


Partition of random graphs into subgraphs of bounded component order

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Connections in Discrete Mathematics
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^{*}Support from . Based on joint work with Nicolas Broutin (INRIA).

Component colouring

Graph $G = (V, E)$

t-component set = subset of V , no induced component more than t vertices

t-component colouring = partition of V into *t*-comp. sets



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$t = 1 \Rightarrow$ ordinary chromatic number

larger $t \Rightarrow$ “more error” allowed and possibly fewer colours needed

Warm up

Graph $G = (V, E)$

t-component set = subset of V , no induced component of more than t vertices

t-component colouring = partition of V into *t*-comp. sets

t-component chromatic number = least number of parts in *t*-comp. colouring
denoted $\chi_c^t(G)$

$$\frac{\chi(G)}{t} \leq \chi_c^t(G) \leq \min \left\{ \chi(G), \left\lceil \frac{|V(G)|}{t} \right\rceil \right\}$$

Previous study

Dynamic/evolving databases:

cf. Kleinberg, Motwani, Raghavan, Venkatasubramanian 1997

H-minor free and bounded maximum degree graphs:

- Alon, Ding, Oporowski, Vertigan 2003
- Berke and Szabó 2007
- Kawarabayashi and Mohar 2007
- Linial, Matoušek, Sheffet, Tardos 2008
- Esperet and Joret 2014
- Liu and Oum 2015+

Random graphs $G_{n,p}$

Binomial random graph $G_{n,p}$, championed by Erdős and Rényi 1959/60:

$V(G_{n,p})$: $[n] = \{1, \dots, n\}$

$E(G_{n,p})$: edges included independently with probability $p = p(n)$

Want properties holding *asymptotically almost surely* (a.a.s.), i.e. with prob. $\rightarrow 1$ as $n \rightarrow \infty$

History of $\chi_c^1(G_{n,p})$

Problem of finding $\chi(G_{n,p})$ mentioned by Erdős and Rényi 1959/60
and conjecture made by Grimmett and McDiarmid 1979

Theorem (Bollobás 1988, Matula 1987, Łuczak 1991)

Suppose $np \rightarrow \infty$ and p bounded from 1 as $n \rightarrow \infty$. Letting $b = 1/(1 - p)$,

$$\chi(G_{n,p}) \sim \frac{n}{2 \log_b np} \text{ a.a.s.}$$

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... fixed np remains open, but close to settled: Coja-Oghlan and Vilenchik 2013

A rough view of $\chi_c^t(G_{n,p})$

Proposition (Broutin and Kang)

Suppose $np \rightarrow \infty$ and p bounded from 1 as $n \rightarrow \infty$. Let $b = 1/(1 - p)$.

- If $t = o(\log_b np)$, then $\chi_c^t(G_{n,p}) = \Theta\left(\frac{n}{\log_b np}\right)$ a.a.s.
- If $t = \omega(\log_b np)$ and $t = o(n)$, then $\chi_c^t(G_{n,p}) \sim \frac{n}{t}$ a.a.s.

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Compare with $\chi_c^t(G) \leq \min\left\{\chi(G), \left\lceil \frac{|V(G)|}{t} \right\rceil\right\}$.

Recall $\chi(G_{n,p}) \sim \frac{n}{2 \log_b np}$ a.a.s. and note $\left\lceil \frac{|V(G_{n,p})|}{t} \right\rceil \sim \frac{n}{t}$ if $t = o(n)$.

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I.e. partition into asymptotically fewer than $\chi(G_{n,p})$ parts yields a monochromatic component of at least average part size.

A rough view of $\chi_c^t(G_{n,p})$, p fixed

Slightly better focus when p fixed as $n \rightarrow \infty$:

Fix $0 < p < 1$. Let $b = 1/(1 - p)$.

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- If $t = \Theta(\log_b np)$, then ...
- If $t = \omega(\log_b np)$ and $t = o(n)$, then $\chi_c^t(G_{n,p}) \sim \frac{n}{t}$ a.a.s.

A finer view of $\chi_c^t(G_{n,p})$, p fixed

Theorem (Broutin and Kang)

Fix $0 < p < 1$. Let $b = 1/(1 - p)$.

If $t \sim \tau \log_b np$ and κ is unique root of $\iota(\tau, \kappa) = 0$, then $\chi_c^t(G_{n,p}) \sim \frac{n}{\kappa \log_b np}$ a.a.s.

