal{A}, \mathcal{B} are weakly monoidally Morita equivalent, denoted $\mathcal{A} \approx \mathcal{B}$, iff there exists a (preadditive, k -linear) Morita context \mathcal{F} such that $\mathcal{A} \stackrel{\otimes}{\cong} \text{END}_{\mathcal{F}}(\mathfrak{U})$ and $\mathcal{B} \stackrel{\otimes}{\cong} \text{END}_{\mathcal{F}}(\mathfrak{B})$. We recall that in the non-additive case this means that there are monoidal equivalences $\tilde{\mathcal{A}}^p \stackrel{\otimes}{\simeq} \text{END}_{\mathcal{F}}(\mathfrak{U})$ and $\tilde{\mathcal{B}}^p \stackrel{\otimes}{\simeq} \text{END}_{\mathcal{F}}(\mathfrak{B})$, whereas for preadditive and k -linear categories we require $\tilde{\mathcal{A}}^{p\oplus} \stackrel{\otimes}{\simeq} \text{END}_{\mathcal{F}}(\mathfrak{U})$ and $\tilde{\mathcal{B}}^{p\oplus} \stackrel{\otimes}{\simeq} \text{END}_{\mathcal{F}}(\mathfrak{B})$. In this situation \mathcal{F} is called a Morita context for \mathcal{A}, \mathcal{B} .

Remark 4.3. 1. If in Definition 4.1 we admit $\mathfrak{U} = \mathfrak{B}$ the implication $\mathcal{A} \stackrel{\otimes}{\cong} \mathcal{B} \Rightarrow \mathcal{A} \approx \mathcal{B}$ is obvious.

2. In [56] a notion of Morita equivalence for module categories of Hopf algebras was considered, which has some similarities with the ours. Furthermore, it was shown that Hopf algebras with Morita equivalent module categories (in the sense of [56]) have the same dimension. This is reminiscent of our Proposition 5.17.

3. If the structure morphisms v, w of the Frobenius algebra are isomorphisms with $v^{-1} = v', w^{-1} = w'$ it is easy to see that the functor $X \mapsto \bar{J}XJ$ is faithful, full, essentially

surjective and monoidal. Thus $\bar{\mathcal{A}}^p \stackrel{\otimes}{\simeq} \bar{\mathcal{B}}^p$, viz. \mathcal{A} and \mathcal{B} are *strongly monoidally Morita equivalent* $\mathcal{A} \stackrel{\otimes}{\simeq} \mathcal{B}$.

4. Additional restrictions on the Morita context will be required if \mathcal{A}, \mathcal{B} are spherical or $*$ -categories.

By Lemma 3.4 a Morita context \mathcal{F} for tensor categories \mathcal{A}, \mathcal{B} provides us with strongly separable Frobenius algebras $(Q = J\bar{J}, \dots)$ and $(\hat{Q} = \bar{J}J, \dots)$ in \mathcal{A} and \mathcal{B} , respectively. Conversely, the construction of the preceding subsection provides us with a means of constructing tensor categories which are weakly Morita equivalent to a given one, together with a Morita context:

Lemma 4.4. *Let \mathcal{A} be a strict tensor category, let Q be a Frobenius algebra satisfying (3.14) and let \mathcal{E} be as constructed in the preceding subsection. Then $\mathcal{B} = \text{Hom}_{\mathcal{E}}(\mathfrak{B}, \mathfrak{B})$ is weakly Morita equivalent to \mathcal{A} , a Morita context being given by \mathcal{E} .*

Proof. Obvious by Theorems 3.11 and Definition 4.2. \square

Proposition 4.5. *Let $\mathcal{A} \approx \mathcal{B}$ with Morita context \mathcal{F} . Let (Q, v, v', w, w') be the Frobenius algebra in \mathcal{A} arising as in Lemma 3.4 and \mathcal{E} as in Theorem 3.11. Then there is an equivalence of bicategories $\mathcal{E} \simeq \mathcal{F}$. In particular, we have an equivalence $\text{Hom}_{\mathcal{E}}(\mathfrak{B}, \mathfrak{B}) \simeq \text{Hom}_{\mathcal{F}}(\mathfrak{B}, \mathfrak{B})$ of tensor categories under which the strongly separable Frobenius algebras $\hat{Q} = \bar{J}J$ of $\text{Hom}_{\mathcal{E}}(\mathfrak{B}, \mathfrak{B})$ and $\text{Hom}_{\mathcal{F}}(\mathfrak{B}, \mathfrak{B})$ go into each other.*

Proof. Obvious in view of Theorem 3.17. \square

Our terminology is justified by the following.

Proposition 4.6. *Weak monoidal Morita equivalence is an equivalence relation.*

Proof. Symmetry and reflexivity of the relation \approx are obvious. Assume $\mathcal{A} \approx \mathcal{B}$ and $\mathcal{B} \approx \mathcal{C}$ with respective Morita contexts $\mathcal{E}_1, \mathcal{E}_2$ whose objects we call $\mathfrak{U}, \mathfrak{B}_1$ and $\mathfrak{B}_2, \mathfrak{C}_2$, respectively. In order to prove transitivity we must find a Morita context for \mathcal{A} and \mathcal{C} . Since the definition of weak monoidal Morita equivalence involves only the subobject-completions, we may assume without restriction of generality that $\mathcal{A}, \mathcal{B}, \mathcal{C}$ have subobjects. We identify \mathcal{A} and $\text{End}_{\mathcal{E}_1}(\mathfrak{U})$. By definition of a Morita context we have

$$\text{End}_{\mathcal{E}_1}(\mathfrak{B}_1) \stackrel{\otimes}{\simeq} \mathcal{B} \stackrel{\otimes}{\simeq} \text{End}_{\mathcal{E}_2}(\mathfrak{B}_2),$$

and replacing $\mathcal{E}_1, \mathcal{E}_2$ by equivalent bicategories we may assume $\text{End}_{\mathcal{E}_1}(\mathfrak{B}_1) \stackrel{\otimes}{\cong} \text{End}_{\mathcal{E}_2}(\mathfrak{B}_2)$. The bicategories $\mathcal{E}_1, \mathcal{E}_2$ come with 1-morphisms $J_1 \in \text{Hom}_{\mathcal{E}_1}(\mathfrak{B}_1, \mathfrak{U})$, $J_2 \in \text{Hom}_{\mathcal{E}_2}(\mathfrak{C}_2, \mathfrak{B}_2)$ and their two-sided duals. We thus have a Frobenius algebra $Q_2 = (J_2 \bar{J}_2, v_2, v'_2, w_2, w'_2)$ in $\text{End}_{\mathcal{E}_2}(\mathfrak{B}_2)$. Using the isomorphism $\text{End}_{\mathcal{E}_1}(\mathfrak{B}_1) \stackrel{\otimes}{\cong} \text{End}_{\mathcal{E}_2}(\mathfrak{B}_2)$ we obtain the Frobenius algebra $(\hat{Q}_2, \hat{v}_2, \hat{v}'_2, \hat{w}_2, \hat{w}'_2)$ in $\text{End}_{\mathcal{E}_1}(\mathfrak{B}_1)$. We define $Q_3 =$

$J_1 \tilde{Q}_2 \bar{J}_1 \in \text{End}_{\mathcal{E}_1}(\mathcal{U})$ and claim that this is part of a Frobenius algebra Q_3 in $\mathcal{A} = \text{END}_{\mathcal{E}_1}(\mathcal{U})$ if we put

$$\begin{aligned} v_3 &= id_{J_1} \otimes \tilde{v}_2 \otimes id_{\bar{J}_1} \circ v_1, \\ v'_3 &= v'_1 \circ id_{J_1} \otimes \tilde{v}'_2 \otimes id_{\bar{J}_1}, \\ w_3 &= id_{J_1 \tilde{Q}_2} \otimes e_{J_1} \otimes id_{\tilde{Q}_2 \bar{J}_1} \circ id_{J_1} \otimes \tilde{w}_2 \otimes id_{\bar{J}_1}, \\ w'_3 &= id_{J_1} \otimes \tilde{w}_2 \otimes id_{\bar{J}_1} \circ id_{J_1 \tilde{Q}_2} \otimes \eta_{J_1} \otimes id_{\tilde{Q}_2 \bar{J}_1}. \end{aligned}$$

The verification of (3.1)–(3.5) is straightforward and therefore omitted. (This is quite similar to [44, Theorem IV.8.1] on the composition of adjoints.) Let \mathcal{E}_3 be the bicategory obtained from \mathcal{A} and Q_3 and $\{\mathcal{U}_3, \mathcal{C}_3\}$ its objects. We denote $\mathcal{C}_3 = \text{END}_{\mathcal{E}_3}(\mathcal{C}_3)$. There are functors $F_1: X \mapsto \bar{J}_2 \bar{J}_1 X J_1 J_2$ from \mathcal{A} to $\text{END}_{\mathcal{E}_2}(\mathcal{C}_1)$ (the composition of $X \mapsto \bar{J}_1 X J_1, \mathcal{A} \rightarrow \text{End}_{\mathcal{E}_1}(\mathcal{B}_1)$ and $X \mapsto \bar{J}_2 X J_2, \text{END}_{\mathcal{E}_2}(\mathcal{B}_2) \rightarrow \text{END}_{\mathcal{E}_2}(\mathcal{C}_2)$) and $F_2: X \mapsto \bar{J}_3 X J_3$ from \mathcal{A} to $\text{END}_{\mathcal{E}_3}(\mathcal{C}_3)$. In view of the definitions of $\mathcal{E}_1, \mathcal{E}_2, \mathcal{E}_3$ and of Q_3 it is clear that the images of F_1, F_2 as full subcategories of $\text{END}_{\mathcal{E}_2}(\mathcal{C}_2)$ and $\text{END}_{\mathcal{E}_3}(\mathcal{C}_3)$, respectively, are equivalent as tensor categories. Since the tensor categories $\text{END}_{\mathcal{E}_2}(\mathcal{C}_2)$ and $\text{End}_{\mathcal{E}_3}(\mathcal{C}_3)$ are equivalent to the subobject closures of the respective full subcategories they are themselves equivalent: $\text{END}_{\mathcal{E}_2}(\mathcal{C}_2) \stackrel{\otimes}{\simeq} \text{End}_{\mathcal{E}_3}(\mathcal{C}_3)$. Together with $\mathcal{C} \stackrel{\otimes}{\simeq} \text{END}_{\mathcal{E}_2}(\mathcal{C}_2)$ and $\text{END}_{\mathcal{E}_3}(\mathcal{C}_3) \approx \mathcal{A}$ this implies $\mathcal{C} \approx \mathcal{A}$. \square

Remark 4.7. 1. Comparing our notion of weak Morita equivalence with the one for rings we see that Definition 4.2 is similar (but not quite, see below) to the property of two rings R, S of admitting an invertible $A - B$ bimodule. Now, it is known [6] that this is the case iff one has either of the equivalences $R - \text{mod} \simeq S - \text{mod}$, $\text{mod} - R \simeq \text{mod} - S$. Since there is a notion of representation bicategory of a tensor category, cf. [55] and [50, Chapter 4], it is very natural to conjecture that two tensor categories are weakly monoidally Morita equivalent iff their representation categories are equivalent bicategories. (This would make the transitivity of the relation \approx obvious.) We hope to go into this question elsewhere. There is, however, one caveat, viz. in Definition 4.2 we do not require the 1-morphisms to be mutually inverse (in the sense $J\bar{J} \cong 1_{\mathcal{U}}$, $\bar{J}J \cong 1_{\mathcal{B}}$) but only to be adjoint (conjugate). Already as applied to rings this yields a weaker equivalence relation.

2. It is interesting to note that the usual Morita equivalence of (non-monoidal) categories can be expressed via the existence of a pair of mutually inverse 1-morphisms in the 2-category of small categories, distributors and their morphisms, cf. [9, Section 7.9]. One might ask whether a useful generalization is obtained by requiring only the existence of a two-sided adjoint pair of distributors.

3. Let \mathcal{A} be a $\text{End}(\mathbf{1})$ -linear tensor category. By the definitions and results of the preceding and the present section, a strongly separable Frobenius algebra Q in \mathcal{A} gives rise to a tensor category $\mathcal{B} \approx \mathcal{A}$ together with a Morita context \mathcal{E} . Conversely, a Morita context \mathcal{E} for $\mathcal{B} \approx \mathcal{A}$ gives rise to a strongly separable Frobenius algebra Q in \mathcal{A} . It is clear that this correspondence can be formalized as a one-to-one correspondence, modulo appropriate equivalence relations on both sides, between strongly separable

Frobenius algebras in \mathcal{A} and tensor categories $\mathcal{B} \approx \mathcal{A}$ together with a Morita context. Here we do not pursue this further for lack of space.

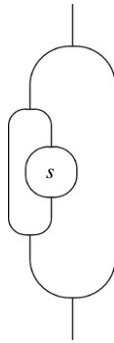
5. Linear categories

5.1. Linear categories

From now on all categories are linear over a field \mathbb{F} and all *hom*-sets are finite-dimensional over \mathbb{F} . If $\mathbf{1}$ is not simple, i.e. $\text{End}(\mathbf{1}) \not\cong \mathbb{F}$, we require strongly separable Frobenius algebras to satisfy $\lambda_1, \lambda_2 \in \mathbb{F}^*$, not just $\lambda_1, \lambda_2 \in \text{End}(\mathbf{1})^*$.

Proposition 5.1. *Let \mathcal{A} be \mathbb{F} -linear with possibly non-simple unit. Let Q be a strongly separable Frobenius algebra. Then the following holds for the bicategory \mathcal{E} of Theorem 3.11:*

- (i) J is simple iff \bar{J} is simple iff $\dim \text{Hom}_{\mathcal{A}}(Q, \mathbf{1}) = 1$ iff $\dim \text{Hom}_{\mathcal{A}}(\mathbf{1}, Q) = 1$.
- (ii) $\mathbf{1}_{\mathfrak{B}}$ is simple iff



(5.1)

is a multiple $F(s)$ of id_Q for every $s \in \text{End}(Q)$.

Furthermore, (i) implies simplicity of $\mathbf{1}_{\mathcal{A}} = \mathbf{1}_{\mathfrak{A}}$ and $\mathbf{1}_{\mathfrak{B}}$.

Proof. J is simple iff $\text{End}_{\mathcal{E}}(J) \cong \mathbb{F}$. By definition of \mathcal{E} this is the case iff $\dim \text{Hom}_{\mathcal{A}}(Q, \mathbf{1}) = 1$. Similarly, \bar{J} is simple iff $\dim \text{Hom}_{\mathcal{A}}(\mathbf{1}, Q) = 1$. The remaining equivalence in (i) follows from $\text{Hom}(J, J) \cong \text{Hom}(J\bar{J}, \mathbf{1}) \cong \text{Hom}(\bar{J}, \bar{J})$, which is a trivial consequence of duality of J, \bar{J} in \mathcal{E} . By Corollary 3.16 all functors $J \otimes -, \bar{J} \otimes -, - \otimes J, - \otimes \bar{J}$ are faithful. Thus, both $\text{End}(\mathbf{1}_{\mathfrak{A}})$ and $\text{End}(\mathbf{1}_{\mathfrak{B}})$ embed as subalgebras into $\text{End}(J)$. Thus $\text{End}(J) \cong \mathbb{F}$ implies $\text{End}(\mathbf{1}_{\mathfrak{A}}) \cong \mathbb{F}$ and $\text{End}(\mathbf{1}_{\mathfrak{B}}) \cong \mathbb{F}$. \square

Remark 5.2. 1. The implication ‘ J simple $\Leftrightarrow \bar{J}$ simple $\Rightarrow \mathbf{1}_{\mathfrak{B}}$ simple’ is reminiscent of the situation for an inclusion $B \subset A$ of von Neumann algebras, where we trivially have

$$A \cap B' = \mathbb{C}\mathbf{1} \quad \Leftrightarrow \quad B' \cap (A')' = \mathbb{C}\mathbf{1} \quad \Rightarrow \quad B' \cap B = \mathbb{C}\mathbf{1} \text{ and } A' \cap A = \mathbb{C}\mathbf{1}.$$

2. Note that we did not need a duality between $\text{Hom}(X, Y)$ and $\text{Hom}(Y, X)$ as it exists in $*$ -categories and non-degenerate spherical categories in order to conclude $\dim \text{Hom}(\mathbf{1}, Q) = 1$ iff $\dim \text{Hom}(Q, \mathbf{1}) = 1$.

Definition 5.3. Let \mathcal{A} be an \mathbb{F} -linear tensor category. A strongly separable Frobenius algebra Q in \mathcal{C} is irreducible if the equivalent conditions of (i) above hold. (By the above this is possible only if $\mathbf{1}_{\mathcal{A}}$ is simple.)

5.2. $*$ -Categories

We recall that a 1-morphism \bar{X} of a $*$ -2-category (or object in a \otimes - $*$ -category) is conjugate to X in the sense of [42] if there are $r_X : \mathbf{1} \rightarrow X\bar{X}$, $\bar{r}_X : \mathbf{1} \rightarrow \bar{X}X$ satisfying (2.4). In the generic notation of Section 2.2 this amounts to $e_X = \bar{r}_X$, $\varepsilon_X = r_X$, $d_X = r_X^*$, $\eta_X = \bar{r}_X^*$, thus $d_X = (\varepsilon_X)^*$, $\eta_X = (e_X)^*$. This implies that the Frobenius algebra $Q = (\bar{X}X, \dots)$ of Lemma 3.4 satisfies the conditions $v' = v^*$, $w' = w^*$ which we mentioned in Remark 3.2. Therefore, conditions (3.14)–(3.15) of strong separability amount to saying that v, w are (non-zero) multiples of isometries thus up to renormalization (Q, v, w) is an ‘abstract Q -system’ in the sense of [42]. In [42] it is shown that, quite remarkably, in this situation (3.5) holds automatically. Therefore, a strongly separable Frobenius algebra in a $*$ -category is the same as an algebra (Q, v, w^*) where v, w are multiples of isometries.

Definition 5.4. A Q -system is a strongly separable Frobenius algebra in a $*$ -category satisfying $v' = v^*$, $w' = w^*$. It is normalized if the Frobenius algebra (Q, v, v^*, w, w^*) is normalized, i.e. if $v^* \circ v = w^* \circ w$.

Proposition 5.5. Let \mathcal{A} be a tensor $*$ -category and $(Q, v, v' = v^*, w, w' = w^*)$ a Q -system in \mathcal{A} . Then \mathcal{E}_0 has a positive $*$ -operation $\#$ which extends the given one on \mathcal{A} . Let \mathcal{E}_* be the full sub-bicategory of \mathcal{E} whose 1-morphisms are (X, p) where $p = p \bullet p = p^\#$. Then \mathcal{E}_* is equivalent to \mathcal{E} and has a positive $*$ -operation $\#$.

Proof. Since the 2-morphisms of \mathcal{E} are given in terms of 1-morphisms in \mathcal{A} , we must denote the $*$ -operation of \mathcal{E}_0 by $\#$ in order to avoid confusion. For morphisms in $\text{END}_{\mathcal{E}}(\mathcal{U}) \equiv \mathcal{A}$ we obviously define $s^\# = s^*$. Let $X, Y \in \text{Obj } \mathcal{A}$ throughout. For $s \in \text{Hom}_{\mathcal{E}}(XJ, YJ) \equiv \text{Hom}_{\mathcal{A}}(XQ, Y)$ we define

$$s^\# = \text{id}_X \otimes r^* \circ s^* \otimes \text{id}_Q \in \text{Hom}_{\mathcal{A}}(YQ, X) \equiv \text{Hom}_{\mathcal{E}}(YJ, XJ)$$

for $s \in \text{Hom}_{\mathcal{E}}(\bar{J}X, \bar{J}Y) \equiv \text{Hom}_{\mathcal{A}}(X, QY)$ we posit

$$s^\# = \text{id}_Q \otimes s^* \circ r \otimes \text{id}_Y \in \text{Hom}_{\mathcal{A}}(Y, QX) \equiv \text{Hom}_{\mathcal{E}}(\bar{J}Y, \bar{J}X)$$

and or $s \in \text{Hom}_{\mathcal{E}}(\bar{J}XJ, \bar{J}YJ) \equiv \text{Hom}_{\mathcal{A}}(XQ, QY)$ we put

$$\begin{aligned} s^\# &= \text{id}_{QX} \otimes r^* \circ \text{id}_Q \otimes s^* \otimes \text{id}_Q \circ r \otimes \text{id}_{YQ} \in \text{Hom}_{\mathcal{A}}(YQ, QX) \\ &\equiv \text{Hom}_{\mathcal{E}}(\bar{J}YJ, \bar{J}XJ). \end{aligned}$$

(Recall that $r = w \circ v$.) Graphically the last definition looks like

$$\left(\begin{array}{c} Q \quad Y \\ | \quad | \\ \boxed{s} \\ | \quad | \\ X \quad Q \end{array} \right)^{\#} = \begin{array}{c} Q \quad X \\ | \quad | \\ \boxed{s^*} \\ | \quad | \\ Y \quad Q \end{array}$$

Antilinearity of these operations is obvious and involutivity follows from the duality equation satisfied by r, r^* . The easy verification of contravariance $((s \circ t)^{\#} = t^{\#} \circ s^{\#})$ and monoidality $((s \times t)^{\#} = s^{\#} \times t^{\#})$ are left to the reader. We limit ourselves to showing that $\#$ is positive. We consider only the case of morphisms between $\mathfrak{B} - \mathfrak{B}$ -morphisms, the others being similar. With $s \in \text{Hom}_{\mathcal{E}}(\bar{J}XJ, \bar{J}YJ) \equiv \text{Hom}_{\mathcal{A}}(XQ, QY)$ we compute

$$s^{\#} \bullet s = \begin{array}{c} Q \quad X \\ | \quad | \\ \boxed{s^*} \\ | \quad | \\ \boxed{s} \\ | \quad | \\ X \quad Q \end{array} = \begin{array}{c} Q \quad X \\ | \quad | \\ \boxed{s^*} \\ | \quad | \\ \boxed{s} \\ | \quad | \\ X \quad Q \end{array}$$

If this vanishes then (by sandwiching between $v' \otimes id_X$ and $id_X \otimes v$) also

$$\begin{array}{c} r^* \\ | \\ \boxed{s^*} \\ | \\ \boxed{s} \\ | \\ r \end{array} \quad \text{and} \quad \begin{array}{c} | \\ | \\ \boxed{s} \\ | \\ | \end{array}$$

vanish, the latter by positivity of the $*$ -operation in \mathcal{A} . Now duality implies $s = 0$, thus $\#$ is a positive $*$ -operation.

We now turn to the bicategory \mathcal{E} . Let $(X, p), (Y, q)$ be parallel 1-morphisms and $s: X \rightarrow Y$. By definition, s is a morphism $(X, p) \rightarrow (Y, q)$ iff $s = s \bullet p = q \bullet s$, which is equivalent to $s^{\#} = p^{\#} \bullet s^{\#} = s^{\#} \bullet q^{\#}$. Thus $s^{\#}: Y \rightarrow X$ is in fact in $\text{Hom}_{\mathcal{E}}((Y, q^{\#}), (X, p^{\#}))$. In the full sub-bicategory \mathcal{E}_* we have $p^{\#} = p, q^{\#} = q$, thus $s^{\#} \in \text{Hom}_{\mathcal{E}}((Y, q), (X, p))$ as it should. Finally, in a finite-dimensional $*$ -algebra (like $\text{End}_{\mathcal{E}}(X)$) every projection is

similar to an orthogonal projection. Thus, every (X, p) is isomorphic to (X, q) where q is an orthogonal projection. This proves $\mathcal{E}_* \simeq \mathcal{E}$. \square

Remark 5.6. 1. Let \mathcal{A} be a $*$ -category which has subobjects and finite-dimensional *hom*-sets. Then positivity of the $*$ -operation implies $\text{End}(X)$ to be a multi matrix algebra for every X and therefore semisimplicity of \mathcal{A} . Since \mathcal{E}_* by construction has retracts of 1-morphisms, finite-dimensional *hom*-sets and a positive $*$ -operation, we conclude that \mathcal{E}_* is semisimple (in the sense that all categories $\text{Hom}_{\mathcal{E}}(?, !)$ are semisimple).

2. As explained in Section 2.4, the notion of (two-sided) duals in $*$ -categories is local in that it does not necessitate a conjugation *map* (or functor) $X \rightarrow \bar{X}$ together with *chosen* morphisms $\mathbf{1} \rightarrow X \otimes \bar{X}$, etc. If $X \in \mathcal{A}$ has a conjugate \bar{X} it is easy to see that $XJ: \mathfrak{B} \rightarrow \mathfrak{U}$ has \bar{X} as conjugate, etc. Thus, if \mathcal{A} has conjugates for all objects then \mathcal{E}_* has conjugates for all 1-morphisms. Therefore, the above construction of a $*$ -structure on \mathcal{E} completes the discussion of $*$ -categories.

3. A self-conjugate object X in a $*$ -category is called *real* (or *orthogonal*) if there exists a solution (X, r_X, \bar{r}_X) of the conjugate equations where $r_X = \bar{r}_X$ and *pseudo-real* (or *symplectic*) if we can put $r_X = -\bar{r}_X$. (Every simple self-conjugate object is either real or pseudo-real, cf. [42].) As already observed in [42] the object of a ‘Q-system’ is real since (Q, r, \bar{r}) with $\bar{r} = r \circ w \circ v$ is a solution of the conjugate equations. By minimality of the intrinsic dimension $d(Q)$, this solution of the conjugate equations is standard iff $r^* \circ r = v^* \circ w^* \circ w \circ v$ equals $d(Q)id_1$. This is automatic when (Q, v, w) is irreducible, i.e. $\dim \text{Hom}(\mathbf{1}, Q) = 1$, as was shown in [42] using the construction of a subfactor from a Q-system. Our construction of the bicategory \mathcal{E} allows to give a simple purely categorical argument. If the Frobenius algebra Q is irreducible then Proposition 5.1 implies irreducibility of $J: \mathfrak{B} \rightarrow \mathfrak{U}$ and $\bar{J}: \mathfrak{B} \rightarrow \mathfrak{U}$ in \mathcal{E} . Then $\text{Hom}_{\mathcal{E}}(\mathbf{1}_{\mathfrak{U}}, J\bar{J})$ and $\text{Hom}_{\mathcal{E}}(\mathbf{1}_{\mathfrak{B}}, \bar{J}J)$ are one dimensional, which implies $v^* \circ v = d(J)id_1$ and $w^* \circ w = d(J)id_Q$. Thus $v^* \circ w^* \circ w \circ v = d(J)^2 id_1 = d(Q)id_1$ and (Q, r, \bar{r}) is standard.

5.3. Spherical categories

In analogy to Frobenius algebras in $*$ -categories, where we required the compatibility condition $v' = v^*, w' = w^*$, we need a compatibility of (Q, v, v', w, w') with the spherical structure of \mathcal{A} . Let Q be a 1-morphism in a strict spherical 2-category \mathcal{F} . Then $Q = J\bar{J}$ is strictly selfdual: $\bar{Q} = \bar{J}\bar{J} = J\bar{J} = Q$. If we consider the Frobenius algebra obtained from Lemma 3.4 with $e_J = \varepsilon(J)$, $\varepsilon_J = \varepsilon(\bar{J})$, $d_J = \bar{\varepsilon}(\bar{J})$, $\eta_J = \bar{\varepsilon}(J)$ then obviously

$$w \circ v = id_J \otimes \varepsilon(\bar{J}) \otimes id_{\bar{J}} \circ \varepsilon(J) = \varepsilon(J\bar{J}) = \varepsilon(Q)$$

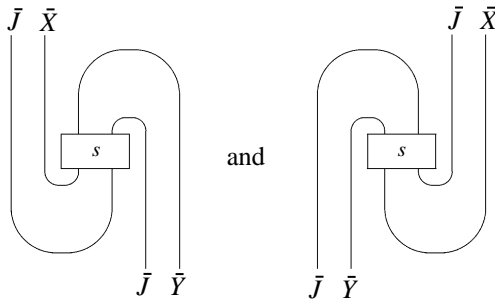
and $v' \circ w' = \bar{\varepsilon}(Q)$. Conversely, we have the following.

Lemma 5.7. *Let \mathcal{A} be a strict pivotal \mathbb{F} -linear tensor category. Let (Q, \dots) be a strongly separable Frobenius algebra such that $\bar{Q} = Q$ and $w \circ v = \varepsilon(Q)$, $v' \circ w' = \bar{\varepsilon}(Q)$. Then the bicategory \mathcal{E} of Theorem 3.11 has a strict pivotal structure which restricts to the one of \mathcal{A} .*

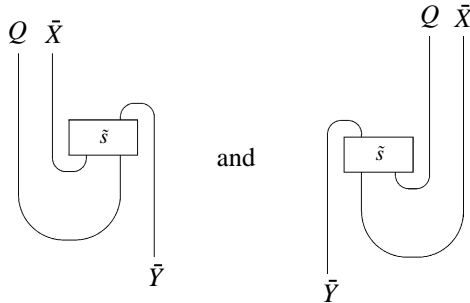
Proof. We extend the conjugation map to \mathcal{E}_0 as follows:

$$\overline{XJ} := \bar{J}\bar{X}, \quad \overline{\bar{J}X} := \bar{X}J, \quad \overline{\bar{J}XJ} := \bar{J}\bar{X}J.$$

Thus $\overline{\overline{XJ}} = \bar{X}\bar{J} = XJ$, etc., and conditions (2.1) are obvious consequences of those for \mathcal{A} . Using the notations of Theorem 3.11 we define $\varepsilon(J) = e_J$, $\varepsilon(\bar{J}) = \varepsilon_J$, $\bar{\varepsilon}(J) = \eta_J$, $\bar{\varepsilon}(\bar{J}) = d_J$. Since we know $\varepsilon(X)$ for $X \in \mathcal{A}$, condition (2) of Definition 2.6 enforces $\varepsilon(XJ) = id_X \otimes \varepsilon(J) \otimes id_{\bar{X}} \circ \varepsilon(X)$ and analogously for $\varepsilon(\bar{J}X)$, $\varepsilon(\bar{J}XJ)$. With this definition conditions (i) and (ii) of Definition 2.6 are clearly satisfied. (The conditions on $w \circ v, v' \circ w'$ are necessary and sufficient for $\varepsilon, \bar{\varepsilon}$ being well defined for all 1-morphisms since they guarantee $\varepsilon(J\bar{J}) = \varepsilon(Q)$.) It remains to verify (iii). Let, e.g. $s \in Hom_{\mathcal{E}}(XJ, YJ)$, represented by $\tilde{s} \in Hom_{\mathcal{A}}(XQ, Y)$. Now an easy computation shows that



are represented in \mathcal{A} by



respectively. These two expressions coincide since \mathcal{A} is pivotal, showing that the dual of the 2-morphism $s \in Hom_{\mathcal{E}_0}(XJ, YJ) \equiv Hom_{\mathcal{A}}(XQ, Y)$ is precisely given by $\bar{s} \in Hom_{\mathcal{A}}(\bar{Y}, Q\bar{X}) \equiv Hom_{\mathcal{E}_0}(\bar{J}Y, \bar{J}X)$, where \bar{s} is computed in \mathcal{A} . The same holds for the other types of 2-morphisms. The completion w.r.t. subobjects in a spherical tensor category behaves nicely w.r.t. duality $((\bar{X}, \bar{p}) := (\bar{X}, \bar{p}))$, and the same holds for 1-morphisms in a 2-category. Thus \mathcal{E} has strict duals and is pivotal. \square

As explained in Section 2.3, the notion of sphericity of a (non-monoidal) 2-category seems to make sense only if all identity 1-morphisms are simple, which is why we need more stringent conditions in the following.

Proposition 5.8. *Let \mathcal{A} be a strict pivotal \mathbb{F} -linear tensor category with simple unit. Let (Q, \dots) be a strongly separable Frobenius algebra such that*

- (i) $\bar{Q} = Q$.
- (ii) $w \circ v = \varepsilon(Q)$ and $v' \circ w' = \bar{\varepsilon}(Q)$.
- (iii) *The Frobenius algebra satisfies the equivalent conditions (ii) of Proposition 5.1.*

Then \mathcal{E} is spherical iff \mathcal{A} is spherical and Q is normalized. If furthermore the trace on \mathcal{A} is non-degenerate then also its natural extension to \mathcal{E} is non-degenerate.

Proof. By the preceding result and Proposition 5.1, \mathcal{E} is a bicategory with strict pivotal structure and simple $\mathbf{1}_{\mathfrak{A}}, \mathbf{1}_{\mathfrak{B}}$. Since \mathcal{A} sits in \mathcal{E} as the corner $END_{\mathcal{E}}(\mathfrak{A})$, sphericity of \mathcal{A} is clearly necessary for \mathcal{E} to be spherical. Condition (ii) implies

$$d(Q) = \bar{\varepsilon}(Q) \circ \varepsilon(Q) = v' \circ w' \circ w \circ v = \lambda_1 \lambda_2.$$

Next we observe that the $F(s)$ in Proposition 5.1 is given by

$$F(s) = \frac{\lambda_1}{\lambda_2} tr_Q(s), \quad (5.2)$$

where $tr_Q(s)$ is the trace in \mathcal{A} of $s \in End(Q)$. (Using (ii), sphericity of \mathcal{A} and $w' \circ w = \lambda_1 id_Q$ the trace (in $End_{\mathcal{A}}(Q)$) of (5.1) is seen to equal $\lambda_1^2 tr_Q(s)$. Since this trace is also equal to $F(s)tr_Q(id_Q) = F(s)d(Q) = F(s)\lambda_1\lambda_2$, the claimed equation follows.)

In order to check sphericity of \mathcal{E} we need to consider the traces on $End_{\mathcal{E}}(X)$ where X is a $\mathfrak{A} - \mathfrak{B}$ -, $\mathfrak{B} - \mathfrak{A}$ - or $\mathfrak{B} - \mathfrak{B}$ -morphism. Let $s \in End_{\mathcal{E}}(XJ, XJ)$ represented by $\tilde{s} \in Hom_{\mathcal{A}}(XQ, X)$. In view of our definition of \mathcal{E} we have

$$\begin{array}{c} X \\ | \\ \boxed{\tilde{s}} \\ | \quad | \\ X \quad Q \end{array} = \begin{array}{c} X \bar{\varepsilon}(J) \\ | \\ \boxed{s} \\ | \quad | \quad | \\ X \quad J \quad \bar{J} \end{array}$$

and with $\varepsilon(J) = e_J = v$ we have

$$tr_L(s) = \begin{array}{c} \bar{\varepsilon}(X) \\ | \\ \boxed{s} \\ | \quad | \quad | \\ \varepsilon(X) \quad \bar{J} \quad \bar{X} \end{array} = \begin{array}{c} \bar{\varepsilon}(X) \\ | \\ \boxed{\tilde{s}} \\ | \quad | \\ \varepsilon(X) \quad \bar{X} \end{array}$$

which expresses $tr_L(s)$ in terms of a formula in \mathcal{A} . The computation of

$$tr_R(s) = \begin{array}{c} \bar{\varepsilon}(\bar{J}) \\ \text{---} \bar{\varepsilon}(\bar{X}) \text{---} \\ \text{---} s \text{---} \\ \varepsilon(\bar{J}) \end{array}$$

is slightly more involved. In order to simplify matters we pretend that $X = \mathbf{1}$ which eliminates the inner trace over X . Thus

$$tr_R(s) = \begin{array}{c} \bar{\varepsilon}(\bar{J}) \\ \text{---} s \text{---} \\ \varepsilon(\bar{J}) \end{array} = \begin{array}{c} \bar{\varepsilon}(\bar{J}) \\ \text{---} \tilde{s} \text{---} \\ \varepsilon(\bar{J}) \quad \varepsilon(\bar{J}) \end{array}$$

Now, $tr_R(s) \in End_{\mathcal{E}}(\mathbf{1}_{\mathfrak{B}}) \subset End_{\mathcal{E}}(\bar{J}J) \equiv End_{\mathcal{A}}(Q)$ is represented by

$$\lambda_1^{-1} \begin{array}{c} \text{---} \\ \text{---} \text{---} \\ \text{---} v \text{---} \\ \text{---} \tilde{s} \text{---} \\ \text{---} \end{array}$$

which by assumption (iii) and (5.2) equals

$$\lambda_1^{-1} \frac{\lambda_1}{\lambda_2} tr_Q(v \circ \tilde{s}) id_Q = \lambda_1^{-1} \frac{\lambda_1}{\lambda_2} (\tilde{s} \circ v) id_Q = \frac{\lambda_1}{\lambda_2} (\tilde{s} \circ v) id_{\mathbf{1}_{\mathfrak{B}}}.$$

Thus, reintroducing the trace over X we have

$$tr_R(s) = \frac{\lambda_1}{\lambda_2} \begin{array}{c} \text{---} \tilde{s} \text{---} \\ \text{---} v \text{---} \end{array}$$

which, using the sphericity of \mathcal{A} coincides with $tr_L(s)$ iff Q is normalized. Completely analogous considerations hold for the traces on $End(\bar{J}X)$ and $End(\bar{J}XJ)$.