A finer view of $\chi_c^t(G_{n,p})$, p fixed

For $\tau, \kappa > 0$, define

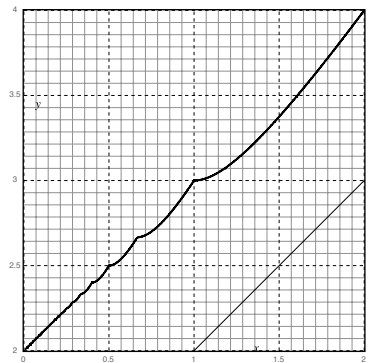
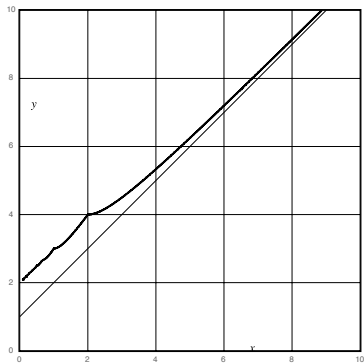
$$\iota(\tau, \kappa) = \frac{1}{2} \left(\left(\kappa - \tau \left\lfloor \frac{\kappa}{\tau} \right\rfloor \right) \left(\kappa - \tau \left\lfloor \frac{\kappa}{\tau} \right\rfloor - \tau \right) - \kappa(\kappa - \tau - 2) \right).$$

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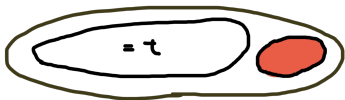
Plot of $\iota(x, y)$



First moment

Non-smooth behaviour governed by expected number of t -component k -sets

$\tau > 2$: two components, one with exactly t vertices



$\tau \leq 2$: more than two components, all but at most one with exactly t vertices



Second moment

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Lemma

If $t \sim \tau \log_b np$, κ is unique root of $\iota(\tau, \kappa) = 0$, and $k \sim (\kappa - \varepsilon) \log_b n$, then $\mathbb{P}(\alpha_c^t(G_{n,p}) < k) \leq \exp(-n^2/(\log n)^5)$.

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Proof of Theorem.

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Start with $S' = [n]$.

While $|S'| > n/(\log n)^2$, extract a t -component k -set from S' as a new colour.

For each of remaining $n/(\log n)^2$ vertices, use a new colour. \square

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Requires an upper bound on

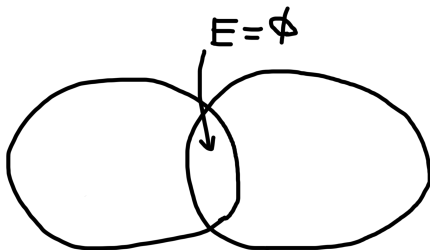
$$\begin{aligned}\Delta &= \sum_{A, B \subseteq [n], 2 \leq |A \cap B| < k} \mathbb{P}(A, B \text{ are } t\text{-comp. } k\text{-sets}) \\ &= \sum_{2 \leq \ell < k} \binom{n}{k} \binom{k}{\ell} \binom{n-k}{k-\ell} p(k, \ell)\end{aligned}$$

where $p(k, \ell)$ is prob. two k -sets, sharing exactly ℓ vertices are t -comp. sets.

Want Δ to be much smaller than squared expected number of t -comp. k -sets.

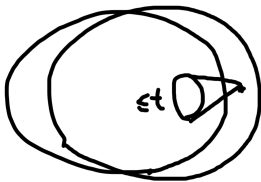
$$\Delta_1$$

When number of common vertices ℓ is small, we reduce to when the graph of the intersection $A \cap B$ has no edges:



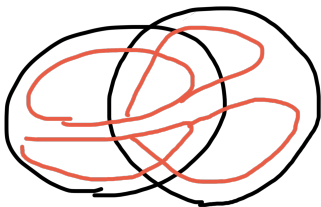
$$\Delta_3$$

When number of common vertices ℓ is large, we use the fact that the maximum degree from $B \setminus A$ to B is at most t if B is a t -component set:



$$\Delta_2$$

An intermediate case consideration needed only when $\tau \geq 2$, so that we basically only need consider when A and B have two components:



Ramsey-type parameter

The *t*-component Ramsey number is the smallest $R^t(k)$ such that any $R^t(k)$ -vertex graph or its complement has a *t*-component *k*-set.

$R^1(k)$ coincides with diagonal 2-colour Ramsey numbers so exponential in *k*, while $R^k(k)$ is trivially *k*.

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Proposition

$$R^{k-\Omega(k)}(k) = 2^{\Omega(k)}. \quad R^{k-\omega(\log k)}(k) = k^{\omega(1)}.$$

Proposition

$$R^{k-O(\log k)}(k) = k^{O(1)}.$$

Further study

Theorem (Broutin and Kang)

Fix $0 < p < 1$. Let $b = 1/(1 - p)$.

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How about the $p = o(1)$ case? (Janson and Thomason 2008)

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What about *edge* partitions? (Bohman, Frieze, Krivelevich, Loh, Sudakov 2011 and Spöhel, Steger, Thomas 2010)

Thank you!