Assume now, that the trace on \mathcal{A} is non-degenerate and let $s \in Hom_{\mathcal{E}_0}(\bar{J}X, \bar{J}Y) \equiv Hom_{\mathcal{A}}(X, QY)$. By non-degeneracy of the trace on \mathcal{A} there is $u \in Hom_{\mathcal{A}}(QY, X)$ such that $tr_X(u \circ s) \neq 0$. Now defining

$$t = \begin{array}{c} Q \quad X \\ \downarrow \quad \downarrow \\ \boxed{u} \\ \downarrow \\ Y \end{array} \in Hom_{\mathcal{A}}(Y, QX) \equiv Hom_{\mathcal{E}_0}(\bar{J}Y, \bar{J}X),$$

one easily verifies $tr_{\bar{J}X}(t \bullet s) = tr_X(u \circ s) \neq 0$. Thus, the pairing $Hom_{\mathcal{E}_0}(\bar{J}X, \bar{J}Y) \times Hom_{\mathcal{E}_0}(\bar{J}Y, \bar{J}X) \rightarrow \mathbb{F}$ provided by the trace is non-degenerate. The other cases are verified similarly.

That the trace on \mathcal{E}_0 extends to a non-degenerate trace on the completion $\mathcal{E} = \overline{\mathcal{E}_0}^p$ follows from the following simple argument. Let tr be a non-degenerate trace, e, f idempotents and $exf \neq 0$. Then there is y such that $tr(exfy) \neq 0$. By cyclicity of the trace $tr(exfy) = tr((exf)(fye))$, thus y can be chosen such that $y = fye$. \square

Remark 5.9. 1. If \mathbb{F} is quadratically closed every strongly separable Frobenius algebra can be turned into a normalized one by renormalization. The sign of $d(J)$ depends on the choice of the renormalization, but in the $*$ -case one can achieve $d(J) = +\sqrt{d(Q)} > 0$.

2. Every simple self-dual object in a spherical category is either orthogonal or symplectic, cf. [4, p. 4018]. A simple object in a spherical category is orthogonal iff we can obtain $\bar{X} = X$ in a suitable strictification of the category. In the latter sense the object of a Frobenius algebra is orthogonal.

As mentioned in Remark 5.6, $*$ -categories are automatically semisimple and therefore semisimplicity of \mathcal{A} entails semisimplicity of \mathcal{E}_* . In order to prove an analogous result for \mathcal{A} semisimple spherical we need the following facts, which we include since we are not aware of a convenient reference.

A trace on a finite-dimensional \mathbb{F} -algebra A is a linear map $A \rightarrow \mathbb{F}$ such that $tr(ab) = tr(ba)$. It is non-degenerate if for every $a \neq 0$ there is b such that $tr(ab) \neq 0$.

Lemma 5.10. *Let A be a finite-dimensional \mathbb{F} -algebra and $tr: A \rightarrow \mathbb{F}$ a non-degenerate trace. If tr vanishes on nilpotent elements then A is semisimple. Conversely, every trace (not necessarily non-degenerate) on a semisimple algebra vanishes on nilpotent elements.*

Proof. Let R be the radical and $0 \neq x \in R$. By non-degeneracy there is $y \in A$ such that $tr(xy) \neq 0$. On the other hand $xy \in R$ and $tr(xy) = 0$ since R is nilpotent. Thus $R = \{0\}$. As to the second statement, observe that every trace on a matrix algebra coincides up to a normalization with the usual trace. (Thus $tr(e_{i,j}) = \alpha \delta_{i,j}$ with $\alpha \in \mathbb{F}$.) The latter vanishes for nilpotent matrices. The trace of a multi matrix algebra is just

a linear combination of such matrix traces on the simple subalgebras, and for general semisimple algebras the result follows by passing to an algebraic closure of \mathbb{F} . \square

Proposition 5.11. *Let \mathbb{F} be algebraically closed and \mathcal{A} strict spherical \mathbb{F} -linear and semi-simple. Let Q as in Proposition 5.8, including $\lambda_1 = \lambda_2$. Then \mathcal{E} is spherical and semisimple.*

Proof. By construction, idempotent 2-morphisms in \mathcal{E} split. Since \mathcal{A} is semisimple, thus has direct sums, the same holds for the 1-morphisms in \mathcal{E} . In order to prove that \mathcal{E} it remains to show that $End_{\mathcal{E}}(X)$ is a multi matrix algebra for every 1-morphism X .

We begin by proving that $End_{\mathcal{E}_0}(XJ) = End_{\mathcal{E}}(XJ)$ is semisimple. By Proposition 5.7 the trace tr_{XJ} on $End_{\mathcal{E}_0}(XJ)$ is non-degenerate. By Corollary 3.16 the algebra homomorphism

$$- \otimes id_{\bar{J}} : End_{\mathcal{E}_0}(XJ) \rightarrow End_{\mathcal{E}_0}(XJ\bar{J}) = End_{\mathcal{A}}(XQ)$$

is injective, such that we can consider $End_{\mathcal{E}_0}(XJ)$ as a subalgebra of $End_{\mathcal{E}_0}(XJ\bar{J})$. Furthermore, by sphericity of \mathcal{E} we have for $s \in End_{\mathcal{E}_0}(XJ)$

$$tr_{XJ\bar{J}}(s \otimes id_{\bar{J}}) = d(J) tr_{XJ}(s).$$

If s is nilpotent then also $s \otimes id_{\bar{J}} \in End_{\mathcal{E}_0}(XQ)$ is nilpotent. Thus, $tr_{XQ}(s \otimes id_{\bar{J}}) = 0$ by Lemma 5.10 and thus $tr_{XJ}(s) = 0$ since $d(J) \neq 0$. Therefore, $End_{\mathcal{E}_0}(XJ)$ is semisimple by Lemma 5.10 and a multi matrix algebra by algebraic closedness of \mathbb{F} . If A is a matrix algebra and $p = p^2 \in A$ then also $pAp \subset A$ is a matrix algebra. Thus also the endomorphism algebras $End_{\mathcal{E}}((XJ, p))$ in the completion $\mathcal{E} = \mathcal{E}_0^p$ are multi matrix algebras. Perfectly similar arguments apply to $End_{\mathcal{E}}((\bar{J}X, p))$ and $End_{\mathcal{E}}((\bar{J}XJ, p))$ for all X . \square

Conditions (i) and (ii) in Proposition 5.8 on the Frobenius algebra are fairly rigid and probably not satisfied in many applications. Furthermore, the above results should be generalized to the situation where neither the tensor product nor the duality of \mathcal{A} are strict. In the following result we limit ourselves to the degree of generality which will be needed for the application in [47]. It is fairly obvious that also the strictness conditions on \mathcal{A} can be dropped by inserting the appropriate isomorphisms wherever needed.

Theorem 5.12. *Let \mathbb{F} be algebraically closed and \mathcal{A} be \mathbb{F} -linear, strict monoidal, strict spherical and semisimple. Let $Q = (Q, v, v', w, w')$ be a normalized strongly separable Frobenius algebra in \mathcal{A} satisfying condition (ii) of Proposition 5.1. Then the bicategory \mathcal{E} of Theorem 3.11 has simple \mathfrak{B} -unit $\mathbf{1}_{\mathfrak{B}}$, is semisimple and has a (non-strict) spherical structure extending that of $\mathcal{A} \cong_{\otimes} END_{\mathcal{E}}(\mathcal{A})$.*

Proof. The \mathbb{F} -linear bicategory \mathcal{E} is defined as in Theorem 3.11. We define a conjugation map on the 1-morphisms as in Lemma 5.7. Thus we still have $\bar{\bar{X}} = X$ for all 1-morphisms. By Lemma 3.6, Q is self-dual, and since duals are unique up to isomorphism the conjugation map $X \mapsto \bar{X}$ of \mathcal{A} satisfies $\bar{\bar{Q}} \cong Q$. In fact, there is a unique

isomorphism $s: Q \rightarrow \bar{Q}$ such that

$$id_Q \otimes s \circ r = id_Q \otimes s \circ w \circ v = \varepsilon(Q): \mathbf{1} \rightarrow Q \otimes \bar{Q}. \quad (5.3)$$

This is seen as follows. If $s: Q \rightarrow \bar{Q}$ satisfies (5.3) then

On the other hand, it is equally easy to see that $s: Q \rightarrow \bar{Q}$ as defined by the second half of this equation satisfies (5.3). In view of $Q = J \circ \bar{J}$ (which is true by construction of \mathcal{E}) we have $\overline{XJ} \bar{J} Y = \overline{XQY} = \bar{Y} \bar{Q} \bar{X}$. This coincides with $\overline{JY} \circ \overline{XJ} = \bar{Y} J \bar{J} \bar{X} = \bar{Y} Q \bar{X}$ only if $Q = \bar{Q}$, which we do not assume. In any case there is an isomorphism

$$\gamma_{XJ}, \bar{J} Y = id_{\bar{Y}} \otimes s \otimes id_{\bar{X}}: \overline{JY} \circ \overline{XJ} \rightarrow \overline{XJ} \circ \overline{JY}$$

and similar for all other pairs of composable 1-morphisms. This makes \mathcal{E}_0 and \mathcal{E} bicategories with dual data in the sense of an obvious generalization of [5]. The definition of $\varepsilon(J), \bar{\varepsilon}(J), \varepsilon(\bar{J}), \bar{\varepsilon}(\bar{J})$ and therefore of ε and $\bar{\varepsilon}$ for all 1-morphisms is as in Proposition 5.8. Yet the verification of the conditions in Definition 2.6 is slightly more involved since we must insert appropriate isomorphisms in the lower lines of the commutative diagrams in condition (2). To illustrate this we consider the diagram

$$\begin{array}{ccc} \mathbf{1}_{\mathcal{A}} & \xrightarrow{\varepsilon(XJ)} & XJ \otimes \bar{XJ} \\ \varepsilon(XJ \otimes \bar{J}Y) \downarrow & & \downarrow id_{XJ} \otimes \varepsilon(\bar{J}Y) \otimes id_{\bar{XJ}} \\ XJ \otimes \bar{J}Y \otimes \overline{XJ \otimes \bar{J}Y} & \xrightarrow{id_{XJ \otimes \bar{J}Y} \otimes \gamma_{XJ, \bar{J}Y}^{-1}} & XJ \otimes \bar{J}Y \otimes \bar{J}Y \otimes \bar{XJ} \end{array}$$

In terms of the category \mathcal{A} , where all this ultimately takes place, this is

$$\begin{array}{ccccc} 1 & \xrightarrow{\varepsilon(X)} & X\bar{X} & \xrightarrow{id_X \otimes v \otimes id_{\bar{X}}} & XQ\bar{X} \\ & & & & \downarrow id_X \otimes w \otimes id_{\bar{X}} \\ & & & & XQQ\bar{X} \\ & & & & \downarrow id_{XQ} \otimes \varepsilon(Y) \otimes id_{Q\bar{X}} \\ \varepsilon(XQY) \downarrow & & & & XQY\bar{Y}Q\bar{X} \\ & & & & \downarrow id_{XQY\bar{Y}} \otimes s^{-1} \otimes id_{\bar{X}} \\ & & & & XQY\bar{Y}Q\bar{X} \end{array}$$

which commutes in view of (5.3) and the assumption that \mathcal{A} is strict pivotal. The other conditions on $\varepsilon, \bar{\varepsilon}$ are verified similarly, concluding that \mathcal{E} is a spherical bicategory. The details are omitted. The proof of semisimplicity is unchanged. \square

We conclude our discussion of spherical categories by showing that a Frobenius algebra in such a category is determined by almost as little data as in the case of $*$ -categories.

Proposition 5.13. *Let \mathcal{A} be a non-degenerate spherical category with simple unit. Let (Q, v, w') be an algebra in \mathcal{A} such that*

- (i) $\dim \text{Hom}(\mathbf{1}, Q) = 1$.
- (ii) *There is an isomorphism $s: Q \rightarrow \bar{Q}$ such that*

$$\bar{\varepsilon}(Q) \circ id_Q \otimes s = \bar{\varepsilon}(\bar{Q}) \circ s \otimes id_Q = v' \circ w' \quad (5.4)$$

with some non-zero $v': Q \rightarrow \mathbf{1}$.

Then there is also $w: Q \rightarrow Q^2$ such that (Q, v, v', w, w') is a strongly separable Frobenius algebra.

Proof. We first remark that by (i) a non-zero $v': Q \rightarrow \mathbf{1}$ exists and is unique up to a scalar. If there is a $s: Q \rightarrow \bar{Q}$ satisfying (5.4) with some v' then this obviously is the case for every choice of v' . Since $\text{Hom}(\mathbf{1}, Q), \text{Hom}(Q, \mathbf{1})$ are one dimensional and in duality we have $v' \circ v = \lambda_2 id_{\mathbf{1}}$ with $\lambda_2 \neq 0$. We write $r' = v' \circ w': Q^2 \rightarrow \mathbf{1}$. Using the fact that v is the unit for the multiplication w' we find

$$r' \circ id_Q \otimes v = v' \circ w' \circ id_Q \otimes v = v'. \quad (5.5)$$

Using (5.4), the duality equations for $\varepsilon, \bar{\varepsilon}$ and property (3) in Definition 2.6 one easily shows $id_Q \otimes s^{-1} \circ \varepsilon(Q) = s^{-1} \otimes id_Q \circ \varepsilon(\bar{Q})$. We take this as the definition of $r: \mathbf{1} \rightarrow Q^2$. One readily verifies that r, r' satisfy the triangular equations. Using the latter and (5.5) we compute

$$v = id_Q \otimes r' \circ r \otimes id_Q \circ v = id_Q \otimes v' \circ r \quad (5.6)$$

and similarly $v' \otimes id_Q \circ r = v$. In the following computation the first and last equalities hold by definition of r, r' and the middle by sphericity, viz. property (3) in Definition 2.6.

This defines a comultiplication $w: Q \rightarrow Q^2$ whose coassociativity is obvious. Together with $r = id_Q \otimes s^{-1} \circ \varepsilon(Q) = s^{-1} \otimes id_Q \circ \varepsilon(\tilde{Q})$ and (5.5) one shows $w \circ v = r$, and (5.5) implies $v' \otimes id_Q \circ w = id_Q \otimes v' \circ w = id_Q$, thus (Q, v', w) is a comonoid. Furthermore,

$$\begin{array}{c} \text{Cup with } w \end{array} \equiv \begin{array}{c} \text{Diagram with } r', w', r \text{ and } Q, Q \text{ labels} \end{array} = \begin{array}{c} \text{Diagram with } r', w', r \text{ and } Q, Q \text{ labels} \end{array} = \begin{array}{c} \text{Diagram with } w', r \end{array}, \quad (5.7)$$

where we have used $r' = v' \circ w'$ and associativity of the multiplication. Similarly,

$$\begin{array}{c} \text{Cup with } w \end{array} = \begin{array}{c} \text{Diagram with } w', r \end{array}, \quad \begin{array}{c} \text{Cap with } w' \end{array} = \begin{array}{c} \text{Diagram with } r', w \end{array} = \begin{array}{c} \text{Diagram with } r', w \end{array}. \quad (5.8)$$

Therefore,

$$\begin{array}{c} \text{Diagram with } w, w' \end{array} = \begin{array}{c} \text{Diagram with } w, w' \end{array} = \begin{array}{c} \text{Diagram with } w, w' \end{array} = \begin{array}{c} \text{Diagram with } w, w' \end{array},$$

and the other part of condition (3.5) is proved analogously. Thus (Q, v, v', w, w') is a Frobenius algebra. Using the relations proved so far we compute further

$$\begin{array}{c} \text{Diagram with } w, w' \end{array} = \begin{array}{c} \text{Diagram with } w, w' \end{array} = \begin{array}{c} \text{Diagram with } w, w' \end{array} = \begin{array}{c} \text{Diagram with } w, w' \end{array} = \begin{array}{c} \text{Diagram with } w, w' \end{array}$$

Now, by assumption (i), $w' \circ r \in \text{Hom}(\mathbf{1}, Q)$ satisfies $w' \circ r = \lambda_1 v$, implying $w' \circ w = \lambda_1 id_Q$. Furthermore, $v' \circ w' \circ r = r' \circ r = \tilde{\varepsilon}(Q) \circ \varepsilon(Q) = d(Q)$, thus $\lambda_1 \lambda_2 = d(Q) \neq 0$. Therefore (Q, v, v', w, w') is a strongly separable Frobenius algebra in \mathcal{A} and we are done. \square

Remark 5.14. 1. This result is in perfect accord with the classical definition according to which an algebra over a field \mathbb{F} is Frobenius iff it is isomorphic to \hat{A} as a left

(equivalently, right) A -module. Further support for our terminology will be supplied Section 6.1.

2. Instead of the existence of $s: Q \rightarrow \bar{Q}$ one might assume the existence of $r: \mathbf{1} \rightarrow Q^2$ satisfying the duality equations together with $r' = v' \circ w'$. Unfortunately, this approach meets a problem. One easily shows the existence of isomorphisms $s_1, s_2: Q \rightarrow \bar{Q}$ such that $id_Q \otimes s_1 \circ r = \varepsilon(Q)$ and $s_2 \otimes id_Q \circ r = \varepsilon(\bar{Q})$. Yet, it is unclear whether $s_1 = s_2$ as required. This condition can in fact be shown if \mathcal{C} is braided, $r' \circ c(Q, Q) = r'$ and $\theta(Q) = id$. Here the twist θ is defined using the spherical structure [77]. These are precisely the defining properties of a ‘rigid algebra’ in the sense of [33]. In view of the preceding remark, we find the terminology ‘Frobenius algebra’ more appropriate.

5.4. More on weak monoidal Morita equivalence

In order to maintain the correspondence between strongly separable Frobenius algebras in \mathcal{A} and tensor categories $\mathcal{B} \approx \mathcal{A}$ it is clear that we need the following.

Definition 5.15. $*$ -categories \mathcal{A}, \mathcal{B} are Morita equivalent if there is a Morita context \mathcal{F} which is a $*$ -bicategory such that the equivalences $\overline{\mathcal{A}}^{p \oplus} \otimes \overline{END_{\mathcal{F}}(\mathfrak{A})}, \overline{\mathcal{B}}^{p \oplus} \otimes \overline{END_{\mathcal{F}}(\mathfrak{B})}$, are equivalences of $*$ -categories. If \mathcal{A}, \mathcal{B} are spherical categories they are weakly monoidally Morita equivalent if there is a Morita \mathcal{F} which is spherical such that the above equivalences are equivalences of spherical categories.

We summarize our findings on $*$ - and spherical categories.

Theorem 5.16. *If \mathcal{A} is a $*$ -category and (Q, v, v^*, w, w^*) is a strongly separable Frobenius algebra then \mathcal{E}_* is a $*$ -bicategory. If \mathcal{A} is spherical (and non-degenerate (and semisimple)) and (Q, v, v', w, w') is a strongly separable Frobenius algebra then \mathcal{E} is spherical (and non-degenerate (and semisimple)). In both cases $\mathcal{B} = END_{\mathcal{E}_*}(\mathfrak{B})$ is weakly monoidally Morita equivalent to \mathcal{A} .*

As a first application of weak monoidal Morita equivalence for spherical or $*$ -categories we prove the analog of a well-known result in subfactor theory, cf. e.g. [26,16]. The proof extends to the present setting without any change.

Proposition 5.17. *Let \mathcal{A}, \mathcal{B} be a finite-dimensional semisimple spherical tensor categories over \mathbb{F} (or $*$ -categories) with simple units. If \mathcal{A}, \mathcal{B} are weakly monoidally Morita equivalent ($\mathcal{A} \approx \mathcal{B}$) then they have the same dimension in the sense of Definition 2.5.*

Proof. Let \mathcal{E} be a Morita context for $\mathcal{A} \approx \mathcal{B}$. Let I, K be (finite) sets labeling the isomorphism classes of simple $\mathfrak{A} - \mathfrak{A}$ -morphisms and $\mathfrak{B} - \mathfrak{A}$ -morphisms, respectively, and let $\{X_i, i \in I\}, \{Y_k, k \in K\}$ be objects in the respective isomorphism classes. The integers

$$N_i^k = \dim \text{Hom}_{\mathcal{E}}(Y_k, X_i J)$$

do not depend on the chosen objects, and by duality we have

$$N_i^k = \dim \operatorname{Hom}_{\mathcal{E}}(Y_k \bar{J}, X_i).$$

Thus $X_i J \cong \bigoplus_k N_i^k Y_k$ and $Y_k \bar{J} \cong \bigoplus_i N_i^k X_i$. Using additivity and multiplicativity of the dimension function we compute

$$\begin{aligned} d(J) \sum_{i \in I} d(X_i)^2 &= \sum_{i \in I} d(X_i) d(X_i J) = \sum_{\substack{i \in I \\ k \in K}} N_i^k d(X_i) d(Y_k) \\ &= \sum_{k \in K} d(Y_k) d(Y_k \bar{J}) = d(\bar{J}) \sum_{k \in K} d(Y_k)^2. \end{aligned}$$

Since $d(J)=d(\bar{J}) \neq 0$ we conclude $\dim \operatorname{Hom}_{\mathcal{E}}(\mathfrak{A}, \mathfrak{A}) = \dim \operatorname{Hom}_{\mathcal{E}}(\mathfrak{B}, \mathfrak{A})$. Entirely analogous arguments yield $\dim \operatorname{Hom}_{\mathcal{E}}(\mathfrak{B}, \mathfrak{A}) = \dim \operatorname{Hom}_{\mathcal{E}}(\mathfrak{B}, \mathfrak{B})$ and therefore $\dim \mathcal{A} = \dim \mathcal{B}$. Of course, also $\dim \operatorname{Hom}_{\mathcal{E}}(\mathfrak{A}, \mathfrak{B})$ has the same dimension. \square

Remark 5.18. 1. Note that the categories $\operatorname{Hom}_{\mathcal{E}}(\mathfrak{A}, \mathfrak{B})$ and $\operatorname{Hom}_{\mathcal{E}}(\mathfrak{B}, \mathfrak{A})$ are not tensor categories, thus the intrinsic notion of dimension of [4,42] does not apply and a priori it does not make sense to speak of their dimensions. But every object of $\operatorname{Hom}_{\mathcal{E}}(\mathfrak{A}, \mathfrak{B})$ or $\operatorname{Hom}_{\mathcal{E}}(\mathfrak{B}, \mathfrak{A})$ is a 1-morphism in \mathcal{E} and as such has a dimension. It is this dimension which is intended in the above statement.

2. Given a linear Morita context \mathcal{E} , the common dimension of the four categories $\operatorname{Hom}_{\mathcal{E}}(\cdot, \cdot)$ is also called the dimension of \mathcal{E} .

3. A less elementary application of Morita equivalence is the fact that weakly Morita equivalent spherical categories define the same state sum invariant (in the sense of [4,20]) for 3-manifolds. The proof is sketched in Section 8.

4. Let \mathcal{A} be a tensor category and (Q, \dots) a strongly separable Frobenius algebra in \mathcal{A} . It should be obvious that the tensor category $\mathcal{B} = \operatorname{Hom}_{\mathcal{E}}(\mathfrak{B}, \mathfrak{B})$ can be defined directly, avoiding the construction of the entire bicategory \mathcal{E} . When we are interested only in \mathcal{B} we may suppress the J, \bar{J} in $(\bar{J} X J, p)$. Thus the objects of \mathcal{B} are pairs (X, p) , where $X \in \operatorname{Obj} \mathcal{A}$ and $p \in \operatorname{Hom}_{\mathcal{A}}(XQ, QX)$ satisfies $p \bullet p = p$. The morphisms are given by

$$\operatorname{Hom}_{\mathcal{B}}((X, p), (Y, q)) = \{s \in \operatorname{Hom}_{\mathcal{A}}(XQ, QY) \mid s = q \bullet s \bullet p\},$$

the tensor product of objects by $(X, p) \otimes (Y, q) = (XQY, p \times q)$, etc. (Here \bullet, \times are defined as in Proposition 3.8.) For many purposes, like the study of the categorical version [48] of ‘ α -induction’ [41,71,8], this is sufficient. But proceeding in this way the weak monoidal Morita equivalence $\mathcal{A} \approx \mathcal{B}$ become obscure and even the proof of Proposition 5.17 (which is an instance of the ‘ 2×2 -matrix trick’) seems very difficult without the Morita context \mathcal{E} .

6. Examples

In this section we will consider two classes of examples: Classical Frobenius algebras over a field, in particular Hopf algebras, and subfactors with finite index. Both

examples are essential for obtaining a deep understanding of categorical Frobenius theory. Whereas most of our discussion essentially amounts to reformulating known facts, in the final subsection we will obtain a new result, viz. the weak monoidal Morita equivalence of $H\text{-mod}$ and $\hat{H}\text{-mod}$ for certain Hopf algebras. This result relies on input from both the classical Frobenius theory and subfactor theory. (That the latter are related is not new and has been discussed, e.g. in [29].)

6.1. Frobenius and Hopf algebras over fields

Here we briefly review a beautiful recent result of Abrams [2] which implies that in the case $\mathcal{A} = \mathbb{F}\text{-Vect}$ (which we treat as strict, following common usage) our notion of Frobenius algebras is equivalent to the classical one. This justifies the terminology. Note, however, that this was not our main motivation for Definition 3.1.

Let A be a finite-dimensional (associative, with unit) algebra over a field \mathbb{F} . The dual vector space \hat{A} comes with two natural coalgebra structures

$$\langle \hat{\Delta}_1(\alpha), x \otimes y \rangle = \langle \alpha, xy \rangle, \quad \langle \hat{\Delta}_2(\alpha), x \otimes y \rangle = \langle \alpha, yx \rangle,$$

both of which have the counit $\hat{\varepsilon}(\alpha) = \langle \alpha, 1 \rangle$. Given an isomorphism $\Phi: A \rightarrow \hat{A}$ of vector spaces we can provide A with a coalgebra structure by

$$\Delta = \Phi^{-1} \otimes \Phi^{-1} \circ \hat{\Delta} \circ \Phi, \quad \varepsilon = \hat{\varepsilon} \circ \Phi,$$

where $\hat{\Delta} = \hat{\Delta}_1$ or $\hat{\Delta} = \hat{\Delta}_2$.

Whenever A admits an isomorphism $\Phi: A \rightarrow \hat{A}$ of left (equivalently right) A -modules (with the natural left or right A -module structures) A is called a Frobenius algebra. We prefer the following equivalent definition, see [29].

Definition 6.1. A finite-dimensional algebra (associative, with unit) over a field \mathbb{F} is a Frobenius algebra if it admits a linear form $\phi: A \rightarrow \mathbb{F}$ which is non-degenerate (in the sense that the bilinear form $b(x, y) = \phi(xy)$ is non-degenerate).

The linear form ϕ gives rise to two isomorphisms between A and its dual \hat{A} via

$$\Phi_1: x \mapsto \phi(x \cdot), \quad \Phi_2: x \mapsto \phi(\cdot x).$$

Clearly, $\Phi_1 = \Phi_2$ iff ϕ is a trace. By the preceding discussion we thus have four canonical ways of providing A with a coalgebra structure, depending on which combination of $\Phi_1/\Phi_2, \hat{\Delta}_1/\hat{\Delta}_2$ we use. (If ϕ is a trace these possibilities reduce to two and if A is commutative we are left with a unique one. The commutative case is discussed in [1].) In any case, the counit is given by $\varepsilon = \phi$.

Theorem 6.2 (Abrams [2]). *Let A be a Frobenius algebra with given $\phi \in \hat{A}$. Let Δ_1, Δ_2 be the coproducts defined as above using the combinations $(\hat{\Delta}_2, \Phi_1)$ and $(\hat{\Delta}_2, \Phi_2)$,*

respectively. Then $\Delta_1 = \Delta_2$ and with $\Delta = \Delta_1$ the following diagrams commute:

$$\begin{array}{ccc}
 A \otimes A & \xrightarrow{m} & A \\
 \downarrow id \otimes \Delta & & \downarrow \Delta \\
 A \otimes A \otimes A & \xrightarrow{m \otimes id} & A \otimes A
 \end{array}
 \quad
 \begin{array}{ccc}
 A \otimes A & \xrightarrow{m} & A \\
 \downarrow \Delta \otimes id & & \downarrow \Delta \\
 A \otimes A \otimes A & \xrightarrow{id \otimes m} & A \otimes A
 \end{array}
 .$$

Thus $(A, \mathbf{1}, \varepsilon, \Delta, m)$ is a Frobenius algebra in the sense of Definition 3.1. Every Frobenius algebra in $\mathbb{F}\text{-Vect}$ arises in this way.

Remark 6.3. 1. In the proof, one first shows that the first diagram commutes with $\Delta = \Delta_1$ and the second with $\Delta = \Delta_2$. Using these facts one proves $\Delta_1 = \Delta_2$, thus $(A, \mathbf{1}, \varepsilon, \Delta, m)$ is a Frobenius algebra in Vect . Our version of the converse statement differs slightly from the one in [2], but it is easily seen to be equivalent.

2. An obvious consequence of the alternative characterization of Frobenius algebras is that the dual vector space of a Frobenius algebra is again a Frobenius algebra.

3. A special case of this had been shown earlier by Quinn in the little noticed appendix of [58]. He defines an ‘ambialgebra’ (in the category Vect) as an algebra and coalgebra satisfying commutativity of the above diagrams plus symmetry conditions on $\Delta(1)$ and $\varepsilon \circ m$. He states that these ambialgebras are the same as symmetric algebras (i.e. algebras admitting a non-degenerate trace). This result is intermediate in generality between those of [1,2].

4. Let X be a finite-dimensional \mathbb{F} -vector space with dual vector space \bar{X} . Then the Frobenius algebra $(X\bar{X}, \dots)$ defined as in Lemma 3.4 is well known to be just the matrix algebra $\text{End } X$. Since there are Frobenius algebras which are not isomorphic to some $M_n(\mathbb{F})$, already the category $\mathbb{F}\text{-Vect}$ provides an example of a tensor category where the Frobenius algebras are not exhausted by those of the form $(X\bar{X}, \dots)$.

5. In view of $\dim \text{Hom}_{\text{Vect}}(\mathbf{1}, Q) = \dim H$, non-trivial Frobenius algebras in Vect are not irreducible.

In order to understand when a Frobenius algebra in $\mathbb{F}\text{-Vect}$ is strongly separable we need the more general notion of a Frobenius extension [29].

Definition 6.4. A ring extension A/S is a Frobenius extension iff there exists a Frobenius system (E, x_i, y_i) , where $E \in \text{Hom}_{S-S}(A, S)$ (i.e. $E(abc) = aE(b)c \ \forall b \in A, a, c \in S$), $|I| < \infty$ and $x_i, y_i \in A$, $i \in I$ such that

$$\sum_{i \in I} x_i E(y_i a) = a = \sum_{i \in I} E(ax_i) y_i \quad \forall a \in A. \quad (6.1)$$

We call $\sum_i x_i \otimes y_i \in A \otimes_S A$ the Frobenius element and $[A : S]_E = \sum_i x_i y_i \in Z(A)$ the E -index. A Frobenius extension A/S of \mathbb{F} -algebras is called strongly separable iff $E(1) = 1$ and $[A : S]_E \in K^* 1$.

We are interested in the case where \mathbb{F} is a field and A is finite-dimensional over \mathbb{F} . Then A/\mathbb{F} is a Frobenius extension iff A is a Frobenius algebra, cf. [29, Proposition 4.8].

In this case $E = \phi$. If $\{x_i\}$ is any basis of A then $\{y_i\}$ satisfies (6.1) iff it is the dual basis: $\phi(y_i x_j) = \delta_{ij}$.

Proposition 6.5. *A Frobenius algebra A in $\mathbb{F}\text{-Vect}$ is strongly separable in the sense of Definition 3.13 iff the Frobenius extension A/\mathbb{F} is strongly separable modulo renormalization.*

Proof. It is obvious that the morphism $v' \circ v$ is given by $c \mapsto c\phi(1)$. Now, the Frobenius property (3.5) implies $\Delta(x) = \Delta(1)(1 \otimes x) = (x \otimes 1)\Delta(1)$ and therefore $m\Delta(x) = (m\Delta(1))x = x(m\Delta(1))$. Thus, the morphism $w' \circ w = m \circ \Delta$ is given by multiplication with the central element $m\Delta(1)$. We will show that $m\Delta(1) = [A : \mathbb{F}]_\phi$. Thus, if A/\mathbb{F} is strongly separable then $v' \circ v \in \mathbb{F}^* id_1$ and $w' \circ w \in \mathbb{F}^* id_A$. The converse holds since $\phi(1) \neq 0$ allows to renormalize such that $\phi(1) = 1$.

Let $\{x_i\}$ be a basis of A and $\{y_i\}$ dual in the sense $\phi(y_i x_j) = \delta_{ij}$. Then $\sum_i x_i \otimes y_i \in A \otimes A$ is the Frobenius element. For $a, b \in A$ we compute

$$\begin{aligned} ((\Phi_1 \otimes \Phi_1)(\Delta(1)))(a \otimes b) &= (\hat{\Delta}_2 \Phi_1(1))(a \otimes b) = \Phi_1(1)(ba) = \phi(ba) \\ &= \sum_i \phi(ax_i) \phi(by_i), \end{aligned}$$

where we used (6.1). Thus, $\Delta(1)$ equals the Frobenius element and $m\Delta(1) = \sum_i x_i y_i = [A : \mathbb{F}]_\phi$. \square

Remark 6.6. In the commutative case Frobenius algebras satisfying the above equivalent conditions were called ‘superspecial’ in [58]. Note furthermore that semisimplicity of A is equivalent to the weaker condition of invertibility of $\sum_i x_i y_i$ (proven in [1] for the commutative and in [58] for the symmetric case). Thus strongly separable Frobenius algebras are semisimple.

It is well known [38] that every finite-dimensional Hopf algebra over a field \mathbb{F} is a Frobenius algebra. Our aim in the remainder of this subsection is to clarify when these Frobenius algebras are strongly separable. We recall some well-known facts. For any finite-dimensional Hopf algebra H one can prove [38] that the subspaces $I_L(H) = \{y \in H \mid xy = \varepsilon(x)y \ \forall x \in H\}$ and $I_R(H) = \{y \in H \mid yx = \varepsilon(x)y \ \forall x \in H\}$ are one dimensional satisfy $S(I_L(H)) = I_R(H)$. Furthermore, every non-zero $\varphi_L \in I_L(\hat{H})$ and $\varphi_R \in I_R(\hat{H})$ is a non-degenerate functional on H .

In view of Theorem 6.2, both φ_L and φ_R give rise to coalgebra structures $(H, \tilde{\Delta}_{L/R}, \tilde{\varepsilon}_{L/R})$ on the vector space H and therefore to Frobenius algebras $Q_{L/R} = (H, m, \eta, \tilde{\Delta}_{L/R}, \tilde{\varepsilon}_{L/R})$. (We denote the Frobenius coproduct by $\tilde{\Delta}_{L/R}$ to avoid confusion with the Hopf algebra coproduct Δ of H .)

Proposition 6.7. *Let H be a finite-dimensional Hopf algebra with non-zero right integrals $\lambda \in I_R(H), \varphi \in I_R(\hat{H})$. Then the following are equivalent:*

- (i) *The Frobenius algebra $Q_L = (H, m, \eta, \tilde{\Delta}_L, \tilde{\varepsilon}_L)$ in $\mathbb{F}\text{-Vect}$ is strongly separable in the sense of Definition 3.13.*

- (ii) $\langle \varphi, 1 \rangle \neq 0$ and $\langle \varepsilon, A \rangle \neq 0$.
 (iii) H is semisimple and cosemisimple.

Proof. (ii) \Leftrightarrow (iii). By [38], H is semisimple iff $\varepsilon(A) \neq 0$ and cosemisimple iff $\varphi(1) \neq 0$.

(i) \Leftrightarrow (ii). By [29, Proposition 6.4], a Frobenius system for H/\mathbb{F} is given by the triple $(\varphi, S^{-1}(A_{(2)}), A_{(1)})$. (This is to say that the Frobenius element is given by $\sum_i x_i \otimes y_i = (S^{-1} \otimes id)A^{\text{op}}(A)$.) But then it is obvious that $[H : \mathbb{F}]_{\varphi} = S^{-1}(A_{(2)})A_{(1)} = \varepsilon(A)1$. Thus $m \circ \tilde{A}_L = [H : \mathbb{F}]_f = \varepsilon(A)1$, and the equivalence (i) \Leftrightarrow (ii) follows from Proposition 6.5. \square

Remark 6.8. In view of $S(I_L) = I_R$ and $S(1) = 1$ we have $\langle \varphi_L, 1 \rangle \neq 0 \Leftrightarrow \langle \varphi_R, 1 \rangle \neq 0$. Thus, it does not matter where the integrals appearing in condition (ii) are left or right integrals. Similarly, in (i) we can write Q_R instead of Q_L . In the strongly separable case these substitutions are vacuous since semisimple Hopf algebras are unimodular, i.e. $I_L = I_R$.

Corollary 6.9. Let H be a finite-dimensional Hopf algebra over \mathbb{F} and Q the corresponding Frobenius algebra (in $\mathbb{F}\text{-Vect}$). Let A, φ be both either left or right integrals in H, \hat{H} , respectively. Then there is the following identity of numerical invariants of Q and H :

$$v' \circ w' \circ w \circ v = \frac{\langle \varphi, 1 \rangle \langle \varepsilon, A \rangle}{\langle \varphi, A \rangle} \in \text{End}(\mathbf{1}) \equiv \mathbb{F}. \quad (6.2)$$

Whenever this number is non-zero H is semisimple and cosemisimple and (6.2) coincides with $\dim H \cdot 1_{\mathbb{F}}$.

Proof. Eq. (6.2) follows from the above computation of $v' \circ v$ and $w' \circ w$. If (6.2) is non-zero then H is semisimple and cosemisimple, and by [14] the antipode is involutive. By [37, Theorem 2.5], (6.2) coincides with $\text{tr}(S^2)$ and therefore with $\dim H \cdot 1_{\mathbb{F}}$. \square

By the preceding result semisimple and cosemisimple Hopf algebras provide examples of strongly separable Frobenius algebras in $\mathbb{F}\text{-Vect}$. By Remark 6.3.4 they are not irreducible. We will now show how one can associate a strongly separable and irreducible Frobenius algebra with a Hopf algebra which is semisimple and cosemisimple.

6.2. Hopf algebras: Frobenius algebras in $H - \text{mod}$

In the theory of quantum groups the following result is known as ‘strong left invariance’ (for $b = 1$ or $c = 1$ it reduces to left invariance: $(id \otimes \varphi)(\Delta(b)) = \varphi(b)1$), but it is also true for all finite-dimensional Hopf algebras. We include the proof since we are not aware of a convenient reference.

Lemma 6.10. *Let H be a finite-dimensional Hopf algebra and $\varphi \in I_L(\hat{H})$. Then*

$$(id \otimes \varphi)((1 \otimes c)\Delta(b)) = (id \otimes \varphi)((S \otimes id)(\Delta(c))(1 \otimes b)) \quad \forall b, c \in H.$$

Proof. For every $x, y \in H$ we have

$$\begin{aligned} x \otimes y &= \sum_i (u_i \otimes 1)\Delta(v_i) \quad \text{with} \quad \sum_i u_i \otimes v_i = (x \otimes 1)(S \otimes id)(1) \\ &= xS(y_{(1)}) \otimes y_{(2)} \end{aligned}$$

as is verified by a trivial computation. Using this representation and left invariance of φ we have $(id \otimes \varphi)(x \otimes y) = \sum_i u_i(id \otimes \varphi)\Delta(v_i) = \sum_i u_i\varphi(v_i) = (id \otimes \varphi)(\sum_i u_i \otimes v_i)$. Applying this to $x \otimes y = (1 \otimes c)\Delta(b) = b_{(1)} \otimes cb_{(2)}$ we find

$$\begin{aligned} \sum_i u_i \otimes v_i &= b_{(1)}S((cb_{(2)})_{(1)}) \otimes (cb_{(2)})_{(2)} = b_{(1)}S(c_{(1)}b_{(2)}) \otimes c_{(2)}b_{(3)} \\ &= b_{(1)}S(b_{(2)})S(c_{(1)}) \otimes c_{(2)}b_{(3)} = S(c_{(1)}) \otimes c_{(2)}b \end{aligned}$$

and therefore,

$$\begin{aligned} (id \otimes \varphi)((1 \otimes c)\Delta(b)) &= (id \otimes \varphi)(S(c_{(1)}) \otimes c_{(2)}b) \\ &= (id \otimes \varphi)((S \otimes id)(\Delta(c))(1 \otimes b)) \end{aligned}$$

as desired. \square

Proposition 6.11. *Let H be a finite-dimensional Hopf algebra and $\varphi_L \in I_L(\hat{H})$. We write \hat{m} for the multiplication of \hat{H} and define $\mathcal{F}: H \rightarrow \hat{H}$ by $\mathcal{F}(a)(\cdot) = \varphi(\cdot a)$. Then the map $\tilde{m} = \mathcal{F}^{-1}\hat{m}(\mathcal{F} \otimes \mathcal{F}): H \otimes H \rightarrow H$ satisfies*

- (i) $\varphi(c\tilde{m}(a \otimes b)) = (\varphi \otimes \varphi)(\Delta(c)(a \otimes b)) = (\varphi \otimes \varphi)((1 \otimes c)(S^{-1} \otimes id)(\Delta(b))(a \otimes 1)) = (\varphi \otimes \varphi)((c \otimes 1)(1 \otimes S^{-1}(b))\Delta(a)) \quad \forall a, b, c \in H.$
- (ii) $\tilde{m}(\Delta(c)x) = c\tilde{m}(x) \quad \forall c \in H, x \in H \otimes H.$
- (iii) $\tilde{m}(a \otimes b) = \varphi(S^{-1}(b_{(1)})a)b_{(2)} = a_{(1)}\varphi(S^{-1}(b)a_{(2)}).$

Proof. (i) We compute

$$\begin{aligned} \varphi(c\tilde{m}(a \otimes b)) &= \langle \mathcal{F}\tilde{m}(a \otimes b), c \rangle = \langle \hat{m}(\mathcal{F}(a) \otimes \mathcal{F}(b)), c \rangle \\ &= \langle \mathcal{F}(a) \otimes \mathcal{F}(b), \Delta(c) \rangle = \underline{(\varphi \otimes \varphi)(\Delta(c)(a \otimes b))} \\ &= \varphi(S^{-1}[(id \otimes \varphi)((S \otimes id)(\Delta(c))(1 \otimes b))]a) \\ &= \varphi(S^{-1}[(id \otimes \varphi)((1 \otimes c)\Delta(b))]a) \\ &= \underline{(\varphi \otimes \varphi)((1 \otimes c)(S^{-1} \otimes id)(\Delta(b))(a \otimes 1))} \\ &= (\varphi \otimes \varphi)((1 \otimes c)\sigma((S \otimes id)(\Delta(S^{-1}(b))))(a \otimes 1)) \\ &= (\varphi \otimes \varphi)((c \otimes 1)(S \otimes id)(\Delta(S^{-1}(b)))(1 \otimes a)) \end{aligned}$$

$$\begin{aligned}
&= \varphi(c[(id \otimes \varphi)((S \otimes id)(\Delta(S^{-1}(b)))(1 \otimes a))]) \\
&= \varphi(c[(id \otimes \varphi)((1 \otimes S^{-1}(b))\Delta(a))]) \\
&= \underline{(\varphi \otimes \varphi)((c \otimes 1)(1 \otimes S^{-1}(b))\Delta(a))}.
\end{aligned}$$

In the first two lines we used the definitions of \mathcal{F} and \tilde{m} , whereas the sixth and 11th equalities follow by application of Lemma 6.10 to the expression in square brackets. The remaining identities result from trivial rearrangements using $(\varphi \otimes \varphi) = \varphi(id \otimes \varphi) = \varphi(\varphi \otimes id)$.

(ii) Let $x = a \otimes b$. Twofold use of the first equality in (i) yields

$$\varphi(d\tilde{m}(\Delta(c)x)) = (\varphi \otimes \varphi)(\Delta(d)\Delta(c)x) = (\varphi \otimes \varphi)(\Delta(dc)x) = \varphi(dc\tilde{m}(x)),$$

which holds for all $c, d \in H, x \in H \otimes H$. The claim now follows by non-degeneracy of φ .

(iii) We can rewrite (i) as

$$\begin{aligned}
\varphi(c\tilde{m}(a \otimes b)) &= \varphi(c(\varphi \otimes id)(S^{-1} \otimes id)(\Delta(b))(a \otimes 1))) \\
&= \varphi(c(id \otimes \varphi)((1 \otimes S^{-1}(b))\Delta(a))).
\end{aligned}$$

Now we appeal again to non-degeneracy of φ and rewrite in Sweedler notation. \square

Theorem 6.12. *Let H be a finite-dimensional Hopf algebra over \mathbb{F} . Let Λ, φ be left integrals in H and \tilde{H} , respectively, normalized such that $\langle \varphi, \Lambda \rangle = 1$. Let $Q \in H - \text{mod}$ be the left regular representation, viz. H acting on itself by $\pi_Q(a)b = ab$. The linear maps*

$$\begin{aligned}
v : \quad \mathbb{F} &\rightarrow Q, & c &\mapsto c\Lambda, \\
v' : \quad Q &\rightarrow \mathbb{F}, & x &\mapsto \varepsilon(x), \\
w : \quad Q &\rightarrow Q \otimes Q, & x &\mapsto \Delta(x), \\
w' : \quad Q \otimes Q &\rightarrow Q, & x \otimes y &\mapsto \tilde{m}(x \otimes y)
\end{aligned} \tag{6.3}$$

are morphisms in $H - \text{mod}$ and (Q, v, v', w, w') is an irreducible Frobenius algebra in $H - \text{mod}$. It is strongly separable iff H is semisimple and cosemisimple.

Proof. In order to show that the maps defined above are morphisms in the category $H - \text{mod}$ we must show that they intertwine the H -actions. This follows from the following diagrams, where $\pi_1 = \varepsilon$ is the tensor unit:

$$\begin{array}{ccc}
c & \xrightarrow{v} & c\Lambda \\
\pi_1(z) \downarrow & & \downarrow \pi_Q(z) \\
\varepsilon(z)c & \xrightarrow{v} & \varepsilon(z)c\Lambda = cz\Lambda.
\end{array}
\qquad
\begin{array}{ccc}
x & \xrightarrow{v'} & \varepsilon(x) \\
\pi_Q(z) \downarrow & & \downarrow \pi_1(z) \\
zx & \xrightarrow{v'} & \varepsilon(zx) = \varepsilon(x)\varepsilon(z).
\end{array}$$

$$\begin{array}{ccc}
x & \xrightarrow{w} & \Delta(x) \\
\pi_Q(z) \downarrow & & \downarrow \pi_{Q \otimes Q}(z) \\
zx & \xrightarrow{w} & \Delta(zx) = \Delta(z)\Delta(x).
\end{array}
\quad
\begin{array}{ccc}
x \otimes y & \xrightarrow{w'} & \tilde{m}(x \otimes y) \\
\pi_{Q \otimes Q}(z) \downarrow & & \downarrow \pi_Q(z) \\
\Delta(z)(x \otimes y) & \xrightarrow{w'} & \tilde{m}(\Delta(z)(x \otimes y)) = z\tilde{m}(x \otimes y)
\end{array}$$

Commutativity of the lower right diagram follows from Proposition 6.11 (ii).

The equations $v' \otimes id_Q \circ w = id_Q \otimes v' \circ w = id_Q$ and $w \otimes id_Q \circ w = id_Q \otimes w \circ w$ are obvious since (Q, w, v') coincides with the coalgebra structure of H . That $w' : Q^2 \rightarrow Q$ is associative is evident in view of $\tilde{m} = \mathcal{F}^{-1}\hat{m}(\mathcal{F} \otimes \mathcal{F})$ and associativity of \hat{m} . Furthermore, $\mathcal{F}(\Lambda)(a) = \varphi(a\Lambda) = \varepsilon(a)\varphi(\Lambda) = \varepsilon(a)$, thus $\mathcal{F}(\Lambda) = \varepsilon = 1_{\hat{H}}$ is the unit for \tilde{m} and (Q, v, w') is a monoid.

Applying ε to Proposition 6.11 (iii) we obtain $(\varepsilon\tilde{m})(a \otimes b) = \varphi(S^{-1}(b)a)$. Comparing with the formulae for \tilde{m} we find

$$\begin{aligned}
\tilde{m}(a \otimes b) &= \varphi(S^{-1}(b_{(1)})a)b_{(2)} = (\varepsilon\tilde{m})(a \otimes b_{(1)})b_{(2)} \\
&= a_{(1)}\varphi(S^{-1}(b)a_{(2)}) = a_{(1)}(\varepsilon\tilde{m})(a_{(2)} \otimes b)
\end{aligned}$$

or in diagrams

$$\begin{array}{c} \text{cup} \end{array} \xrightarrow{\varepsilon} \Delta = \begin{array}{c} \text{cup} \end{array} \xrightarrow{\tilde{m}} = \begin{array}{c} \text{cup} \end{array} \xrightarrow{\Delta} \varepsilon$$

Using the first of these equalities twice to compute

$$\begin{array}{c} \text{cup} \end{array} \xrightarrow{\Delta} \varepsilon = \begin{array}{c} \text{cup} \end{array} \xrightarrow{\varepsilon} \Delta = \begin{array}{c} \text{cup} \end{array} \xrightarrow{\Delta} \varepsilon = \begin{array}{c} \text{cup} \end{array} \xrightarrow{\varepsilon} \Delta$$

we have proven one of the Frobenius conditions (3.5) and the other one follows in the same vein.

That the Frobenius algebra Q is irreducible follows from the obvious isomorphism of vector spaces $\text{Hom}_{H\text{-mod}}(\mathbf{1}, Q) \cong I_L$ together with $\dim I_L = 1$. Finally, we compute

$$(\tilde{m}\Delta)(a) = \tilde{m}(a_{(1)} \otimes a_{(2)}) = a_{(1)}\varphi(S^{-1}(a_{(3)})a_{(2)}) = a\varphi(1).$$

Thus Q is strongly separable iff $\varepsilon(\Lambda) \neq 0$ and $\varphi(1) \neq 0$, which is the case iff H is semisimple and cosemisimple. \square

Example 6.13. Let $H = \mathbb{F}(G)$ be the algebra of \mathbb{F} -valued functions on a finite group G with the usual Hopf algebra structure. With $H = \text{span}\{\delta_g, g \in G\}$, where $\delta_g(h) = \delta_{g,h}$, the integrals $A \in H, \varphi \in \hat{H}$ are

$$A = \delta_e,$$

$$\langle \varphi, \delta_g \rangle = 1.$$

Then we find

$$\langle \mathcal{F}(\delta_g), \delta_h \rangle = \langle \varphi, \delta_g \delta_h \rangle = \delta_{g,h} \langle \varphi, \delta_g \rangle = \delta_{g,h}$$

and thus $\mathcal{F}(\delta_g) = u_g$, where $\{u_g, g \in G\}$ is the usual basis in $\hat{H} = \mathbb{F}G$. We obtain

$$\mathbf{1}_Q = \delta_e = A,$$

$$m_Q(\delta_g, \delta_h) = \delta_{gh}$$

and thus the Frobenius algebra Q associated to $H = \mathbb{F}(G)$ by our prescription coincides with the one given in [58, A.4.5]. In a similar fashion one sees that the Frobenius algebra associated with $\mathbb{F}G$ has the coalgebra structure of $\mathbb{F}G$ and the algebra structure of $\mathbb{F}(G)$ under the correspondence $u_g \leftrightarrow \delta_g$.

Remark 6.14. Our original proof of Theorem 6.12 has improved considerably as a consequence of discussions with L. Tuset. In the joint work [49] we examine to which extent the above results carry over to not necessarily finite-dimensional algebraic quantum groups. If an algebraic quantum group (A, Δ) is discrete, there exists a monoid $(\pi_L, \tilde{m}, \tilde{\eta})$ in the category $\text{Rep}(A, \Delta)$ of non-degenerate $*$ -representations. If (A, Δ) is compact (=unital) then there is a comonoid $(\pi_L, \tilde{\Delta}, \tilde{\varepsilon})$ in $\text{Rep}(A, \Delta)$. Both the monoid and comonoid structures exist only if (A, Δ) is finite-dimensional, in which case they coincide with those considered above. Therefore, in the infinite-dimensional situation one does not obtain a Frobenius algebra in $\text{Rep}(A, \Delta)$, but a ‘regularized’ version of the Frobenius relation (3.5) can still be proven.

6.3. Morita equivalence of $H - \text{mod}$ and $\hat{H} - \text{mod}$

In this subsection, \mathbb{F} is an arbitrary algebraically closed field. If H is a finite-dimensional semisimple and cosemisimple Hopf algebra over \mathbb{F} , Theorem 6.12 gives rise to a strongly separable Frobenius algebra Q in $H - \text{mod}$. Applying Theorem 3.11, we obtain a Morita context \mathcal{E} , and it is natural to ask what can be said about the tensor category $\mathcal{B} = \text{END}_{\mathcal{E}}(\mathfrak{B})$, which is Morita equivalent to $H - \text{mod}$ by construction. We may and will assume \mathcal{E} to be strict, i.e. a 2-category. The aim of this subsection is to prove the following.

Theorem 6.15. *Let H be a finite-dimensional semisimple and cosemisimple Hopf algebra over an algebraically closed field \mathbb{F} and let Q be the associated strongly separable Frobenius algebra in $H - \text{mod}$. If \mathcal{E} is as in Theorem 3.1 and $\mathcal{B} = \text{HOM}_{\mathcal{E}}(\mathfrak{B}, \mathfrak{B})$ then we have the equivalence $\mathcal{B} \stackrel{\otimes}{\simeq} \hat{H} - \text{mod}$ of spherical tensor categories.*

The theorem will be an easy consequence of the more general Theorem 6.20. In view of Definition 4.2 we obtain the remarkable.

Corollary 6.16. *Let H be a finite-dimensional semisimple and cosemisimple Hopf algebra. Then we have the weak monoidal Morita equivalence $H - \text{mod} \approx \hat{H} - \text{mod}$ of spherical tensor categories.*

We begin with a semisimple spherical \mathbb{F} -linear Morita context \mathcal{E} . (Recall that we require $\mathbf{1}_{\mathfrak{A}}$ and $\mathbf{1}_{\mathfrak{B}}$ to be simple.) We denote $A = \text{End}(J\bar{J})$, $B = \text{End}(\bar{J}J)$, $C = \text{End}(J\bar{J}J)$, and we write $\text{Tr}_A, \text{Tr}_B, \text{Tr}_C$ instead of $\text{Tr}_{J\bar{J}}, \text{Tr}_{\bar{J}J}, \text{Tr}_{J\bar{J}J}$. We define a linear map, the ‘Fourier transform’, by

$$\mathcal{F} : A \rightarrow B, \quad \begin{array}{c} \text{---} \\ | \\ \boxed{a} \\ | \\ J \quad \bar{J} \end{array} \mapsto \begin{array}{c} \text{---} \\ | \\ \boxed{a} \\ \varepsilon(\bar{J}) \quad \bar{J} \quad J \end{array} \quad \bar{\varepsilon}(\bar{J})$$

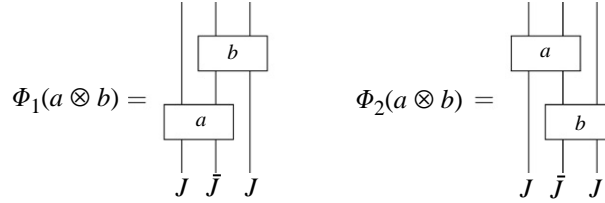
and $\hat{\mathcal{F}} : B \rightarrow A$ is defined by the same diagram with the obvious changes. The Fourier transforms are clearly invertible. Furthermore, we define ‘antipodes’ $S : A \rightarrow A, \hat{S} : B \rightarrow B$ by $S = \hat{\mathcal{F}} \circ \mathcal{F}, \hat{S} = \mathcal{F} \circ \hat{\mathcal{F}}$. The antipodes are easily seen to be antimultiplicative: $S(ab) = S(b)S(a)$ and analogously for \hat{S} . As a consequence of axiom (3) in Definition 2.6, we have $S \circ S = \text{id}_A, \hat{S} \circ \hat{S} = \text{id}_B$. Using the Fourier transforms we define ‘convolution products’ on A and B by $a \star b = \mathcal{F}^{-1}(\mathcal{F}(a)\mathcal{F}(b))$ for $a, b \in A$, and similarly for B . One easily verifies

$$a \star b = \begin{array}{c} \text{---} \\ | \\ \boxed{a} \quad \boxed{b} \\ | \quad | \\ J \quad \bar{J} \end{array}$$

The Fourier transform further allows to define a bilinear form $\langle \cdot, \cdot \rangle : A \otimes B \rightarrow \mathbb{F}$ by $\langle a, b \rangle = d(J)^{-1} \text{Tr}_A(a\mathcal{F}^{-1}(b))$. Since \mathcal{F} is bijective and Tr_A is non-degenerate, this bilinear form establishes a duality between A and B . One verifies

$$d(J) \langle a, b \rangle = \text{Tr}_B(\hat{\mathcal{F}}^{-1}(a)b) = \text{Tr}_J \begin{array}{c} \text{---} \\ | \\ \boxed{b} \\ | \\ \boxed{a} \\ | \\ J \end{array} = \text{Tr}_{\bar{J}} \begin{array}{c} \text{---} \\ | \\ \boxed{a} \\ | \\ \boxed{b} \\ | \\ \bar{J} \end{array}$$

For later use we remark that with $a, b \in A$ we have $\langle a, \mathcal{F}(b) \rangle = d(J)^{-1} \text{Tr}_A(ab) = d(J)^{-1} \text{Tr}_A(ba) = \langle b, \mathcal{F}(a) \rangle$. The duality between A and B enables us to define

Fig. 5. $\Phi_1(a \otimes b)$ and $\Phi_2(a \otimes b)$.

coproducts $\Delta : A \rightarrow A \otimes A$, $\hat{\Delta} : B \rightarrow B \otimes B$ by

$$\langle \Delta(a), x \otimes y \rangle = \langle a, xy \rangle, \quad a \in A, x, y \in B,$$

$$\langle a \otimes b, \hat{\Delta}(x) \rangle = \langle ab, x \rangle, \quad a, b \in A, x \in B.$$

Associativity of \hat{m} (m) implies coassociativity of Δ ($\hat{\Delta}$). With

$$\varepsilon(a) = \langle a, 1 \rangle, \quad \hat{\varepsilon}(x) = \langle 1, x \rangle, \quad a \in A, x \in B$$

it is clear that (A, Δ, ε) and $(B, \hat{\Delta}, \hat{\varepsilon})$ are coalgebras. We note that $\varepsilon(1) = \langle 1, \hat{1} \rangle = d(J)^{-1} \text{Tr}_J \text{id}_J = 1$, which explains the normalization of $\langle \cdot, \cdot \rangle$.

The above considerations are valid without further assumptions on the Morita context. In order to establish A, B as mutually dual Hopf algebras it remains to show that the maps $\Delta, \hat{\Delta}, \varepsilon, \hat{\varepsilon}$ are multiplicative and that the antipodes are coinverses. It is here that further assumptions are needed.

Definition 6.17. A semisimple \mathbb{F} -linear Morita context \mathcal{E} has ‘depth two’ if every simple summand of $J\bar{J}J \in \text{Hom}(\mathfrak{B}, \mathfrak{A})$ appears as a simple summand of J . \mathcal{E} is called irreducible if the distinguished 1-morphism $J : \mathfrak{B} \rightarrow \mathfrak{A}$ is simple.

If \mathcal{E} is irreducible and has depth two then $J\bar{J}J$ is a multiple of J . Here we restrict ourselves to the irreducible depth two case, which is all we need to prove Theorem 6.15. Note that we do not assume that J and \bar{J} generate \mathcal{E} , as is the case in subfactor theory. In a depth two Morita context where this is the case, every simple $\mathfrak{B} - \mathfrak{A}$ -morphism is isomorphic to J . For results on the reducible depth two case—which leads to finite quantum groupoids—see [66], where, however, not all proofs are given.

Lemma 6.18. *In addition to the above assumptions, let \mathcal{E} be irreducible of depth two. Then*

1. *The maps*

$$\Phi_1 : A \otimes B \rightarrow C, \quad a \otimes b \mapsto \text{id}_{\bar{J}} \otimes b \circ a \otimes \text{id}_J,$$

$$\Phi_2 : A \otimes B \rightarrow C, \quad a \otimes b \mapsto a \otimes \text{id}_J \circ \text{id}_{\bar{J}} \otimes b$$

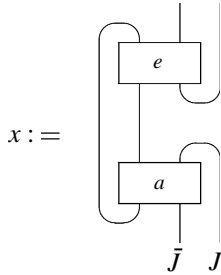
depicted in Fig. 5, are bijections.

2. For all $a \in A, b \in B$ we have

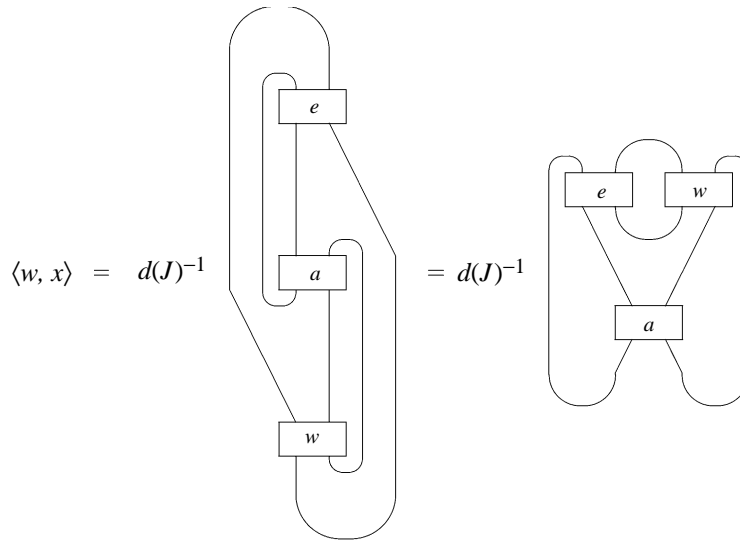
$$\Phi_1(a \otimes b) = \langle a_{(2)}, b_{(1)} \rangle \Phi_2(a_{(1)} \otimes b_{(2)}). \quad (6.4)$$

Proof. 1. Since \mathcal{E} is semisimple and J is the only simple $\mathfrak{B} - \mathfrak{A}$ -morphism up to isomorphism, the composition $\otimes : \text{Hom}(J\bar{J}J, J) \otimes \text{Hom}(J, J\bar{J}J) \rightarrow \text{End}(J\bar{J}J)$ is an isomorphism. Combining this with the isomorphisms $\text{End}(J\bar{J}) \cong \text{Hom}(J\bar{J}J, J)$, etc., provided by the spherical structure, this easily implies that Φ_1, Φ_2 are isomorphisms.

2. For $a, e \in A$ we define $x \in B$ by



and claim that $x = d(J)^{-1} \text{Tr}_A(ea_{(1)})\mathcal{F}(a_{(2)})$. To prove this, we compute



$$= d(J)^{-1} \text{Tr}_A((e \star w)a) = \langle a, \mathcal{F}(e \star w) \rangle = \langle a, \mathcal{F}(e)\mathcal{F}(w) \rangle$$

$$= \langle a_{(1)}, \mathcal{F}(e) \rangle \langle a_{(2)}, \mathcal{F}(w) \rangle = d(J)^{-1} \text{Tr}_A(ea_{(1)})\langle w, \mathcal{F}(a_{(2)}) \rangle.$$

The equality of the two diagrams follows from a simple computation using the axioms of a spherical category, which we omit. The claim now holds by non-degeneracy of

the pairing $\langle \cdot, \cdot \rangle$. Using our formula for x we find

$$\begin{array}{c} \text{Tr}_C \begin{array}{c} \begin{array}{|c|} \hline f \\ \hline \end{array} \\ \begin{array}{|c|} \hline e \\ \hline \end{array} \\ \begin{array}{|c|} \hline b \\ \hline \end{array} \\ \begin{array}{|c|} \hline a \\ \hline \end{array} \\ \hline \begin{array}{ccc} J & \bar{J} & J \end{array} \end{array} = \text{Tr}_C \begin{array}{c} \begin{array}{|c|} \hline e \\ \hline \end{array} \\ \begin{array}{|c|} \hline b \\ \hline \end{array} \\ \begin{array}{|c|} \hline a \\ \hline \end{array} \\ \begin{array}{|c|} \hline f \\ \hline \end{array} \\ \hline \begin{array}{ccc} J & \bar{J} & J \end{array} \end{array} = \begin{array}{c} \begin{array}{|c|} \hline x \\ \hline \end{array} \quad \begin{array}{|c|} \hline b \\ \hline \end{array} \\ \diagdown \quad \diagup \\ \begin{array}{|c|} \hline f \\ \hline \end{array} \\ \diagup \quad \diagdown \\ \begin{array}{|c|} \hline \text{loop} \\ \hline \end{array} \end{array} \end{array} \quad (6.5)$$

$$= d(J)^{-1} \text{Tr}_A(ea_{(1)}) \begin{array}{c} \begin{array}{|c|} \hline a_{(2)} \quad b \\ \hline \end{array} \\ \diagdown \quad \diagup \\ \begin{array}{|c|} \hline f \\ \hline \end{array} \\ \diagup \quad \diagdown \\ \begin{array}{|c|} \hline \text{loop} \\ \hline \end{array} \end{array} = \text{Tr}_A(ea_{(1)}) \langle a_{(2)} \hat{\mathcal{F}}(f), b \rangle$$

$$\begin{aligned}
 &= \text{Tr}_A(ea_{(1)}) \langle a_{(2)}, b_{(1)} \rangle \langle \hat{\mathcal{F}}(f), b_{(2)} \rangle \\
 &= d(J)^{-1} \text{Tr}_A(ea_{(1)}) \text{Tr}_B(fb_{(2)}) \langle a_{(2)}, b_{(1)} \rangle.
 \end{aligned}$$

Since J is simple, the partial trace $\text{Tr}_J(a) \in \text{End } \bar{J}$ of $a \in \text{End } J\bar{J}$ is given by $d(J)^{-1} \text{Tr}_A(a) \text{id}_{\bar{J}}$, which implies

$$\begin{array}{c} \langle a_{(2)}, b_{(1)} \rangle \text{Tr}_C \begin{array}{c} \begin{array}{|c|} \hline f \\ \hline \end{array} \\ \begin{array}{|c|} \hline e \\ \hline \end{array} \\ \begin{array}{|c|} \hline a_{(1)} \\ \hline \end{array} \\ \begin{array}{|c|} \hline b_{(2)} \\ \hline \end{array} \\ \hline \begin{array}{ccc} J & \bar{J} & J \end{array} \end{array} = d(J)^{-1} \langle a_{(2)}, b_{(1)} \rangle \text{Tr}_A(ea_{(1)}) \text{Tr}_B(fb_{(2)}).
 \end{array}$$

Comparing this with (6.5), we see that multiplying both sides of (6.4) with $\Phi_1(e \otimes f) = id_{\bar{J}} \otimes f \circ e \otimes id_J$ and taking traces we obtain an identity for all $e \in A, f \in B$. In view of 2 and the non-degeneracy of the trace in a spherical category, we conclude that (6.4) holds for all $a \in A, b \in B$. \square

Proposition 6.19. *Let \mathcal{E} be a Morita context which is semisimple, irreducible and has depth two. Let $\varepsilon, \hat{\varepsilon}, \Delta, \hat{\Delta}$ be defined as above. Then*

1. $\varepsilon, \hat{\varepsilon}$ are multiplicative.
2. $\Delta, \hat{\Delta}$ are multiplicative.
3. S, \hat{S} are coinverses, i.e. $m(S \otimes id)\Delta = m(id \otimes S)\Delta = \eta\varepsilon$, etc.
4. A and B are semisimple Hopf algebras in duality, and C is the Weyl algebra of A , cf. e.g. [51].

Proof. 1. Let $a, b \in A$. Since J is simple, $Hom(\mathbf{1}_{\mathfrak{A}}, J\bar{J})$ is one dimensional and we have

$$a \circ \varepsilon(J) = d(J)^{-1}(\overline{\varepsilon(J)}) \circ a \circ \varepsilon(J) \varepsilon(J) = \langle a, \hat{1} \rangle \varepsilon(J) = \varepsilon(a) \varepsilon(J).$$

(There should be no danger of confusion between the duality morphisms $\varepsilon(J), \overline{\varepsilon(J)}$, which are part of the spherical structure, and the counits $\varepsilon, \hat{\varepsilon}$ of A and B .) Thus, $\varepsilon(ab)\varepsilon(J) = (ab)\varepsilon(J) = a(b\varepsilon(J)) = a\varepsilon(b)\varepsilon(J) = \varepsilon(a)\varepsilon(b)\varepsilon(J)$, and ε is multiplicative.

2. Let $a, b \in A$. Using Lemma 6.18 we compute

$$\begin{aligned} \text{Diagram 1} &= \langle a_{(2)}, b_{(1)} \rangle \text{Diagram 2} = \langle a_{(2)}, b_{(1)} \rangle \hat{\varepsilon}(b_{(2)}) \text{Diagram 3} \\ &= \langle a_{(2)}, b \rangle \text{Diagram 4} \end{aligned} \quad (6.6)$$

In an entirely analogous fashion one shows

$$\text{Diagram 5} = \langle a, b_{(1)} \rangle \text{Diagram 6} \quad (6.7)$$

Let now $a, b \in A, c, d \in B$. We compute

$$\begin{aligned}
 \langle \Delta(ab), c \otimes d \rangle &= \langle ab, cd \rangle = d(J)^{-1} \text{Tr} \quad \begin{array}{c} \text{---} \\ | \\ \boxed{c} \\ | \\ \boxed{d} \\ | \\ \boxed{a} \\ | \\ \boxed{b} \\ | \\ \text{---} \end{array} = d(J)^{-1} \langle a_{(2)}, d_{(1)} \rangle \text{Tr}_J \quad \begin{array}{c} \text{---} \\ | \\ \boxed{c} \\ | \\ \boxed{a_{(1)}} \\ | \\ \boxed{d_{(2)}} \\ | \\ \boxed{b} \\ | \\ \text{---} \end{array} \\
 &= d(J)^{-1} \langle a_{(2)}, d_{(1)} \rangle \langle a_{(1)}, c_{(1)} \rangle \langle b_{(2)}, d_{(2)} \rangle \text{Tr}_J \quad \begin{array}{c} \text{---} \\ | \\ \boxed{c_{(2)}} \\ | \\ \boxed{b_{(1)}} \\ | \\ \text{---} \end{array} \\
 &= \langle a_{(2)}, d_{(1)} \rangle \langle a_{(1)}, c_{(1)} \rangle \langle b_{(2)}, d_{(2)} \rangle \langle b_{(1)}, c_{(2)} \rangle \\
 &= \langle a_{(1)} b_{(1)}, c \rangle \langle a_{(2)} b_{(2)}, d \rangle = \langle \Delta(a) \Delta(b), c \otimes d \rangle.
 \end{aligned}$$

(The first and sixth equality hold by definition of Δ and $\hat{\Delta}$, respectively. The second and fifth are just the definition of $\langle \cdot, \cdot \rangle$. The third equality follows by Lemma 6.18 and the fourth is due to (6.6) and (6.7).) Since this holds for all c, d , we conclude by duality that $\Delta(ab) = \Delta(a)\Delta(b)$, as desired.

3. Appealing once more to Lemma 6.18, we have

$$\text{Tr}_J \quad \begin{array}{c} \text{---} \\ | \\ \boxed{b} \\ | \\ \boxed{a} \\ | \\ \text{---} \end{array} = \langle a_{(2)}, d_{(1)} \rangle \text{Tr}_J \quad \begin{array}{c} \text{---} \\ | \\ \boxed{a_{(1)}} \\ | \\ \boxed{b_{(2)}} \\ | \\ \text{---} \end{array} \quad \forall a \in A, b \in B.$$

The left-hand side equals

$$\varepsilon(a) \text{Tr}_J \quad \begin{array}{c} \text{---} \\ | \\ \boxed{b} \\ | \\ \text{---} \end{array} = d(J) \varepsilon(a) \hat{\varepsilon}(b).$$

The right-hand side equaling $d(J) \langle a_{(2)}, b_{(1)} \rangle \langle a_{(1)}, \hat{S}(b_{(2)}) \rangle$, we obtain

$$\varepsilon(a) \hat{\varepsilon}(b) = \langle a_{(2)}, b_{(1)} \rangle \langle a_{(1)}, \hat{S}(b_{(2)}) \rangle = \langle a, \hat{S}(b_{(2)}) b_{(1)} \rangle,$$

which is equivalent to $\hat{S}(b_{(2)})b_{(1)} = \hat{\varepsilon}(b)\hat{1}$. Since \hat{S} is involutive, we find $\hat{S}(b_{(1)})b_{(2)} = \hat{\varepsilon}(b)\hat{1}$. The other identities are proved similarly.

4. The first statement just summarizes our results so far, and the second is obvious by definition of the Weyl algebra [51]. \square

Given a semisimple irreducible depth two Morita context, the preceding theorem provides us with a pair A, B of mutually dual Hopf algebras. It remains to relate these Hopf algebras to the categorical structure. Here we use Tannaka theory for Hopf algebras together with a result from [49]. The strategy is (i) to construct a faithful tensor functor $E: \mathcal{A} \rightarrow \text{Vect}_{\mathbb{F}}$, (ii) to deduce that \mathcal{A} is monoidally equivalent to $H - \text{mod}$ for some Hopf algebra H and (iii) to prove that $H \cong A$.

Theorem 6.20. *Let \mathcal{E} be a semisimple spherical irreducible depth two Morita context. Consider the full tensor subcategories*

$$\mathcal{A}_0 \subset \mathcal{A} = \text{END}_{\mathcal{E}}(\mathfrak{A}), \quad \mathcal{B}_0 \subset \mathcal{B} = \text{END}_{\mathcal{E}}(\mathfrak{B})$$

consisting of the tensor powers of $J\bar{J}$ and $\bar{J}J$, respectively, and their retracts. Then we have the equivalences

$$\mathcal{A}_0 \cong A^{\text{cop}} - \text{mod}, \quad \mathcal{B}_0 \cong B^{\text{cop}} - \text{mod}$$

of spherical tensor categories, where A, B are the Hopf algebras constructed in Proposition 6.19.

Proof. Let $Q = J\bar{J}$ and consider the Frobenius algebra (Q, v, v', w, w') with $v = \varepsilon(J)$, $w' = id_J \otimes \bar{\varepsilon}(\bar{J}) \otimes id_{\bar{J}}$, etc. In particular, (Q, w', v) is a monoid in \mathcal{A}_0 . By the depth 2 property, we have $\bar{J}X \cong d(X)\bar{J}$ for all $X \in \mathcal{A}_0$, i.e. there are morphisms $r_i: \bar{J} \rightarrow \bar{J}X$, $r'_i: \bar{J}X \rightarrow \bar{J}$ satisfying the usual conditions. Thus, the morphisms $s_i = id_J \otimes r_i: Q = J\bar{J} \rightarrow J\bar{J}X = QX$ establish an isomorphism $QX \cong d(X)Q$. One easily verifies that the s_i are Q -module morphisms, i.e. satisfy $s_i \circ w' = w' \otimes id_X \circ id_Q \otimes s_i$, and similarly the $s'_i = id_J \otimes r'_i$. These facts imply that the functor $E: \mathcal{A}_0 \rightarrow \text{Vect}_{\mathbb{F}}$ defined by $E(X) = \text{Hom}(\mathbf{1}, Q \otimes X)$ and $E(s)\phi = (id_Q \otimes s) \circ \phi$ for $s: X \rightarrow Y$ is a faithful (strong) tensor functor, where the isomorphisms $d_{X,Y}: E(X) \boxtimes E(Y) \rightarrow E(X \otimes Y)$ are given by $d_{X,Y}(\phi \boxtimes \psi) = w' \otimes id_{X \otimes Y} \circ id_Q \otimes \phi \otimes id_Y \circ \psi$ for $\phi \in E(X), \psi \in E(Y)$. See [49, Section 3] for the details. It then follows from Tannaka theory there exists a finite-dimensional Hopf algebra H and an equivalence $F: \mathcal{A}_0 \rightarrow H - \text{mod}$ such that $K \circ F = E$, where $K: H - \text{mod} \rightarrow \mathbb{F}\text{-Vect}$ is the forgetful functor. Here $H = \text{Nat } E$ is the \mathbb{F} -algebra of natural transformations from E to itself. Consider the map α which to $a \in \text{End } Q = A$ associates the family $\alpha(a) = \{\alpha(a)_X \in \text{End } E(X), X \in \mathcal{A}_0\}$, where $\alpha(a)_X \phi = (a \otimes id_X) \circ \phi$ for $\phi \in E(X)$. It is obvious that $\alpha(a) \in \text{Nat } E =: H$ and that $\alpha: A \rightarrow H$ is an algebra homomorphism. Semisimplicity of \mathcal{A}_0 and finite dimensionality of H imply that α is an isomorphism, which we now suppress. It remains to show that the coproduct Δ' of H provided by Tannaka theory coincides with the one constructed in Proposition 6.19. By definition

of $\Delta'(a) = a_{(1)} \otimes a_{(2)}$, the diagram

$$\begin{array}{ccc} E(X) \otimes E(Y) & \xrightarrow{d_{X,Y}} & E(X \otimes Y) \\ \downarrow a_{(1),X} \otimes a_{(2),Y} & & \downarrow a_{X \otimes Y} \\ E(X) \otimes E(Y) & \xrightarrow{d_{X,Y}} & E(X \otimes Y) \end{array}$$

commutes for all $a \in A$ and $X, Y \in \mathcal{A}_0$. In view of the definition of $d_{X,Y}$, this is equivalent to $a \circ w' = w' \circ a_{(2)} \otimes a_{(1)}$ for all $a \in A = \text{End } Q$. For arbitrary $b, c \in A$, this implies

$$\text{Tr}_Q(a \circ w' \circ b \otimes c \circ w) = \text{Tr}_Q(w \circ a_{(2)} \otimes a_{(1)} \circ b \otimes c \circ w).$$

Consistent with previous terminology we write $b \star c = w' \circ b \otimes c \circ w \in A$ for $b, c \in A$, and the fact that Q contains 1 with multiplicity one implies $\text{Tr}_Q(b \star c) = d(J)^{-1} \text{Tr}_Q(b) \text{Tr}_Q(c)$. Therefore,

$$\text{Tr}_Q(a(b \star c)) = d(J)^{-1} \text{Tr}_Q(a_{(2)}b) \text{Tr}_Q(a_{(1)}c) \quad \text{where } a_{(1)} \otimes a_{(2)} = \Delta'(a).$$

On the other hand, the definition $\langle \Delta(a), b \otimes c \rangle = \langle a, bc \rangle$ of Δ as given above satisfies

$$\text{Tr}_Q(a(b \star c)) = d(J)^{-1} \text{Tr}_Q(a_{(1)}b) \text{Tr}_Q(a_{(2)}c) \quad \text{where } a_{(1)} \otimes a_{(2)} = \Delta(a).$$

Thus $\Delta' = \Delta^{\text{cop}}$, and we are done. \square

We briefly recall some facts concerning the (left) regular representation $Q_l \in A - \text{mod}$ of a semisimple Hopf algebra A . We have $Q_l \cong \bigoplus_X d(X)X$, where the X are the irreducible representations, and therefore $\dim \text{Hom}(X, Q_l) = d(X)$ for all simple X . Furthermore, the regular representation is absorbing: $X \otimes Q_l \cong Q_l \otimes X \cong d(X)Q_l$ for every $X \in A - \text{mod}$.

Proof of Theorem 6.15. By Barrett and Westburry [5], the category $H - \text{mod}$ is spherical and by the coherence theorem [5] we may consider $H - \text{mod}$ as strict monoidal and strict spherical. By Theorem 6.12 we have a strongly separable and irreducible Frobenius algebra Q in $H - \text{mod}$, which we can normalize such that $\lambda_1 = \lambda_2$. Since Q is irreducible, by Proposition 5.1 the same is true for the Morita context \mathcal{E} of Theorem 3.11. By Theorem 5.12 there thus is a spherical structure on \mathcal{E} extending that of \mathcal{A} . The claim now follows from Theorem 6.20 and the fact $H^{\text{op}, \text{cop}} \cong H$, provided we can show that \mathcal{E} has depth 2.

By definition of \mathcal{E} , every $Y: \mathfrak{B} \rightarrow \mathfrak{A}$ is a retract of XJ for some $X \in \text{END}(\mathfrak{A}) \simeq \mathcal{A}$. By semisimplicity is thus sufficient to show that XJ is a multiple of J for every simple $X \in \mathcal{A}$. We have

$$\text{Hom}(J, XJ) \cong \text{Hom}(J\bar{J}, X) \cong \text{Hom}_{\mathcal{A}}(Q, X),$$

$$\text{End}(XJ) \cong \text{Hom}(XJ\bar{J}, X) \cong \text{Hom}_{\mathcal{A}}(XQ, X).$$

By the properties of Q recalled above, we have $\dim \operatorname{Hom}(J, XJ) = d(X)$, implying that XJ contains J with multiplicity $d(X)$, thus $\operatorname{End}(XJ)$ contains the matrix algebra $M_{d(X)}(\mathbb{F})$. In view of $\dim \operatorname{End}(XJ) = \dim \operatorname{Hom}(XQ, X) = \dim \operatorname{Hom}(d(X)Q, X) = d(X)^2$ we conclude $\operatorname{End}(XJ) \cong M_{d(X)}(\mathbb{F})$ and therefore $XJ \cong d(X)J$ as desired. \square

Remark 6.21. If \mathcal{E} in Proposition 6.19 is a $*$ -bicategory then A, B come with canonical $*$ -operations. It is then not difficult to show that ε, Δ, S are $*$ -homomorphisms, thus A and B are Hopf $*$ -algebras. (E.g. the property $\varepsilon(\overline{J}) = \varepsilon(J)^*$ immediately implies $\varepsilon(a) = \varepsilon(a^*)$.) In the Theorems 6.20 and 6.15 we then have equivalences of tensor $*$ -categories. We omit the proofs.

6.4. Subfactors

The entire analysis of this paper is motivated by the mathematical structures which are implicit in subfactor theory. In this subsection we make the link between subfactor theory and our categorical setting explicit, shedding light on both subjects. The main aim of this section is in fact to improve the communication between subfactor theorists and category minded people, the only new result being Theorem 6.28. We begin with a very brief definition of the notions we will use. For everything else see any textbook on von Neumann algebras, e.g. [64,67,60] and subfactors [26,16].

A von Neumann algebra (vNa) is a unital subalgebra $M \subset \mathcal{B}(\mathcal{H})$ of the algebra $\mathcal{B}(\mathcal{H})$ of bounded operators on some Hilbert space \mathcal{H} which is closed w.r.t. the hermitian conjugation $x \mapsto x^*$ and w.r.t. weak convergence. Equivalently, by von Neumann's double commutant theorem a vNa is a set $M \subset \mathcal{B}(\mathcal{H})$ which is closed under conjugation and satisfies $M'' = M$, where $S' = \{x \in \mathcal{B}(\mathcal{H}) \mid xy = yx \ \forall y \in S\}$ is the commutant of S . A factor is a vNa with trivial center ($M \cap M' = \mathbb{C}\mathbf{1}$) and if N, M are factors such $N \subset M$ then N is called a subfactor. (By abuse of notation ‘subfactor’ occasionally refers to the inclusion $N \subset M$.) Every factor M is of one of the types I, II or III, where M is of type I iff $M \simeq \mathcal{B}(\mathcal{H})$ for some Hilbert space \mathcal{H} . (Every finite-dimensional factor is of type I.) If $N \subset M$ are both of type I then also $M \cap N'$ is of type I and $M \simeq N \bar{\otimes} (M \cap N')$. Under this isomorphism the embedding $N \hookrightarrow M$ becomes $x \mapsto x \bar{\otimes} \mathbf{1}$, and nothing more of interest is to be said. In our discussion of the remaining cases we restrict ourselves to $vNas$ on a separable Hilbert space which simplifies the definitions. Then a factor M is of type III iff every orthogonal projection $e = e^2 = e^* \in M$ is the range of some isometry $v \in M$, i.e. $vv^* = e$, $v^*v = \mathbf{1}$. A factor M which is neither types I or III is of type II, of which there are two subclasses: II_1 and II_∞ . A factor is of type II_1 iff it admits a tracial state (trace, for short) $\operatorname{tr}: M \rightarrow \mathbb{C}$, i.e. a weakly continuous linear functional which is positive ($A > 0 \Rightarrow \operatorname{tr} A > 0$), normalized ($\operatorname{tr}(\mathbf{1}) = 1$) and vanishes on commutators. (It follows that every isometry in a type II_1 factor must be unitary.) A type II_∞ factor is isomorphic to the tensor product of some II_1 factor with $\mathcal{B}(\mathcal{H})$ where $\dim \mathcal{H} = \aleph_0$. The tensor product of any factor with a type III factor is of type III.

In the literature on subfactors the focus has been on type II_1 factors, which are technically easiest to deal with thanks to the existence of a trace, cf. the textbooks [26,16]. Yet, in our discussion we concentrate on the type III case, the technical aspects

of which have been clarified in particular in the work of Longo [39,40]. Those aspects of subfactor theory [40,23] which most directly inspired the present investigation and [47] were in fact done in the type III setting. Anyway, by Popa’s results [57] the classifications of amenable inclusions of hyperfinite types II_1 and III_1 factors amount to the same thing.

The following is implicit in much of the literature on type III subfactors and explicit in [42, Section 7].

Definition 6.22. We denote by \mathcal{T} the 2-category whose objects are type III factors with separable predual. The 1-morphisms are normal unital $*$ -homomorphisms with the obvious composition. For parallel 1-morphisms $\rho, \sigma : M \rightarrow N$ the 2-morphisms are given by

$$\text{Hom}_{\mathcal{T}}(\rho, \sigma) = \{s \in N \mid s\rho(x) = \sigma(x)s \ \forall x \in M\}.$$

The vertical composition of 2-morphisms is multiplication in N and the horizontal composite

is given by

$$s \times t = t\rho_2(s) = \sigma_2(s)t : \rho_2\rho_1 \rightarrow \sigma_2\sigma_1.$$

Lemma 6.23. The 2-category \mathcal{T} has direct sums of 1-morphisms idempotent 2-morphisms split. All identity 1-morphisms are simple.

Proof. Let $\rho, \sigma : M \rightarrow N$. Pick an orthogonal projection $e \in N$, put $f = 1 - e$ and choose isometries $p, q \in N$ such that $pp^* = e, qq^* = f$. Then

$$\eta(\cdot) = p\rho(\cdot)p^* + q\sigma(\cdot)q^*$$

is a direct sum. Let $\rho : M \rightarrow N$ and $e = e^2 = e^* \in \text{End}_{\mathcal{T}}(\rho) \subset N$. Picking an isometry $q \in N$ such that $qq^* = e$ and setting

$$\sigma(\cdot) = q\rho(\cdot)q^*$$

one obviously has $q \in \text{Hom}_{\mathcal{T}}(\sigma, \rho)$. The last claim follows from factoriality. \square

Remark 6.24. This lemma fails for finite factors, which is why one works with bi-modules in the type II_1 case.

Since the kernel of a normal $*$ -homomorphism is a closed two-sided ideal and a type III factor with separable predual is simple, all 1-morphisms in \mathcal{T} are injective. Thus, a

morphism $\rho: N \rightarrow M$ provides an isomorphism of N with the subalgebra $\rho(N) \subset M$. Now Longo's main result in [39] can be rephrased as follows:

Proposition 6.25. *Let $\rho: N \rightarrow M$ be a 1-morphism in \mathcal{T} . The ρ has a two-sided adjoint $\bar{\rho}: M \rightarrow N$ in \mathcal{T} iff the index $[M : \rho(N)]$ is finite.*

Proof. We identify N with the subalgebra $\rho(N) \cong N$ of M , but we insist on writing an embedding map $\iota: N \hookrightarrow M$. Longo proves the following: the index $[M : N]$ of $N \subset M$ is finite iff there is a triple (γ, v, w) where γ is a normal $*$ -endomorphism of M with that $\rho(M) \subset N$ such that there are isometries $v \in M, w \in N$ satisfying

$$v \in \text{Hom}_{\mathcal{T}}(id_M, \gamma), \quad w \in \text{Hom}_{\mathcal{T}}(id_N, \gamma \upharpoonright N), \quad (6.8)$$

$$v^* w = \gamma(v^*) w = c \mathbf{1}. \quad (6.9)$$

Then $c = [M : N]^{-1/2}$, and we refer to [39] for the original definition of the index $[M : N]$. In order to translate Longo's result into categorical language we write $\bar{\iota} = \gamma$ and observe that γ maps M into N , so that only $\iota \circ \bar{\iota}$ gives an morphism of M . On the other hand we see $\gamma \upharpoonright N = \bar{\iota} \circ \iota$. Now (6.8) becomes

$$v \in \text{Hom}_{\mathcal{T}}(id_M, \bar{\iota}), \quad w \in \text{Hom}_{\mathcal{T}}(id_N, \bar{\iota})$$

and v^*, w^* are morphisms in the opposite directions. Finally in view of the definition of the horizontal composition of 2-morphisms in \mathcal{T} the equations (6.9) and their $*$ -conjugates are—up to a numerical factor which can be absorbed in v or w —the four triangular equations which make $\iota, \bar{\iota}$ two-sided duals. If we normalize v, w such that $v^* v = w^* w = [M : N]^{1/2}$ then $d(\iota) = d(\bar{\iota}) = [M : N]^{1/2}$. \square

Lemma 6.26. *Every inclusion of type III factors with finite index defines a 2- $*$ -category $\mathcal{T}_{N \subset M}$ which is a Morita context. The dimension of \mathcal{T} (which is well defined by Proposition 5.17) is finite iff the subfactor has finite depth.*

Proof. $\mathcal{T}_{N \subset M}$ is the subcategory of \mathcal{T} whose objects are $\{N, M\}$ and whose 1-morphisms are generated by $\iota, \bar{\iota}$. More precisely, $\text{Hom}_{\mathcal{T}_{N \subset M}}(N, M)$ is the replete full subcategory of $\text{Hom}_{\mathcal{T}}(N, M)$ whose morphisms are retracts of some $\iota(\bar{\iota})^N$, $N \in \mathbb{N} \cup \{0\}$, and similarly for the other categories of 1-morphisms. In this category the functors $- \otimes \iota$, etc. are clearly dominant, thus $\mathcal{T}_{N \subset M}$ is a Morita context for the tensor categories $\text{End}_{\mathcal{T}_{N \subset M}}(N), \text{End}_{\mathcal{T}_{N \subset M}}(M)$. The $*$ -involution obviously it the $*$ -operation of the algebras. The last claim follows from Proposition 5.17 and the definition according to which $N \subset M$ has finite depth iff the powers of $\bar{\iota}$ contain finitely many simple $M - M$ morphisms up to equivalence. \square

Remark 6.27. 1. In subfactor theory the dimension of \mathcal{T} is called the global index (as opposed to the index $[M : N] = d(Q)$).

2. $\mathcal{T}_{N \subset M}$ having a $*$ -structure it can be made into a spherical category, though not in a completely unique way. See Section 2.4.

Now, given an endomorphism γ of finite index of a type III factor M it is natural to ask whether there is a subfactor $N \subset M$ such that $\gamma = v\bar{v}$, where \bar{v} is a two-sided conjugate of the embedding morphism v . The answer, given in [40], is positive iff there are isometries

$$v \in \text{Hom}(id_M, \gamma), \quad w \in \text{Hom}(\gamma, \gamma^2)$$

satisfying

$$w^2 = \gamma(w)w, \tag{6.10}$$

$$ww^* = \gamma(w^*)w, \tag{6.11}$$

$$v^*w = \gamma(v)^*w = c\mathbf{1}, \quad c \in \mathbb{C}^*. \tag{6.12}$$

(It turned out that (6.11) is redundant, cf. [42].) Then the subfactor is given by

$$N = w^*\gamma(M)w. \tag{6.13}$$

Eqs. (6.10)–(6.12) together with the requirement that v, w are isometries are, of course, saying precisely that (γ, v, v^*, w, w^*) is a strongly separable Frobenius algebra in $\text{End } M$. In a sense, this entire paper is about finding a categorical analog for the simple formula (6.13), which turned out to be more tedious than one might expect. This is precisely due to the fact that as seen above subfactor theory comes with a rich and beautiful inherent categorical structure which we had to model by Theorem 3.11. The reward for our work is the following result.

Theorem 6.28. *Let M be a type III factor, let $\gamma \in \text{End}(M)$ satisfy (6.10)–(6.12) and let \mathcal{A} be the replete full subcategory with subobjects of $\text{End}(M)$ generated by γ . Let $\mathcal{T}_{N \subset M}$ be the bicategory associated with the subfactor $N \subset M$, where N is given by (6.13). (Obviously, $\mathcal{A} = \text{HOM}_{\mathcal{T}}(\mathcal{A}, \mathcal{A})$.) If \mathcal{E} is obtained from (\mathcal{A}, γ) by Theorem 3.11 then have an equivalence of bicategories $\mathcal{E} \simeq \mathcal{T}_{N \subset M}$.*

Proof. By Lemma 6.26 the 2-category $\mathcal{T}_{N \subset M}$ is a Morita context. Thus the claim follows directly from Proposition 4.5. \square

Remark 6.29. 1. The importance of this theorem is that it allows us to compute the 2-category $\mathcal{T}_{N \subset M}$ (up to equivalence) from the data (\mathcal{A}, γ) without explicitly working with subfactors. In the case where \mathcal{A} is the subcategory generated by a Frobenius object γ in some $\text{End } M$ this may seem a rather complicated detour. Yet, we have gained two things. On one hand we see that the bicategory associated with a subfactor with finite index is a structure which appears also in other contexts. Equally important is the fact that our constructions work for arbitrary tensor categories \mathcal{A} which are not subcategories of some $\text{End } M$ generated by one object γ , in fact for arbitrary (algebraically closed) ground field. This will be exploited in [47], the results of which seem hard to prove without our machinery.

2. The results of Section 6.3 are related to subfactor theory in a very direct way whenever the Hopf algebra H is a finite-dimensional C^* -Hopf algebra. (This means H is a complex multi-matrix algebra and Δ, ε respect the natural $*$ -operation.) As shown in [76] every finite-dimensional C^* -Hopf algebra H admits an action on a type II_1 factor M . This action is outer, i.e. $(M^H)' \cap M = \mathbb{C}\mathbf{1}$. It has long been known [52] that in this situation the Jones extension M_1 of the subfactor $M^H \subset M$ carries an outer action of \hat{H} . Together with the well-known material in the present section this provides a von Neumann algebraic proof of the weak Morita equivalence $H - \text{mod} \approx \hat{H} - \text{mod}$. From the perspective of this paper this proof is, however, quite unsatisfactory. On the one hand it is restricted to C^* -Hopf algebras, on the other it is rather indirect since it involves infinite dimensional von Neumann algebras. \square

7. Morita invariance of state sum invariants

In this section, we give an interesting and non-trivial application of our notion of weak monoidal Morita equivalence to the study of triangulation (or state sum) invariants of closed 3-manifolds. We begin with a very brief sketch of the works which are relevant to our discussion, apologizing to everyone whose contribution is being glossed over.

In [69] Turaev and Viro used the 6j-symbols of the quantum group $SU_q(2)$ to define a numerical invariant $TV_{SU_q(2)}(M, T)$ for any closed 3-manifold together with a triangulation T . They went on to prove that it does not depend on T and thus gives rise to a topological invariant $TV_{SU_q(2)}(M)$. In [68] this construction was generalized to an invariant $TV(M, \mathcal{C})$ associated with any modular category \mathcal{C} . (Modular categories are braided ribbon categories satisfying a certain non-degeneracy condition.) Recently Gelfand and Kazhdan [20] and Barrett and Westbury [5] defined triangulation invariants for 3-manifolds on the basis of certain tensor categories which are not required to be braided. In our discussion we focus on the invariant of Barrett and Westbury, which we call $BW(M, \mathcal{C})$, since it is based on spherical categories and therefore close in spirit to our work. With appropriate normalizations one has $TV(M, \mathcal{C}) = BW(M, \mathcal{C})$ if \mathcal{C} is modular. The utility of the notion of weak monoidal Morita equivalence is now illustrated by the following.

Theorem 7.1. *Let \mathcal{A}, \mathcal{B} be (strict) semisimple spherical categories with simple unit and finitely many simple objects. If \mathcal{A}, \mathcal{B} are weakly monoidally Morita equivalent and $\dim \mathcal{A} \neq 0$ then we have $BW(M, \mathcal{A}) = BW(M, \mathcal{B})$ for all closed orientable 3-manifolds M .*

Remark 7.2. 1. Before we sketch the proof of this result we point out that it resolves a (minor) puzzle concerning the BW invariant. In [36] Kuperberg had defined a 3-manifold invariant $Ku(M, H)$ for every finite-dimensional Hopf algebra H over an algebraically closed field \mathbb{F} which is involutive ($S^2 = id$) and whose characteristic does not divide the dimension of H . (That these conditions are equivalent to semisimplicity of H and \hat{H} was not known then.) In [3] it was proven that the invariant

BW is a generalization of Ku in the sense that $Ku(M, H) = BW(M, H - \text{mod})$, again assuming appropriate normalizations. For Kuperberg's invariant it had been known that $Ku(M, H) = Ku(M, \hat{H})$, but this becomes obscure when it is expressed in terms of the invariant BW . This puzzle is resolved by Theorem 7.1 together with Corollary 6.16, according to which $H - \text{mod}$ and $\hat{H} - \text{mod}$ are weakly monoidally Morita equivalent.

Example 7.3. The preceding application of weak monoidal Morita equivalence is not really new in that the result $BW(M, H - \text{mod}) = BW(M, \hat{H} - \text{mod})$ can be derived from the connection between the invariants BW and Ku . A less obvious example is provided in the companion paper [47]. There we prove that the center $\mathcal{Z}(\mathcal{C})$ [45, 27, 62], which is the categorical version of Drinfel'd's quantum double, of a semisimple spherical category with non-zero dimension is again spherical and semisimple (and modular in the sense of Turaev). Furthermore, we prove the weak monoidal Morita equivalence $\mathcal{Z}(\mathcal{C}) \approx \mathcal{C} \boxtimes \mathcal{C}^{\text{op}}$, which by the above theorem implies

$$\begin{aligned} BW(M, \mathcal{Z}(\mathcal{C})) &= BW(M, \mathcal{C} \boxtimes \mathcal{C}^{\text{op}}) = BW(M, \mathcal{C}) \cdot BW(M, \mathcal{C}^{\text{op}}) \\ &= BW(M, \mathcal{C}) \cdot BW(-M, \mathcal{C}) \\ &= |BW(M, \mathcal{C})|^2 \text{ if } \mathcal{C} \text{ is a } * - \text{category}. \end{aligned}$$

The relation between a category \mathcal{C} and its quantum double being quite non-trivial we are not aware of a simpler proof of this equality.

Sketch of Proof. The proof relies strongly on ideas of Ocneanu which, unfortunately, found expression only in the unpublished (and unfinished) manuscript [53]. Therefore, the more complete accounts [15, 34] are very useful. In [53] Ocneanu defined a triangulation invariant $Oc(M, A \subset B)$ of 3-manifolds departing from an inclusion $A \subset B$ of type II_1 factors with finite index and finite depth. We recall from Section 6.4 that subfactors $A \subset B$ with finite index give rise to a Morita context \mathcal{E} , whose dimension is finite iff the subfactor has finite depth. (\mathcal{E} is given by bimodules associated with the subfactor or, alternatively, by $*$ -algebra homomorphisms in the type III case.) Ocneanu's invariant is easily seen to depend only on \mathcal{E} and not on other, in particular analytic properties of the subfactor. Furthermore, as is quite evident from [15], it generalizes to any spherical (or $*$ -) Morita context of finite, non-zero dimension over an algebraically closed field. As to the definition of the invariant we only say that one chooses a triangulation T with directed edges and an assignment $V \in \{\mathfrak{A}, \mathfrak{B}\}^V$ of labels $\{\mathfrak{A}, \mathfrak{B}\}$ to the vertices V . Then to every edge of T one assigns an isomorphism class of 1-morphisms in $\text{HOM}_{\mathcal{E}}(\mathfrak{X}, \mathfrak{Y})$, where $\mathfrak{X}, \mathfrak{Y}$ are the labels attached to the initial and terminal vertices of the edge. $Oc(M, \mathcal{E}, T, V)$ is now defined as the sum over the edge labelings of a product of 6j-symbols. Note that there is no summation over the labeling V of the vertices! In fact it is shown in [53, 15] that $Oc(M, \mathcal{E}, T, V)$ depends neither on the labeling V (for fixed triangulation T) nor on T . If one labels all vertices of T with \mathfrak{A} one finds that $Oc(M, \mathcal{E}, T) = BW(M, \text{END}_{\mathcal{E}}(\mathfrak{A}), T)$, i.e. the invariant reduces to the invariant of Barrett and Westbury for the spherical category $\text{END}_{\mathcal{E}}(\mathfrak{A})$. Similarly, by labeling all vertices with \mathfrak{B} one obtains $Oc(M, \mathcal{E}, T) = BW(M, \text{END}_{\mathcal{E}}(\mathfrak{B}), T)$.

By independence of BW and Oc of the triangulation one concludes

$$Oc(M, \mathcal{E}) = BW(M, \text{END}_{\mathcal{E}}(\mathfrak{A})) = BW(M, \text{END}_{\mathcal{E}}(\mathfrak{B})). \quad (7.1)$$

Our claim thus follows if we take \mathcal{E} to be a Morita context for $\mathcal{A} \approx \mathcal{B}$. \square

Remark 7.4. 1. The argument sketched above should of course be spelled out in more detail. In particular, this requires a construction of the TQFT associated with the invariant $BW(\cdot, \mathcal{A})$. (In doing so extreme care is required when the tensor category \mathcal{A} contains simple objects which are self-dual and pseudo-real. Unfortunately, this is neglected in the bulk of the literature on the subject, with the notable exception of [68] and the remarks in [4].) We hope to do this in a future part of this series.

2. Let us emphasize the lesson we draw from the above considerations. In view of (7.1) the invariant $Oc(M, \mathcal{E})$ is already determined by considerably smaller amounts of data, as contained in the tensor categories $\text{END}_{\mathcal{E}}(\mathfrak{A})$ or $\text{END}_{\mathcal{E}}(\mathfrak{B})$. (Therefore, the observation [4,20] that the Turaev–Viro invariant generalizes to tensor categories without braiding could have been made already by the authors of [53,15].) Despite the greater generality of [4,20] there is a lasting significance of the invariant Oc which clearly goes beyond [4,20], viz. precisely the Morita invariance of the invariant BW which we pointed out above.

8. Discussion and outlook

In various places we have already mentioned closely related works by other authors. We summarize these references and comment on several other recent works. The relation between classical Frobenius algebras and Frobenius algebras in $\mathbb{F}\text{-Vect}$ is due to Quinn [58] and Abrams [1,2]. The literature on Frobenius algebras in categories other than Vect is quite small but has begun to grow recently. As mentioned earlier, strongly separable Frobenius algebras in C^* -categories (‘Q-systems’) were first considered in [42], motivated by subfactor theory. In an algebraic-topological context commutative Frobenius algebras in symmetric tensor categories appear in [63]. The relation between Frobenius algebras and two-sided duals (only in \mathcal{CAT} , though) is hinted at in [32] but not developed very far. The discussion in [65, Section 3.3] has some relations to our work, and [66, Section 3] has some overlap with our Section 6.3. Note, however, that in these references most proofs are omitted, and the discussion in [66] is limited to C^* -bicategories. In [33], module categories of ‘rigid \mathcal{C} -algebras’ in braided tensor categories are considered with the aim of categorifying the considerations on modular invariants in [8]. In view of Proposition 5.13, rigid \mathcal{C} -algebras are nothing but Frobenius algebras. (Since the Frobenius algebras are assumed to be commutative (i.e. $w' \circ c(Q, Q) = w'$) only ‘type I’ modular invariants are covered by this analysis.) Similar matters are considered in somewhat greater generality in [18], where Frobenius algebras in the sense of Definition 3.1 appear explicitly, influenced by a talk of the author. A recent construction by Yamagami [73] bears some relation to our construction of the bicategory \mathcal{E} . Given a tensor category \mathcal{A} and a full subcategory $\mathcal{A}_0 \simeq H\text{-mod}$, he constructs a bicategory which seems to be equivalent to our \mathcal{E} in the special case

of Section 6.2, where the Frobenius algebra arises from the regular representation of a semisimple and cosemisimple Hopf algebra H .

When the present work was essentially finished we learned that the definition of the bicategory \mathcal{E} via (bi)modules was discovered independently by Yamagami. Furthermore, in a very interesting sequel [54] of [33], Ostrik considers modules \mathcal{M} over a tensor category \mathcal{C} . (Note that, given a bicategory \mathcal{E} with $\mathfrak{A}, \mathfrak{B} \in \text{Obj } \mathcal{E}$, the categories $\text{HOM}_{\mathcal{E}}(\mathfrak{A}, \mathfrak{B}), \text{HOM}_{\mathcal{E}}(\mathfrak{B}, \mathfrak{A})$ are left and right, respectively, modules over $\mathcal{A} = \text{END}_{\mathcal{E}}(\mathfrak{A})$.) He shows that every module over \mathcal{C} arises from an algebra A in \mathcal{C} and constructs a dual algebra \mathcal{C}^* . The latter is equivalent to our $\mathcal{B} = \text{END}_{\mathcal{E}}(\mathfrak{A})$. He also considers applications to modular invariants and succeeds in phrasing most results of [8] in categorical terms, albeit without many proofs.

Results like Theorem 7.1 lead us to believe that all existing (and future) applications of subfactor theory to low-dimensional topology ‘factor through category theory’—as is by now well known for the knot invariants of Jones and HOMFLY. More generally, we are convinced that essentially all algebraic aspects and results of subfactor theory (at finite index) permit generalization to a considerably wider categorical setting. This is further vindicated by the subsequent parts of this series whose main results we briefly outline.

As already mentioned, in part II [47] we prove that the center $\mathcal{Z}(\mathcal{C})$ of a finite semisimple spherical tensor category \mathcal{C} of non-zero dimension is weakly monoidally Morita equivalent to $\mathcal{C} \boxtimes \mathcal{C}^{\text{op}}$. Furthermore, it is a modular category in the sense of [68]. In view of the relation [30] between the categorical and the Hopf algebraic quantum double this should be interpreted as a generalization of the fact [13] that quantum doubles of nice Hopf algebras have modular representation categories.

In Part III [48] we will consider the bicategory $\tilde{\mathcal{E}}$ mentioned in Remark 3.18. It will be shown to satisfy the assumptions of Theorem 3.17, which implies its equivalence with \mathcal{E} . In particular, we have $\mathcal{A} \approx \mathcal{B}$ iff there exists a Frobenius algebra Q in \mathcal{A} such that $\mathcal{B} \cong {}^{\mathcal{Q}}\text{-mod-}Q$. As mentioned in Remark 3.18, the latter implies the equivalence of the braided tensor categories $\mathcal{L}_1(\mathcal{A}) \simeq \mathcal{L}_1(\mathcal{B})$. It is natural to ask whether the converse is true.

The programme of identifying the connections between subfactor theory (at finite index) and category theory is certainly vindicated by the applications to the classification of modular invariants, cf. [33, 18, 54], and to topology, as considered in Section 7 and [47]. The rapprochement of subfactor theory and ‘mainstream’ mathematics which this foreshadows will undoubtedly be helpful also in the classification programme of subfactors.

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